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		25.222	CR	034	Curr	ent Version	on: 3.2.0	
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Form: CR cover sheet, version 2 for 3GPP and SMG  The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc  Proposed change affects: (at least one should be marked with an X)  The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc  WE UTRAN / Radio X Core Network								
Source:	Siemens A	3				Date:	07.04.2000	
Subject:	Alignment of	of Multiplexing for	TDD					
Work item:								
Category: A (only one category B shall be marked C with an X)	Addition of Functional	modification of fea		rlier releas		elease:	Phase 2 Release 96 Release 97 Release 98 Release 99 Release 00	X
Reason for change:	the FDD sp removed fro	fications have bee ec The case of om the CCTrCH is bid ambiguities in	TFCS re now de	configurat scribed. S	ion where a ome formula	transport	channel is	vith
Clauses affected	d: 4.2.4.	4.2.6, 4.2.7.1, 4.	.2.7.2. 4	l.2.12				
Other specs affected:	<u> </u>	e specifications ore ions ifications cifications	-	→ List of C	CRs: CRs: CRs:			
Other comments:								

<----- double-click here for help and instructions on how to create a CR.

## 4.2.4 Radio frame size equalisation

Radio frame size equalisation is padding the input bit sequence in order to ensure that the output can be segmented in  $F_i$  data segments of same size as described in the subclause 4.2.6.

The input bit sequence to the radio frame size equalisation is denoted by  $c_{i1}, c_{i2}, c_{i3}, \ldots, c_{iE_i}$ , where i is TrCH number and  $E_i$  the number of bits. The output bit sequence is denoted by  $t_{i1}, t_{i2}, t_{i3}, \ldots, t_{iT_i}$ , where  $T_i$  is the number of bits. The output bit sequence is derived as follows:

$$t_{ik} = c_{ik}$$
, for  $k = 1 \dots E_i$  and

$$t_{ik} = \{0 + 1\}$$
 for  $k = E_i + 1 \dots T_i$ , if  $E_i < T_i$ 

where

$$T_i = F_i * N_i$$
 and

$$N_i = [(E_i - 1)/F_i] + 1$$
  $N_i = [E_i/F_i]$  is the number of bits per segment after size equalisation.

### 4.2.6 Radio frame segmentation

When the transmission time interval is longer than 10 ms, the input bit sequence is segmented and mapped onto consecutive  $\underline{F_i}$  radio frames. Following radio frame size equalisation the input bit sequence length is guaranteed to be an integer multiple of  $F_i$ .

The input bit sequence is denoted by  $x_{i1}, x_{i2}, x_{i3}, \dots, x_{iX_i}$  where i is the TrCH number and  $X_i$  is the number bits. The  $F_{i_i}$  output bit sequences per TTI are denoted by  $y_{i,n_i1}, y_{i,n_i2}, y_{i,n_i3}, \dots, y_{i,n_iY_i}$  where  $n_i$  is the radio frame number in current TTI and  $Y_i$  is the number of bits per radio frame for TrCH i. The output sequences are defined as follows:

$$y_{i,n,k} = x_{i,((n_i-1)Y_i)+k}, n_i = 1...F_i, k = 1...Y_i$$

where

 $Y_i = (X_i / F_i)$  is the number of bits per segment,

 $x_{ik}$  is the  $k^{th}$ -bit of the input bit sequence and

-y<sub>i,n,k</sub> is the k<sup>th</sup>-bit of the output bit sequence corresponding to the n<sup>th</sup>-radio frame

The  $n_i$  –th segment is mapped to the  $n_i$  –th radio frame of the transmission time interval.

The input bit sequence to the radio frame segmentation is denoted by  $d_{i1}, d_{i2}, d_{i3}, \dots, d_{iT_i}$ , where i is the TrCH number and  $T_i$  the number of bits. Hence,  $x_{ik} = d_{ik}$  and  $X_i = T_i$ .

The output bit sequence corresponding <u>to</u> radio frame  $n_i$  is denoted by  $e_{i1}, e_{i2}, e_{i3}, \dots, e_{iN_i}$ , where i is the TrCH number and  $N_i$  is the number of bits. Hence,  $e_{i,k} = y_{i,n,k}$  and  $N_i = Y_i$ .

