#### Document R1-00-512 3GPP TSG RAN WG1 Meeting #12 e.g. for 3GPP use the format TP-99xxx or for SMG, use the format P-99-xxx Korea, Seoul, 10-13 April 2000 Please see embedded help file at the bottom of this CHANGE REQUEST page for instructions on how to fill in this form correctly. Current Version: 3.2.0 25.223 CR 008 GSM (AA.BB) or 3G (AA.BBB) specification number ↑ ↑ CR number as allocated by MCC support team For submission to: RAN #8 for approval strategic (for SMG list expected approval meeting # here use only) for information non-strategic Form: CR cover sheet, version 2 for 3GPP and SMG The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc Proposed change affects: (U)SIM ME UTRAN / Radio X Core Network (at least one should be marked with an X) **Date:** 07.04.2000 Source: Siemens AG Editorial Modifications for 25.223 Subject: Work item: Category: Correction Release: Phase 2 Corresponds to a correction in an earlier release Release 96 (only one category Addition of feature Release 97 В shall be marked Functional modification of feature Release 98 C with an X) Editorial modification Release 99 X Release 00 Some editorial modifications to 25.223 are required after approval of removal of synch-Reason for case 3 and change of signal point constellation. The scope of the spec was somehow change: changed in a former release of the document. Clauses affected: 1, 7.1, 7.2 Other specs Other 3G core specifications → List of CRs: affected: Other GSM core → List of CRs:

Other comments:



specifications
MS test specifications

BSS test specifications

**O&M** specifications

<----- double-click here for help and instructions on how to create a CR.

→ List of CRs:

→ List of CRs:

→ List of CRs:

# 1 Scope

The present document describes spreading and modulation for UTRA Physical Layer TDD mode.

The present document describes multiplexing, channel coding and interleaving for UTRA Physical Layer TDD mode.

## 7 Synchronisation codes

## 7.1 Code Generation

The Primary code sequence,  $C_p$  is constructed as a so-called generalised hierarchical Golay sequence. The Primary SCH is furthermore chosen to have good aperiodic auto correlation properties.

Define 
$$a = \langle x_1, x_2, x_3, ..., x_{16} \rangle = \langle 1, 1, 1, 1, 1, 1, -1, -1, 1, -1, 1, -1, 1, -1, 1, 1 \rangle$$

The PSC code word is generated by repeating the sequence 'a' modulated by a Golay complementary sequence and creating a complex-valued sequence with identical real and imaginary components.

The PSC code word  $C_p$  is defined as  $C_p = \langle y(0), y(1), y(2), ..., y(255) \rangle$ 

and the left most index corresponds to the chip transmitted first in each time slot.

The  $\frac{16}{12}$  secondary synchronization code words,  $\{C_0, ..., C_{15}, C_{11}\}$  are complex valued with identical real and imaginary components, and are constructed from the position wise multiplication of a Hadamard sequence and a sequence z, defined as

$$b = \langle x_1, ..., x_8, -x_9, ..., -x_{16} \rangle = \langle 1, 1, 1, 1, 1, 1, -1, -1, -1, 1, -1, 1, -1, 1, -1 \rangle.$$

The Hadamard sequences are obtained as the rows in a matrix  $H_8$  constructed recursively by:

$$H_{0} = (1)$$

$$H_{k} = \begin{pmatrix} H_{k-1} & H_{k-1} \\ H_{k-1} & -H_{k-1} \end{pmatrix} \quad k \ge 1$$

The rows are numbered from the top starting with row  $\theta$  (the all zeros sequence).

Denote the *n*:th Hadamard sequence  $\underline{h}_n$  as a row of  $H_\delta$  numbered from the top, n = 0, 1, 2, ..., 255, in the sequel.

Furthermore, let  $h_m(i)$  and z(i) denote the i:th symbol of the sequence  $h_m$  and z, respectively where i = 0, 1, 2, ..., 255 and i = 0 corresponds to the leftmost symbol.

