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Sophia Antipolis, France

CR-Form-v4
CHANGE REQUEST
⌘ 35.201 CR ⌘ ev - ⌘ Current version: 4.0.0 ⌘

For **HELP** on using this form, see bottom of this page or look at the pop-up text over the ⌘ symbols.

Proposed change affects: ⌘ (U)SIM ME/UE Radio Access Network Core Network

Title:	⌘ Correct the maximum input message length for f8 and f9		
Source:	⌘ Siemens Atea		
Work item code:	⌘ Security	Date:	⌘ 29 November 2001
Category:	⌘ A	Release:	⌘ REL-4
	Use <u>one</u> of the following categories: F (correction) A (corresponds to a correction in an earlier release) B (addition of feature), C (functional modification of feature) D (editorial modification) Detailed explanations of the above categories can be found in 3GPP TR 21.900 .		Use <u>one</u> of the following releases: 2 (GSM Phase 2) R96 (Release 1996) R97 (Release 1997) R98 (Release 1998) R99 (Release 1999) REL-4 (Release 4) REL-5 (Release 5)

Reason for change:	⌘ The f8 maximum input message length is too low: 5114 shall be enhanced to 20000. There is no message size limitation for f9.
Summary of change:	⌘ Bring TS 35.201 inline with TS 33.105
Consequences if not approved:	⌘ Inconsistent specifications and risk of f8/f9 implementations that take 5114 bit as an upper limit.

Clauses affected:	⌘ 2.1; 2.3; 3; 4		
Other specs affected:	⌘ <input type="checkbox"/> Other core specifications <input type="checkbox"/> Test specifications <input type="checkbox"/> O&M Specifications	⌘	
Other comments:	⌘		

***** First Modification *****

2 Introductory information

2.1 Introduction

Within the security architecture of the 3GPP system there are two standardised algorithms: A confidentiality algorithm f_8 , and an integrity algorithm f_9 . These algorithms are fully specified here. Each of these algorithms is based on the **KASUMI** algorithm that is specified in a companion document[4]. **KASUMI** is a block cipher that produces a 64-bit output from a 64-bit input under the control of a 128-bit key.

The confidentiality algorithm f_8 is a stream cipher that is used to encrypt/decrypt blocks of data under a confidentiality key **CK**. The block of data may be between 1 and 200005444 bits long. The algorithm uses **KASUMI** in a form of output-feedback mode as a keystream generator.

The integrity algorithm f_9 computes a 32-bit MAC (Message Authentication Code) of a given input message using an integrity key **IK**. The approach adopted uses **KASUMI** in a form of CBC-MAC mode.

***** Next Modification *****

2.3 List of Variables

A, B	are 64-bit registers that are used within the f_8 and f_9 functions to hold intermediate values.
BEARER	a 5-bit input to the f_8 function.
BLKCNT	a 64-bit counter used in the f_8 function.
BLOCKS	an integer variable indicating the number of successive applications of KASUMI that need to be performed, for both the f_8 and f_9 functions.
CK	a 128-bit confidentiality key.
COUNT	a 32-bit time variant input to both the f_8 and f_9 functions.
DIRECTION	a 1-bit input to both the f_8 and f_9 functions indicating the direction of transmission (uplink or downlink).
FRESH	a 32-bit random input to the f_9 function.
IBS	the input bit stream to the f_8 function.
IK	a 128-bit integrity key.
KM	a 128-bit constant that is used to modify a key. This is used in both the f_8 and f_9 functions. (It takes a different value in each function).
KS[i]	is the i^{th} bit of keystream produced by the keystream generator.

KSB _{<i>i</i>}	is the <i>i</i> th block of keystream produced by the keystream generator. Each block of keystream comprises 64 bits.
LENGTH	is an input to the <i>f8</i> and <i>f9</i> functions. It specifies the number of bits in the input bitstream (1-5114).
MAC-I	is the 32-bit message authentication code (MAC) produced by the integrity function <i>f9</i> .
MESSAGE	is the input bitstream of LENGTH bits that is to be processed by the <i>f9</i> function.
OBS	the output bit streams from the <i>f8</i> function.
PS	is the input padded string processed by the <i>f9</i> function.
REGISTER	is a 64-bit value that is used within the <i>f8</i> function.

