

Agenda Item:

Source: SAMSUNG Electronics Co.

Title: Parameters (a,b) selection for rate matching algorithm

Document for: Discussion

1. Introduction

This document supplements Tdoc [1] and [2] proposed by Samsung and LGIC with some analysis and some additional simulation results over AWGN channel. In the Ad Hoc 5 report, it has been mentioned that (a,b) parameter is to be studied for optimization. In this document, we show that (a,b) parameter is only related to the initial offset of puncturing pattern. In order to make arbitrary offset for a given puncturing pattern, it is sufficient to control (a,b) parameter. We also show that a specific (a,b) parameter set determines a regular puncturing pattern.

2. Initial offset and period of rate matched symbol stream

Assume that N_{cs} is the number of bits before rate matching for each rate matching block and N_{is} is number of bits after rate matching for each rate matching block. For given rate matching parameters, the number punctured bits y is determined as $y = N_{cs} - N_{is}$ and puncturing is performed by the following rate matching algorithm.

```
e=(a*s(k)*y+bNcs) mod aNcs;  
if e=0 then e=aNcs;  
m=1  
do while m<=Ncs  
  e=e-a*y;  
  if e <=0 then  
    punctured  
    puncture bit m from set S0  
    e=e+a*Ncs  
  end if  
  m=m+1;  
end do
```

For simplicity, let's define the initial offset **Offset_m** of rate matching block to be the number of symbols from the first symbol to the first punctured symbol. Also, let's define period **P** to be the puncturing period. In figure 1, definition of **Offset_m** and period **P** are given. If N_{cs}/y is an integer, then initial offset and period **P** of the rate matching block can be easily obtained from the above algorithm and they are given by

$$\text{Offset}_m = \lceil (b/a) * (N_{cs}/y) \rceil = \lceil (b/a) * P \rceil \quad (1)$$

$$P = N_{cs}/y \quad (2)$$

where $1 \leq a \leq b * (N_{cs}/y)$ and $1 \leq b \leq a$.

The current rate matching algorithm uses $(a,b)=(2,1)$. This means that the initial offset **Offset_m** of the current algorithm is just 1/2 of the period $P(=N_{cs}/y)$. For a given parameter b , the initial offset decreases as the parameter a increases. If $a^3 N_{cs}/y$, the initial offset always becomes 1. Therefore, the first bit should be punctured.

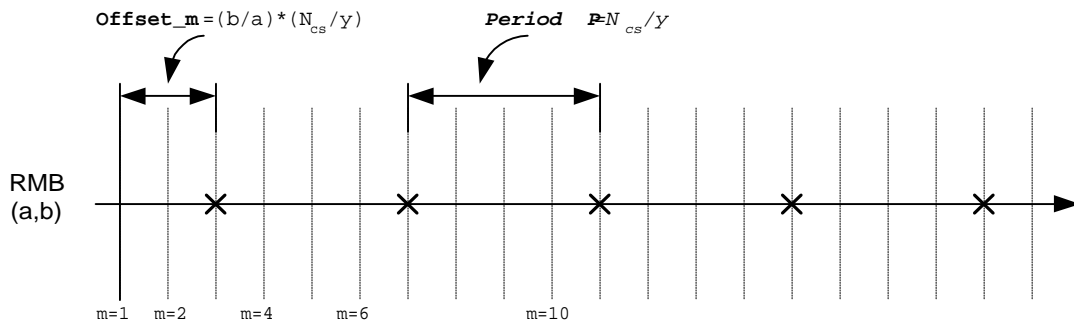


Figure 1. Definitions of initial offset **Offset_m** and period **P** in the rate matching block.

The following two examples show the role of (a,b) parameter for the initial offset control for each rate matching block. In the following cases, we assume that *b* is 1 according to the current rate matching scheme. In the following figures, 'X' means the punctured bit.

Example 1. (a,b)=(2, 1) for RMB2 and (a,b)=(1, 1) for RMB3. Period=4 (=N_{cs}/y).