The i:th SCH code word,  $C_{SCH,i}$ ,  $i=0,\,...,\,\frac{15}{11}$  is then defined as

$$C_{\text{SCH,i}} = (1+j) \times \langle h_m(0) \times z(0), h_m(1) \times z(1), h_m(2) \times z(2), \dots, h_m(255) \times z(255) \rangle$$

where  $m = (16 \times i)$  and the leftmost chip in the sequence corresponds to the chip transmitted first in time.

This code word is chosen from every  $16^{th}$  row of the matrix  $H_{sc}$ , which yields 16 possible codewords.

The Secondary SCH code words are defined in terms of  $C_{SCH,i}$  and the definition of  $\{C_0,...,C_{15},C_{11}\}$  now follows as:

$$C_i = C_{SCH, i}, i=0,...,115$$

## 7.2 Code Allocation

Three SCH codes are QPSK modulated and transmitted in parallel with the primary synchronization code. The QPSK modulation carries the following information.

- The code group that the base station belongs to (5 bits; Cases 1, 2)
- The position of the frame within an interleaving period of 20 msec (1 bit, Cases 1, 2)

- The position of the slot within the frame (1 bit, Case 2)

The modulated codes are also constructed such that their cyclic-shifts are unique, i.e. a non-zero cyclic shift less than 2 (Case 1) and 4 (Case 2) of any of the sequences is not equivalent to some cyclic shift of any other of the sequences. Also, a non-zero cyclic shift less than 2 (Case 1) and 4 (Case 2) of any of the sequences is not equivalent to itself with any other cyclic shift less than 8. The secondary synchronization codes are partitioned into two code sets for Case 1 and four code sets for Case 2. The set is used to provide the following information:

#### Case 1:

Table 2: Code Set Allocation for Case 1

Code Set	Code Group			
1	0-15			
2	16-31			

The code group and frame position information is provided by modulating the secondary codes in the code set.

#### Case 2:

Table 3: Code Set Allocation for Case 2

Code Set	Code Group			
1	0-7			
2	8-15			
3	16-23			
4	24-31			

The slot timing and frame position information is provided by the comma free property of the code word and the Code group is provided by modulating some of the secondary codes in the code set.

The following SCH codes are allocated for each code set:

#### Case 1

Code set 1:  $C_0$ ,  $C_1$ ,  $C_2$ 

Code set 2:  $C_3$ ,  $C_4$ ,  $C_5$ .

#### Case 2

Code set 1:  $C_0$ ,  $C_1$ ,  $C_2$ 

Code set 2: C<sub>3</sub>, C<sub>4</sub>, C<sub>5</sub>.

Code set 3:  $C_6$ ,  $C_7$ ,  $C_8$ .

Code set 4: C<sub>9</sub>, C<sub>10</sub>, C<sub>11</sub>.

The following subsections 7.2.1 to 7.2.2 refer to the two cases of SCH/P-CCPCH usage as described in [7].

Note that in the Tables 4-6-5 corresponding to Cases 1,2, and 3, respectively, Frame 1 implies the frame with an odd SFN and Frame 2 implies the frame with an even SFN.

## 7.2.1 Code allocation for Case 1:

NOTE: Modulation by "j" indicates that the code is transmitted on the Q channel.