#### 4.2.7.1 Determination of rate matching parameters

The following relations, defined for all TFC j, are used when calculating the rate matching pattern:

$$Z_{0,i} = 0$$

$$\frac{Z_{ij} = \left[ \frac{\sum_{m=1}^{i} RM_m \cdot N_{mj}}{\sum_{m=1}^{I} RM_m \cdot N_{mj}} \cdot N_{data,j} \right]}{\sum_{m=1}^{I} RM_m \cdot N_{mj}} Z_{ij} = \left[ \frac{\left\{ \left( \sum_{m=1}^{i} RM_m \cdot N_{mj} \right) \cdot N_{data,j} \right\}}{\sum_{m=1}^{I} RM_m \cdot N_{mj}} \right] \text{ for all } i = 1 \dots I$$

$$\Delta N_{ij} = Z_{ij} - Z_{i-1,j} - N_{ij} \quad \text{ for all i = 1 .. I}$$

Puncturing can be used to minimise the required transmission capacity. The maximum amount of puncturing that can be applied is signalled from higher layers and denoted by PL. The possible values for  $N_{data}$  depend on the number of physical channels  $P_{max}$ , allocated to the respective CCTrCH, and on their characteristics (spreading factor, length of midamble and TFCI, usage of TPC and multiframe structure), which is given in [7].

Denote the number of data bits in each physical channel by  $U_{p,Sp}$ , where p refers to the sequence number  $I \pounds p \pounds P_{max}$  of this physical channel in the allocation message, and the second index Sp indicates the spreading factor with the possible values  $\{16, 8, 4, 2, 1\}$ , respectively. For each physical channel an individual minimum spreading factor  $Sp_{min}$  is transmitted by means of the higher layer. Then, for  $N_{data}$  one of the following values in ascending order can be chosen:

$$\left\{\!U_{1,16}, \ldots, \!U_{1,S1_{\min}}, \!U_{1,S1_{\min}} + \!U_{2,16}, \ldots, \!U_{1,S1_{\min}} + \!U_{2,S2_{\min}}, \ldots, \!U_{1,S1_{\min}} + \!U_{2,S2_{\min}} + \ldots + \!U_{P_{\max},16}, \ldots, \!U_{1,S1_{\min}} + \!U_{2,S2_{\min}} + \ldots + \!U_{P_{\max},(SP_{\max})_{\min}} \right\}$$

 $N_{data, j}$  for the transport format combination j is determined by executing the following algorithm:

SET1 = { 
$$N_{data}$$
 such that  $N_{data} - PL \cdot \sum_{x=1}^{I} \frac{RM_x}{\min_{1 \le y \le I} \{RM_y\}} \cdot N_{x,j} = \min_{1 \le y \le I} \{RM_y\} \cdot N_{data} - PL \cdot \sum_{x=1}^{I} RM_x \cdot N_{x,j}$  is non negative }

$$N_{data, j} = min SET1$$

The number of bits to be repeated or punctured,  $\Delta N_{ij}$ , within one radio frame for each TrCH i is calculated with the relations given at the beginning of this subclause for all possible transport format combinations j and selected every radio frame.

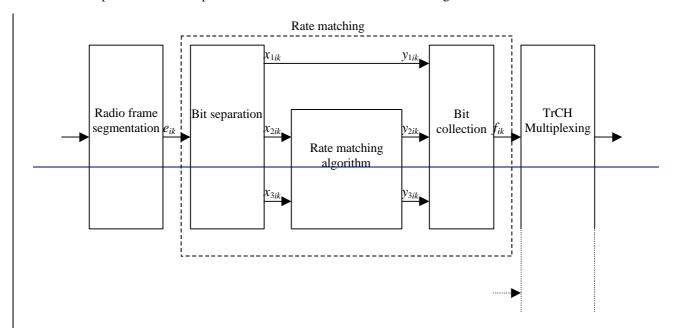
If  $\Delta N_{ij} = 0$  then the output data of the rate matching is the same as the input data and the rate matching algorithm of subclause 4.2.7.3 does not need to be executed.

Otherwise, the rate matching pattern is calculated with the algorithm described in subclause 4.2.7.3. For this algorithm the parameters  $e_{ini}$ ,  $e_{plus}$ ,  $e_{minus}$ , and  $X_i$  are needed, which are calculated according to the equations in subclauses 4.2.7.1.1 and 4.2.7.1.2.

## 4.2.7.2 Bit separation and collection for rate matching

The systematic bits (excluding bits for trellis termination) of turbo encoded TrCHs shall not be punctured. The systematic bit, first parity bit, and second parity bit in the bit sequence input to the rate matching block are therefore separated from each other. Puncturing is only applied to the parity bits and systematic bits used for trellis termination.

The bit separation function is transparent for uncoded TrCHs, convolutionally encoded TrCHs, and for turbo encoded TrCHs with repetition. The bit separation and bit collection are illustrated in figures 4 and 5.