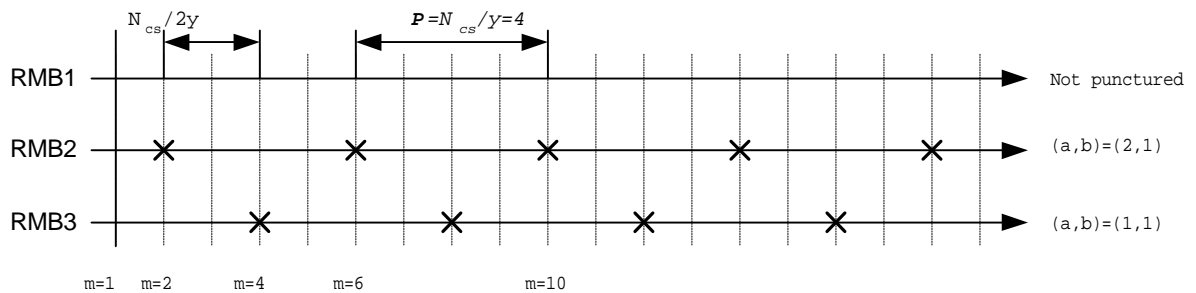


Figure 2. **Offset_m=2** for RMB2, **Offset_m=4** for RMB3. **P=4**.

Example 2. (a,b)=(2, 1) for RMB2 and (a,b)=(2, 1) for RMB3. Period=4 (=N_{cs}/y).

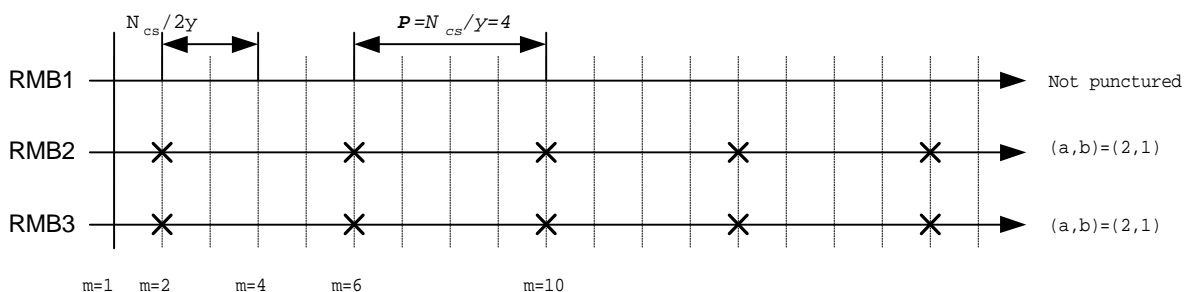
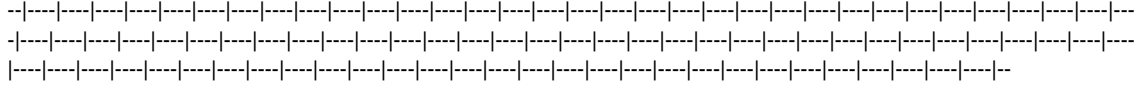


Figure 3. **Offset_m=2** for RMB2, **Offset_m=2** for RMB3. **P=4**.

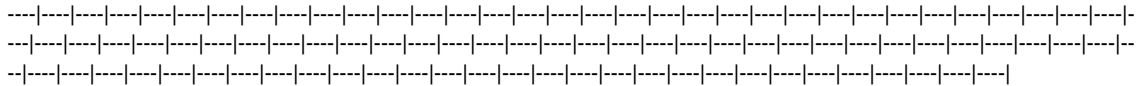
Example 3. $N_{cs}=480, N_{is}=384, \gamma=96: a=2, b=1,2.$

$(a,b)=(2,1)$



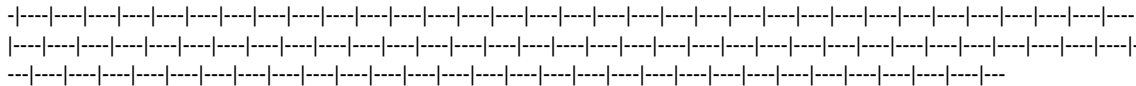
● Note that ‘-’ and ‘|’ mean the unpunctured bit and the punctured bits, respectively.