**Table 4: Code Allocation for Case 1** 

Code Group	Code Set	Frame 1			Frame 2			Associated toffset
0	1	$C_0$	C <sub>1</sub>	$C_2$	$C_0$	C <sub>1</sub>	-C <sub>2</sub>	t <sub>0</sub>
1	1	$C_0$	-C <sub>1</sub>	C <sub>2</sub>	C <sub>0</sub>	-C <sub>1</sub>	-C <sub>2</sub>	t <sub>1</sub>
2	1	-C <sub>0</sub>	C <sub>1</sub>	C <sub>2</sub>	-C <sub>0</sub>	C <sub>1</sub>	-C <sub>2</sub>	t <sub>2</sub>
3	1	-C <sub>0</sub>	-C <sub>1</sub>	C <sub>2</sub>	-C <sub>0</sub>	-C <sub>1</sub>	-C <sub>2</sub>	t <sub>3</sub>
4	1	jC <sub>0</sub>	<del>JC</del> ₁ <u>jC</u> ₁	C <sub>2</sub>	jC <sub>0</sub>	jC₁	-C <sub>2</sub>	t <sub>4</sub>
5	1	jC <sub>0</sub>	-jC₁	C <sub>2</sub>	jC <sub>0</sub>	-jC₁	-C <sub>2</sub>	t <sub>5</sub>
6	1	-jC <sub>0</sub>	<del>JC</del> ₁ <u>jC</u> ₁	C <sub>2</sub>	-jC₀	jC₁	-C <sub>2</sub>	t <sub>6</sub>
7	1	-jC <sub>0</sub>	-jC₁	$C_2$	-jC <sub>0</sub>	-jC₁	-C <sub>2</sub>	t <sub>7</sub>
8	1	jC <sub>0</sub>	JC <sub>2</sub> jC <sub>2</sub>	C <sub>1</sub>	jC <sub>0</sub>	jC <sub>2</sub>	-C <sub>1</sub>	t <sub>8</sub>
9	1	jC <sub>0</sub>	-jC <sub>2</sub>	C <sub>1</sub>	jC <sub>0</sub>	-jC <sub>2</sub>	-C <sub>1</sub>	t <sub>9</sub>
10	1	-jC <sub>0</sub>	JC <sub>2</sub> jC <sub>2</sub>	C <sub>1</sub>	-jC <sub>0</sub>	jC <sub>2</sub>	-C <sub>1</sub>	t <sub>10</sub>
11	1	-jC <sub>0</sub>	-jC <sub>2</sub>	C <sub>1</sub>	-jC₀	-jC <sub>2</sub>	-C <sub>1</sub>	t <sub>11</sub>
12	1	jC₁	<del>JC</del> 2jC2	C <sub>0</sub>	<del>JC</del> ₁ <u>jC</u> ₁	jC <sub>2</sub>	-C <sub>0</sub>	t <sub>12</sub>
13	1	jC₁	-jC <sub>2</sub>	C <sub>0</sub>	<del>JC</del> ₁ <u>iC</u> ₁	-jC <sub>2</sub>	-C <sub>0</sub>	t <sub>13</sub>
14	1	-jC₁	<del>JC</del> 2 <u>jC</u> 2	C <sub>0</sub>	-jC₁	jC <sub>2</sub>	-C <sub>0</sub>	t <sub>14</sub>
15	1	-jC₁	-jC <sub>2</sub>	$C_0$	-jC₁	-jC <sub>2</sub>	-C <sub>0</sub>	t <sub>15</sub>
16	2	C <sub>3</sub>	C <sub>4</sub>	C <sub>5</sub>	$C_3$	C <sub>4</sub>	-C <sub>5</sub>	t <sub>16</sub>
17	2	C <sub>3</sub>	-C <sub>4</sub>	C <sub>5</sub>	C <sub>3</sub>	-C <sub>4</sub>	-C <sub>5</sub>	t <sub>17</sub>
				•••				
20	2	jC₃	<del>JC</del> 4 <u>jC</u> 4	C <sub>5</sub>	jC₃	jC <sub>4</sub>	-C <sub>5</sub>	t <sub>20</sub>
•••				•••				
24	2	jC <sub>3</sub>	<del>JC</del> ₅ <u>jC</u> ₅	C <sub>4</sub>	jC <sub>3</sub>	<del>JC</del> ₅ <u>jC</u> ₅	-C <sub>4</sub>	t <sub>24</sub>
	•••							
31	2	-jC₄	-jC₅	С3	-jC₄	-jC₅	-C₃	t <sub>31</sub>

NOTE: The code construction for code groups 0 to 15 using only the SCH codes from code set 1 is shown. The construction for code groups 16 to 31 using the SCH codes from code set 2 is done in the same way.