3 Confidentiality algorithm *f8*

3.1 Introduction

The confidentiality algorithm *f8* is a stream cipher that encrypts/decrypts blocks of data between 1 and 200005114 bits in length.

3.2 Inputs and Outputs

The inputs to the algorithm are given in table 1, the output in table 2:

Table 1: *f8* inputs

Parameter	Size (bits)	Comment
COUNT	32	Frame dependent input COUNT[0]...COUNT[31]
BEARER	5	Bearer identity BEARER[0]...BEARER[4]
DIRECTION	1	Direction of transmission DIRECTION[0]
CK	128	Confidentiality key CK[0]...CK[127]
LENGTH	X18 ¹	The number of bits to be encrypted/decrypted (1-200005114)
IBS	1-200005114	Input bit stream IBS[0]...IBS[LENGTH-1]

Table 2: *f8* output

Parameter	Size (bits)	Comment
OBS	1-200005114	Output bit stream OBS[0]...OBS[LENGTH-1]

3.3 Components and Architecture

(See fig 1 Annex A)

¹X18 is a parameter whose value is yet to be defined. In the sample C-code we treat LENGTH as a 32-bit integer.

The keystream generator is based on the block cipher **KASUMI** that is specified in [4]. **KASUMI** is used in a form of output-feedback mode and generates the output keystream in multiples of 64-bits.

The feedback data is modified by static data held in a 64-bit register **A**, and an (incrementing) 64-bit counter **BLKCNT**.

3.4 Initialisation

In this section we define how the keystream generator is initialised with the key variables before the generation of keystream bits.

We set the 64-bit register **A** to **COUNT || BEARER || DIRECTION || 0...0**

(left justified with the right most 26 bits set to 0).

i.e. **A = COUNT[0]...COUNT[31] BEARER[0]...BEARER[4] DIRECTION[0] 0...0**

We set counter **BLKCNT** to zero.

We set the key modifier **KM** to 0x55555555555555555555555555555555

We set **KSB₀** to zero.

One operation of **KASUMI** is then applied to the register **A**, using a modified version of the confidentiality key.

$$\mathbf{A} = \mathbf{KASUMI}[\mathbf{A}]_{\mathbf{CK} \oplus \mathbf{KM}}$$

3.5 Keystream Generation

Once the keystream generator has been initialised in the manner defined in section 3.4, it is ready to be used to generate keystream bits. The plaintext/ciphertext to be encrypted/decrypted consists of **LENGTH** bits (1-200005114) whilst the keystream generator produces keystream bits in multiples of 64 bits. Between 0 and 63 of the least significant bits are discarded from the last block depending on the total number of bits required by **LENGTH**.

So let **BLOCKS** be equal to (**LENGTH**/64) rounded up to the nearest integer. (For instance, if **LENGTH** = 128 then **BLOCKS** = 2; if **LENGTH** = 129 then **BLOCKS** = 3.)

To generate each keystream block (**KSB**) we perform the following operation:

For each integer **n** with $1 \leq n \leq \mathbf{BLOCKS}$ we define:

$$\mathbf{KSB}_n = \mathbf{KASUMI}[\mathbf{A} \oplus \mathbf{BLKCNT} \oplus \mathbf{KSB}_{n-1}]_{\mathbf{CK}}$$

where **BLKCNT** = n-1

The individual bits of the keystream are extracted from **KSB₁** to **KSB_{BLOCKS}** in turn, most significant bit first, by applying the operation:

For **n** = 1 to **BLOCKS**, and for each integer **i** with $0 \leq i \leq 63$ we define:

$$\mathbf{KS}[(n-1)*64+i] = \mathbf{KSB}_n[i]$$

3.6 Encryption/Decryption

Encryption/decryption operations are identical and are performed by the exclusive-OR of the input data (IBS) with the generated keystream (KS).