$(a,b)=(2,2)$

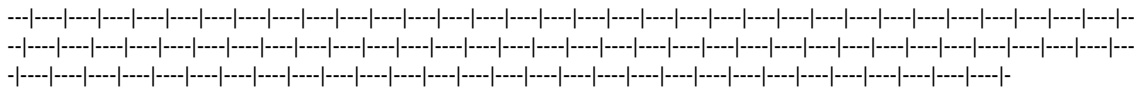


Example 4. $N_{cs}=480, N_{is}=384, \gamma=96: a=3, b=1,2,3.$

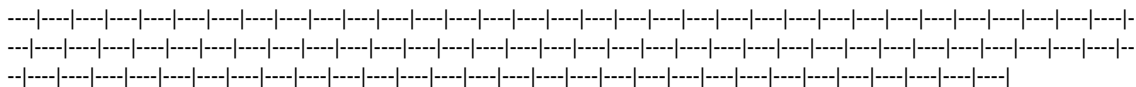
$(a,b)=(3,1)$



$(a,b)=(3,2)$



$(a,b)=(3,3)$



3. Simulation conditions

For turbo codes, the simulation conditions are as follows;

- Block sizes: 320, 640
- Puncturing rates: $p=20, 15, 10\%$
- Decoding algorithm: Log MAP decoder
- Turbo interleaver: PIL
- Number of iterations: 12
- Number of frame errors: greater than 100
- Channel model: AWGN
- Algorithms: SEC & LGIC puncturing

In the following figures, $_20$, $_15$, and $_10$ mean puncturing with $p=20, 15$, and 10 , respectively [1],[2]. Notation of (a,b,a,b) means the following parameters for each rate matching block (RMB).

(a,b,a,b)	RMB1	RMB2	RMB3
(2,1,2,1)	Not used	(a,b)=(2,1)	(a,b)=(2,1)
(2,1,1,1)	Not used	(a,b)=(2,1)	(a,b)=(1,1)
(2,1,3,1)	Not used	(a,b)=(2,1)	(a,b)=(3,1)

4. Results and conclusion

Figure 1 to 3 show the performances with interleaver size of 320, and Figure 4 to 6 show the performances of interleaver size of 640, both with different set of (a,b,a,b) parameter. According to the simulation results in Figure 1 to 6, the performances with (2,1,1,1) are slightly better consistently than those with others. Therefore, we propose to use (a,b,a,b) parameter set as (2,1,1,1) for RMB2 and RMB3.

5. References

- [1] "Unified rate matching scheme for Turbo/convolutional codes and up/down links", Samsung Electronics Co., TSGR1#6(99)919
- [2] "Unified rate matching scheme for turbo codes in both uplink and downlink", Samsung Electronics Co. and LGIC, TSGR1#6(99)a30

Contact inform:

Kimmingu@samsung.co.kr

Bjkim@telecom.samsung.co.kr

interleaver size=320, PIL, iteration=12, Frame Error count=100, Downlink, P=20%

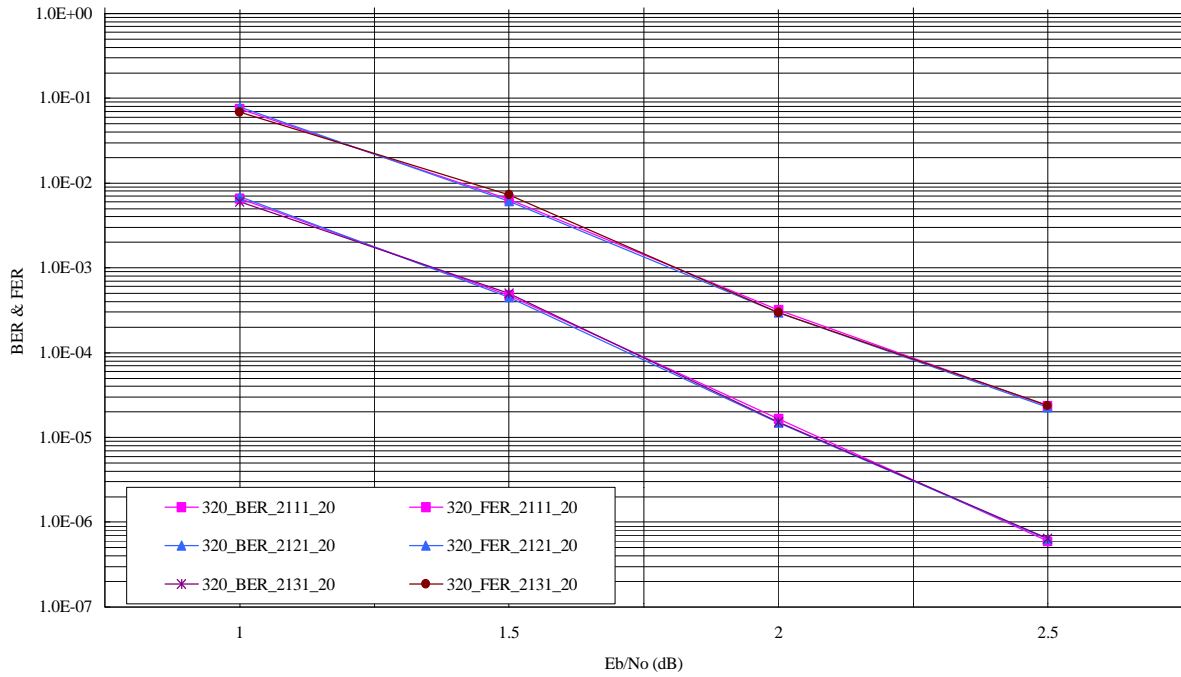


Figure 1. BER and FER of rate matching algorithm with various (a,b). N=320, p=20%.

interleaver size=320, PIL, iteration=12, Frame Error count=100, Downlink, P=15%

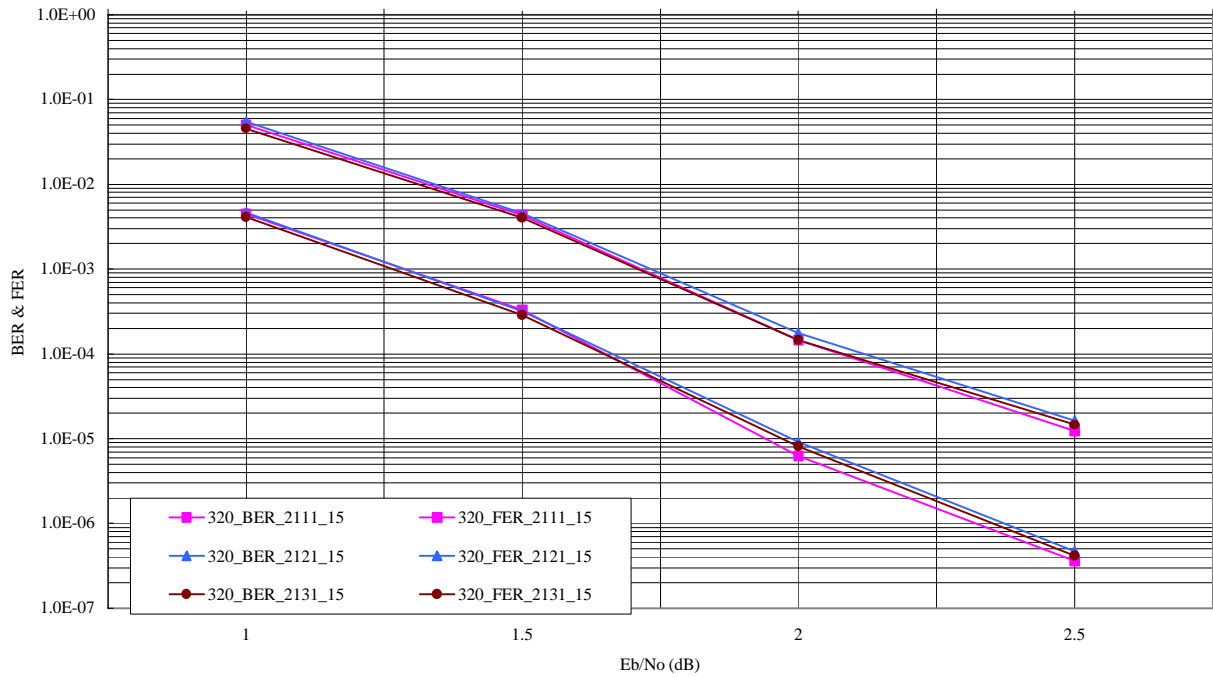


Figure 2. BER and FER of rate matching algorithm with various (a,b). N=320, p=15%.