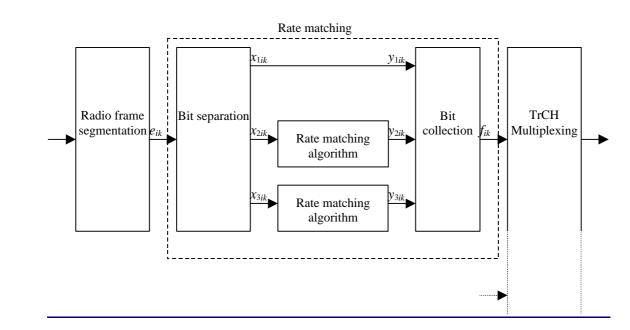


Figure 4: Puncturing of turbo encoded TrCHs

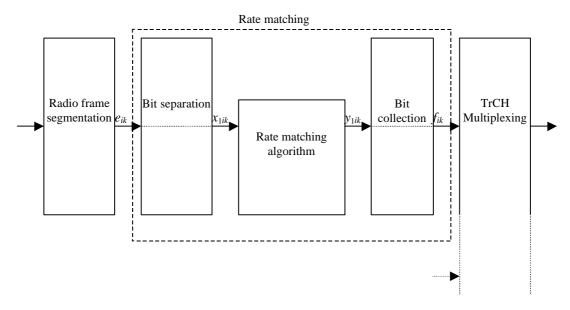


Figure 5: Rate matching for uncoded TrCHs, convolutionally encoded TrCHs, and for turbo encoded TrCHs with repetition

The bit separation is dependent on the 1<sup>st</sup> interleaving and offsets are used to define the separation for different TTIs. The offsets  $a_b$  for the systematic (b=1) and parity bits  $(b \in \{2,3\})$  are listed in table 4.

Table 4: TTI dependent offset needed for bit separation

TTI (ms)	<b>a</b> <sub>1</sub>	$a_2$	<b>a</b> <sub>3</sub>
10, 40	0	1	2
20, 80	0	2	1

The bit separation is different for different radio frames in the TTI. A second offset is therefore needed. The radio frame number for TrCH i is denoted by  $n_i$ , and the offset by  $\boldsymbol{b}_{n_i}$ .

Table 5: Radio frame dependent offset needed for bit separation

TTI (ms)	<b>b</b> <sub>0</sub>	<b>b</b> <sub>1</sub>	<b>b</b> <sub>2</sub>	<b>b</b> <sub>3</sub>	<b>b</b> <sub>4</sub>	<b>b</b> <sub>5</sub>	<b>b</b> <sub>6</sub>	<b>b</b> <sub>7</sub>
10	0	NA						
20	0	1	NA	NA	NA	NA	NA	NA
40	0	1	2	0	NA	NA	NA	NA
80	0	1	2	0	1	2	0	1

# 4.2.12 Multiplexing of different transport channels onto one CCTrCH, and mapping of one CCTrCH onto physical channels

Different transport channels can be encoded and multiplexed together into one Coded Composite Transport Channel (CCTrCH). The following rules shall apply to the different transport channels which are part of the same CCTrCH:

1) Transport channels multiplexed into one CCTrCh shall have co-ordinated timings. When the TFCS of a CCTrCH is changed because a one or more transport channels i are added to the CCTrCH or reconfigured within the CCTrCH, or removed from the CCTrCH, the change may only be made at the TTI of transport channel i may only start in of a radio frames with CFN fulfilling the relation

```
CFN_i \mod F_{max} = 0,
```

where  $F_{max}$  denotes the maximum number of radio frames within the transmission time intervals of all transport channels which are multiplexed into the same CCTrCH, including <u>any</u> transport channels *i* which <u>isare</u> added, or reconfigured <u>or have been removed</u>, and CFN<sub>i</sub> denotes the connection frame number of the first radio frame <u>of</u> the changed CCTrCH, within the transmission time interval of transport channel *i*.

After addition or reconfiguration of a transport channel *i* within a CCTrCH, the TTI of transport channel *i* may only start in radio frames with CFN fulfilling the relation

 $CFN_i \mod F_i = 0.$ 

- 2) Different CCTrCHs cannot be mapped onto the same physical channel.
- 3) One CCTrCH shall be mapped onto one or several physical channels.
- 4) Dedicated Transport channels and common transport channels cannot be multiplexed into the same CCTrCH.
- 5) For the common transport channels, only the FACH and PCH may belong to the same CCTrCH.
- 6) Each CCTrCH carrying a BCH shall carry only one BCH and shall not carry any other Transport Channel.
- 7) Each CCTrCH carrying a RACH shall carry only one RACH and shall not carry any other Transport Channel.

Hence, there are two types of CCTrCH.

CCTrCH of dedicated type, corresponding to the result of coding and multiplexing of one or several DCH.

CCTrCH of common type, corresponding to the result of the coding and multiplexing of a common channel, i.e. RACH and USCH in the uplink and DSCH, BCH, FACH or PCH in the downlink, respectively.

Transmission of TFCI is possible for CCTrCH containing Transport Channels of:

- dedicated type;
- USCH type;
- DSCH type;
- FACH and/or PCH type.