For each integer i with $0 \leq i \leq \text{LENGTH}-1$ we define:

$$\text{OBS}[i] = \text{IBS}[i] \oplus \text{KS}[i]$$

4 Integrity algorithm f_9

4.1 Introduction

The integrity algorithm f_9 computes a Message Authentication Code (MAC) on an input message under an integrity key **IK**. ~~The message may be between 1 and 5114 bits in length. There is no limitation on the input message length of the f_9 algorithm.~~

For ease of implementation the algorithm is based on the same block cipher (**KASUMI**) as is used by the confidentiality algorithm f_8 .

4.2 Inputs and Outputs

The inputs to the algorithm are given in table 3, the output in table 4:

Table 3: f_9 inputs

Parameter	Size (bits)	Comment
COUNT-I	32	Frame dependent input COUNT-I[0]...COUNT-I[31]
FRESH	32	Random number FRESH[0]...FRESH[31]
DIRECTION	1	Direction of transmission DIRECTION[0]
IK	128	Integrity key IK[0]...IK[127]
LENGTH	$\times 19^2$	The number of bits to be 'MAC'd
MESSAGE	LENGTH	Input bit stream

Table 4: f_9 output

Parameter	Size (bits)	Comment
MAC-I	32	Message authentication code MAC-I[0]...MAC-I[31]

4.3 Components and Architecture

(See fig 2 Annex A)

The integrity function is based on the block cipher **KASUMI** that is specified in [4]. **KASUMI** is used in a chained mode to generate a 64-bit digest of the message input. Finally the leftmost 32-bits of the digest are taken as the output value **MAC-I**.

² ~~$\times 19$ is a parameter whose value is yet to be defined.~~ In the sample C-code we treat LENGTH as a 32-bit integer.

4.4 Initialisation

In this section we define how the integrity function is initialised with the key variables before the calculation commences.

We set the working variables: $\mathbf{A} = \mathbf{0}$
and $\mathbf{B} = \mathbf{0}$

We set the key modifier \mathbf{KM} to 0xAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA

We concatenate \mathbf{COUNT} , \mathbf{FRESH} , $\mathbf{MESSAGE}$ and $\mathbf{DIRECTION}$. We then append a single '1' bit, followed by between 0 and 63 '0' bits so that the total length of the resulting string \mathbf{PS} (padded string) is an integral multiple of 64 bits, i.e.:

$$\mathbf{PS} = \mathbf{COUNT}[0] \dots \mathbf{COUNT}[31] \mathbf{FRESH}[0] \dots \mathbf{FRESH}[31] \mathbf{MESSAGE}[0] \dots \mathbf{MESSAGE}[\mathbf{LENGTH}-1] \mathbf{DIRECTION}[0] \mathbf{1} \mathbf{0}^*$$

Where $\mathbf{0}^*$ indicates between 0 and 63 '0' bits.

4.5 Calculation

We split the padded string \mathbf{PS} into 64-bit blocks \mathbf{PS}_i where:

$$\mathbf{PS} = \mathbf{PS}_0 \parallel \mathbf{PS}_1 \parallel \mathbf{PS}_2 \parallel \dots \parallel \mathbf{PS}_{\mathbf{BLOCKS}-1}$$

We perform the following operations for each integer n with $0 \leq n \leq \mathbf{BLOCKS}-1$:

$$\begin{aligned} \mathbf{A} &= \mathbf{KASUMI}[\mathbf{A} \oplus \mathbf{PS}_n]_{\mathbf{IK}} \\ \mathbf{B} &= \mathbf{B} \oplus \mathbf{A} \end{aligned}$$

Finally we perform one more application of \mathbf{KASUMI} using a modified form of the integrity key \mathbf{IK} .

$$\mathbf{B} = \mathbf{KASUMI}[\mathbf{B}]_{\mathbf{IK} \oplus \mathbf{KM}}$$

The 32-bit $\mathbf{MAC-I}$ comprises the left-most 32 bits of the result.

$$\mathbf{MAC-I} = \text{lefthalf}[\mathbf{B}]$$

i.e. For each integer i with $0 \leq i \leq 31$ we define:

$$\mathbf{MAC-I}[i] = \mathbf{B}[i].$$

Bits $\mathbf{B}[32] \dots \mathbf{B}[63]$ are discarded.