interleaver size=320, PIL, iteration=12, Frame Error count=100, Downlink, P=10%

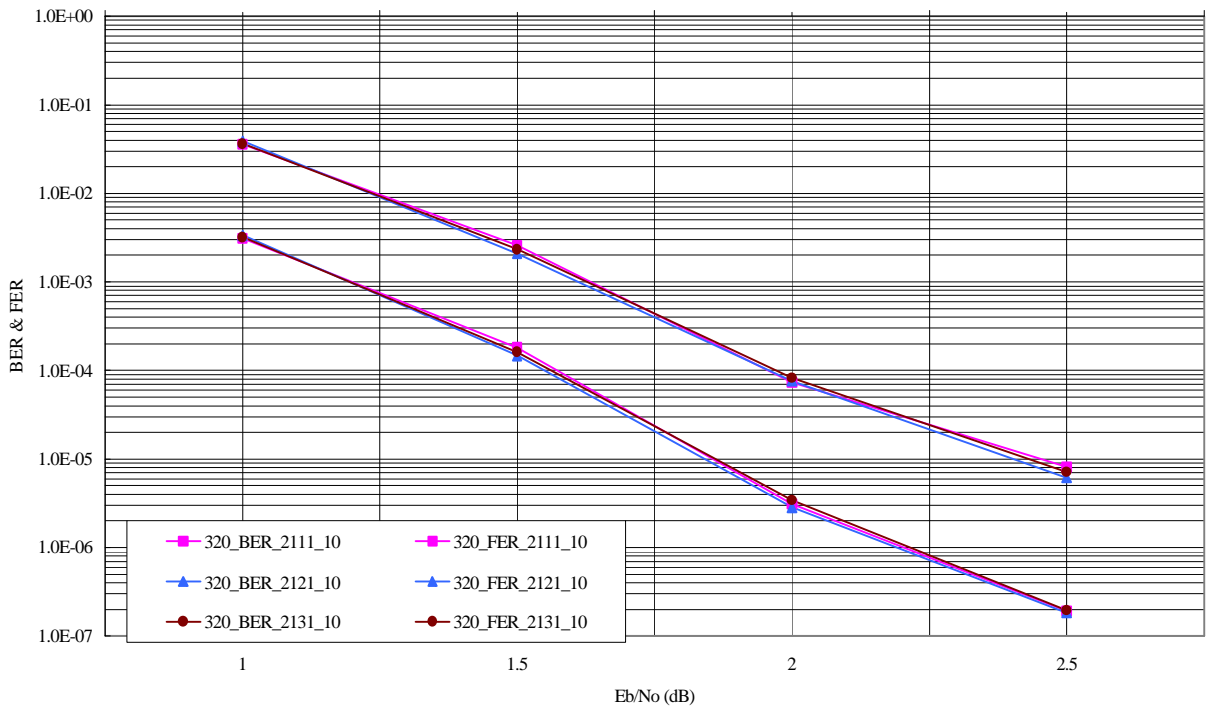


Figure 3. BER and FER of rate matching algorithm with various (a,b). N=320, p=20%.

