3GPP TSG SA WG 3 (Security) meeting #11 Mainz, 22—24 February, 2000

S3-000159

Document (Rev. of S3-000111)
e.g. for 3GPP use the format TP-99xxx
or for SMG, use the format P-99-xxx

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Form: CR cover sheet, version 2 for 3GPP and SMG The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc Proposed change affects: (at least one should be marked with an X) The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc U)SIM X ME X UTRAN / Radio Core Network								
Source:	S3					Date:	2000-02-23	
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6.3.3 Authentication and key agreement

The purpose of this procedure is to authenticate the user and establish a new pair of cipher and integrity keys between the VLR/SGSN and the <u>MSUSIM</u>. During the authentication, the <u>user-USIM</u> verifies the freshness of the authentication vector that is used.

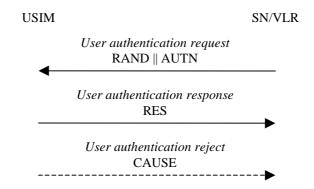


Figure 8: Authentication and key establishment

The VLR/SGSN invokes the procedure by selecting the next unused authentication vector from the ordered array of authentication vectors in the VLR database. The VLR/SGSN sends to the <u>user-USIM</u> the random challenge RAND and an authentication token for network authentication AUTN from the selected authentication vector.

Upon receipt the user proceeds as shown in Figure 9.

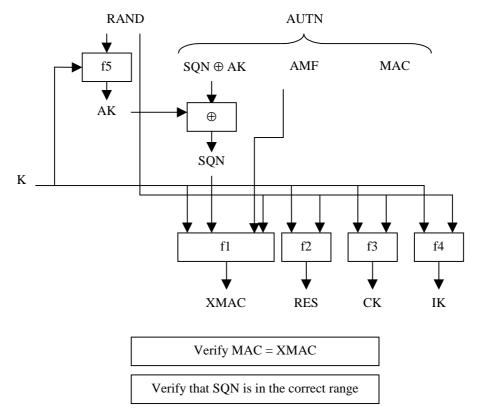


Figure 9: User authentication function in the USIM

Upon receipt of RAND and AUTN the <u>user_USIM</u> first computes the anonymity key $AK = f5_K$ (RAND) and retrieves the sequence number $SQN = (SQN \oplus AK) \oplus AK$.

Next the <u>user-USIM</u> computes XMAC = $f1_K$ (SQN || RAND || AMF) and compares this with MAC which is included in AUTN. If they are different, the user sends *user authentication reject* back to the VLR/SGSN with an indication of the cause and the user abandons the procedure.

Next the USIM verifies that the received sequence number SQN is in the correct range.

If the <u>user-USIM</u> considers the sequence number to be not in the correct range, <u>he-it</u> sends *synchronisation failure* back to the VLR/SGSN including an appropriate parameter, and abandons the procedure.

The synchronisation failure message contains the parameter AUTS. It is AUTS = $Conc(SQN_{MS}) \parallel MACS$. $Conc(SQN_{MS}) = SQN_{MS} \oplus f_{SK}(MACS)$ is the concealed value of the counter SEQ_{MS} in the MS, and MACS = $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel RAND \parallel AMF$) where RAND is the random value received in the current user authentication request. $f_{SK}(SEQ_{MS}) \parallel AMF$) where $f_{SK}(SEQ_{MS}) \parallel AMF$ is a message authentication code (MAC) function with the property that no valuable information can be inferred from the function values of $f_{SK}(SEQ_{MS}) \parallel AMF$.

The AMF used to calculate MACS assumes a dummy value of all zeros so that it does not need to be transmitted in the clear in the re-synch message.

The construction of the parameter AUTS in shown in the following Figure 10:

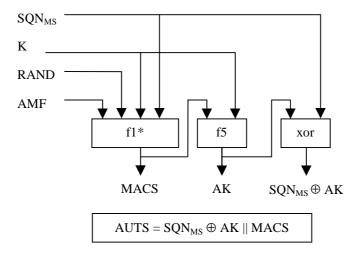


Figure 10: Construction of the parameter AUTS

If the sequence number is considered to be in the correct range however, the <u>user-USIM</u> computes RES = $f2_K$ (RAND) and includes this parameter in a *user authentication response* back to the VLR/SGSN. Finally the <u>user-USIM</u> computes the cipher key CK = $f3_K$ (RAND) and the integrity key IK = $f4_K$ (RAND). Note that if this is more efficient, RES, CK and IK could also be computed earlier at any time after receiving RAND. If the <u>USIM</u> also supports GSM AKA, it shall derive the GSM cipher key Kc from the <u>UMTS</u> cipher/integrity keys CK and IK using conversion function c3. <u>UMTS</u> keys are sent to the MS along with the derived GSM key for <u>UMTS-GSM</u> interoperability purposes. <u>USIM</u> shall store original CK, IK until the next successful execution of AKA. The <u>MS-USIM</u> also stores RAND until completion of current AKA, for re-synchronisation purposes.

Upon receipt of *user authentication response* the VLR/SGSN compares RES with the expected response XRES from the selected authentication vector. If XRES equals RES then the authentication of the user has passed. The VLR/SGSN also selects the appropriate cipher key CK and integrity key IK from the selected authentication vector.