interleaver size=640, PIL, iteration=12, Frame Error count=100, Downlink, P=20%

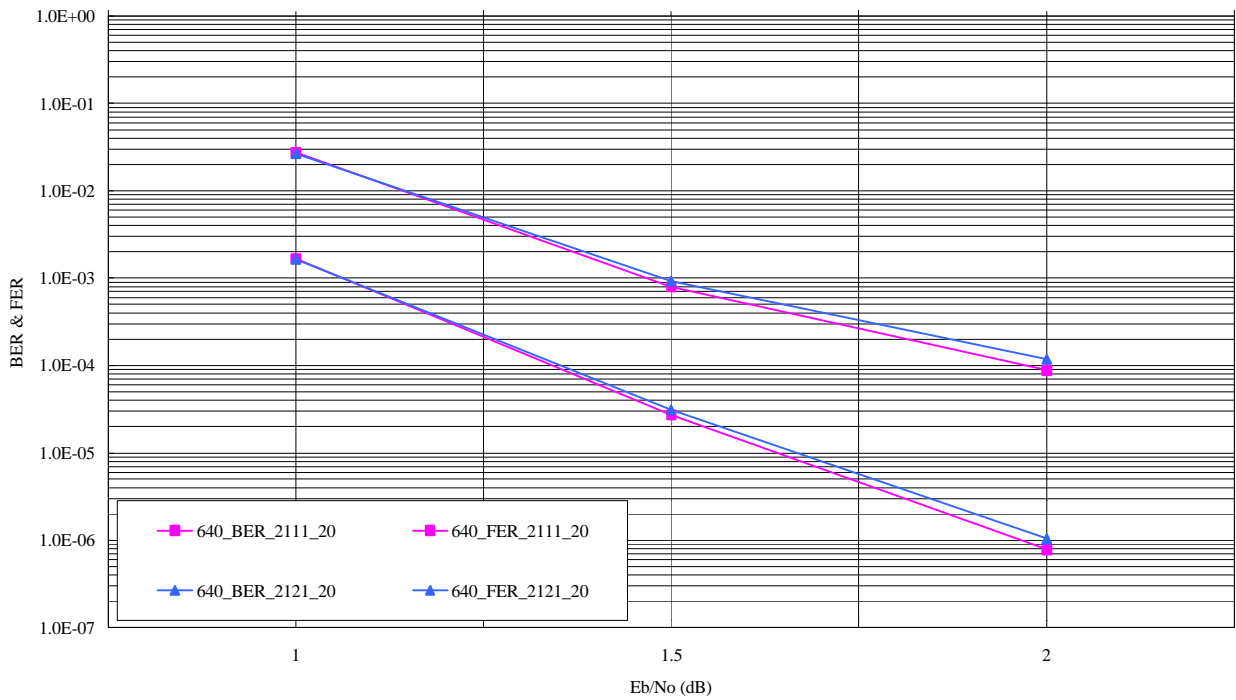


Figure 4. BER and FER of rate matching algorithm with various (a,b). N=640, p=20%.

interleaver size=640, PIL, iteration=12, Frame Error count=100, Downlink, P=15%

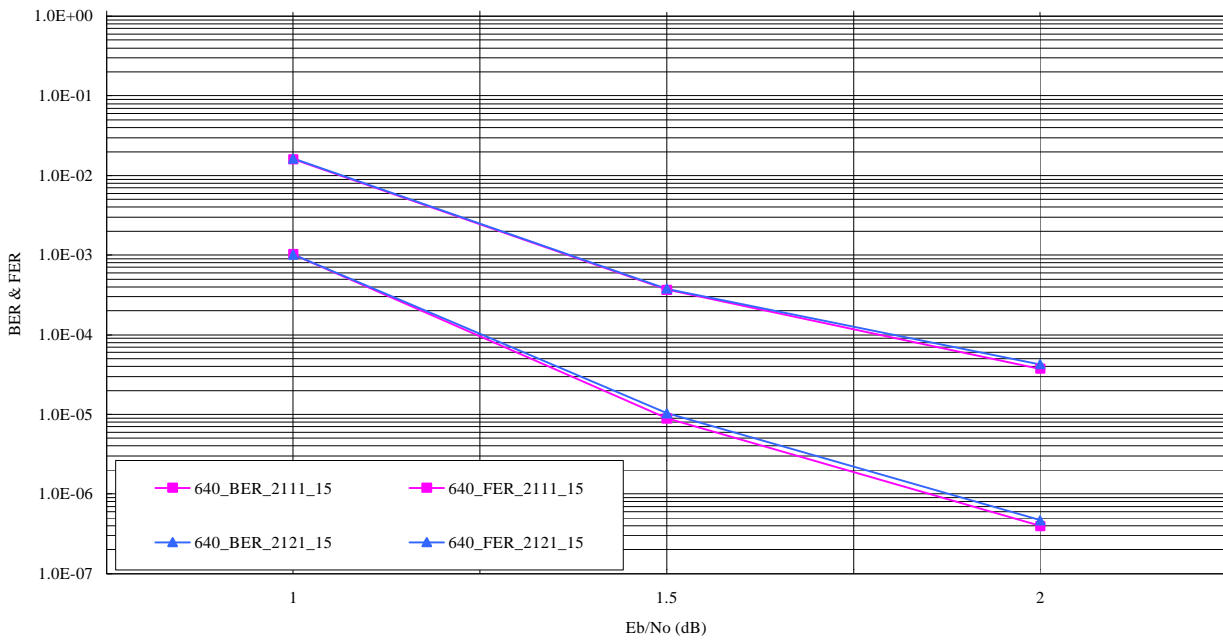


Figure 5. BER and FER of rate matching algorithm with various (a,b). N=640, p=15%.

interleaver size=640, PIL, iteration=12, Frame Error count=100, Downlink, P=10%

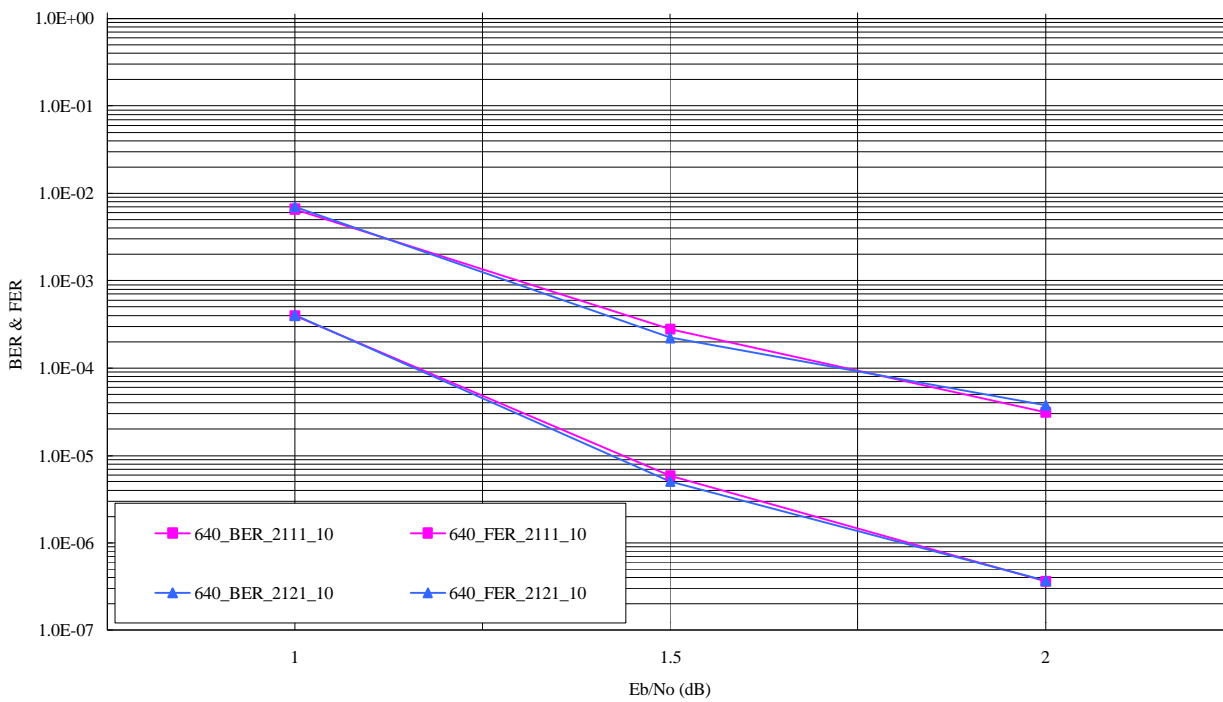


Figure 6. BER and FER of rate matching algorithm with various (a,b). N=640, p=10%.