
[*secondary scrambling codes, in order to limit the signalling requirements.*](#) →

<Editor's note: it is not standardised how many scrambling codes a UE must decode in parallel.>

The scrambling code sequences are constructed by combining two real sequences into a complex sequence. Each of the two real sequences are constructed as the position wise modulo 2 sum of [40960 chip segments of] two binary m -sequences generated by means of two generator polynomials of degree 18. The resulting sequences thus constitute segments of a set of Gold sequences. The scrambling codes are repeated for every 10 ms radio frame. Let x and y be the two sequences respectively. The x sequence is constructed using the primitive (over GF(2)) polynomial $1+X^7+X^{18}$. The y sequence is constructed using the polynomial $1+X^6+X^7+$

$$X^{10} + X^{18}$$

<Editor's note: [] is due to the fact that only 4.096 Mcps is a working assumptions. 1.024, 8.192, and 16.384 Mcps are ffs.>

Let $n_{17} \dots n_0$ be the binary representation of the scrambling code number n (decimal) with n_0 being the least significant bit. The x sequence depends on the chosen scrambling code number n and is denoted x_n in the sequel. Furthermore, let $x_n(i)$ and $y(i)$ denote the i .th symbol of the sequence x_n and y , respectively

The m -sequences x_n and y are constructed as:

Initial conditions:

$$x_n(0)=n_0, x_n(1)=n_1, \dots, x_n(16)=n_{16}, x_n(17)=n_{17}$$

$$y(0)=y(1)=\dots=y(16)=y(17)=1$$

Recursive definition of subsequent symbols:

$$x_n(i+18) = x_n(i+7) + x_n(i) \text{ modulo } 2, i=0, \dots, 2^{18}-20,$$

$$y(i+18) = y(i+10)+y(i+7)+y(i+5)+y(i) \text{ modulo } 2, i=0, \dots, 2^{18}-20.$$

The n :th Gold code sequence z_n is then defined as

$$z_n(i) = x_n(i) + y(i) \text{ modulo } 2, i=0, \dots, 2^{18}-2.$$

x_0 is constructed with $x_0(0) = x_0(1) = \dots x_0(16) = 0$ $x_0(17) = 1$ as initial conditions.

x_n is constructed with n phase shift from x_0 .

These binary code words are converted to real valued sequences by the transformation '0' ->

Finally, the n :th complex scrambling code sequence C_{scramb} is defined as (the lowest index corresponding to the chip scrambled first in each radio frame): (see Table 1 for definition of N and M)

$$C_{scramb}(i) = z'_n(i) + j z'_n(i+M), i=0, 1, \dots, N-1.$$

<Editor's note: the values 40960 is based on an assumption of a chip rate of 4.096 Mcps.>

Note that the pattern from phase 0 up to the phase of 10 msec is repeated.

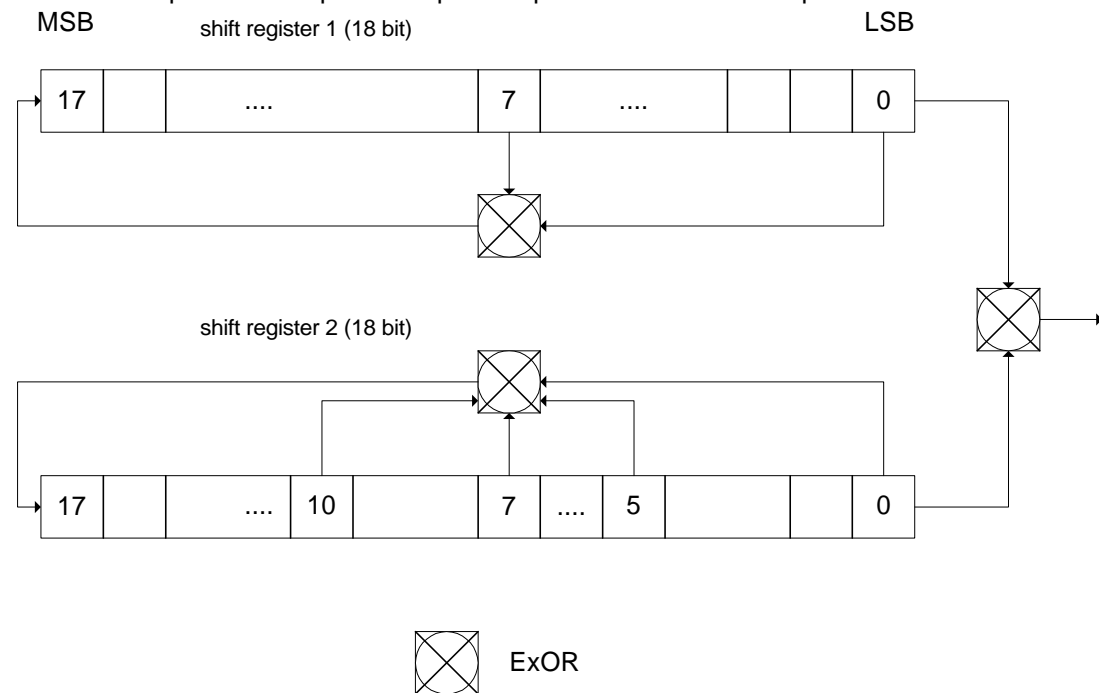


Figure 1. Configuration of downlink scrambling code generator

<Editor's note: a replacement figure for the above is to be prepared showing both I & Q generation.>

chip rate (Mcps)	Period N	IQ Offset M	Range of phase (chip)	
			for in-phase component	for quadrature component
[1.024]	[10240]	[131072]	0 – N-1	M – N+M-1
4.096	40960	131072		
[8.192]	[81920]	[131072]		
[16.384]	[163840]	[131072]		

Table 1. Correspondence between chip rate and downlink scrambling code phase range

Reference

- [1] 3GPP TSGR1#6 (99)924, 'Multiple scrambling code', Source: Samsung
- [2] 3GPP TSGR1#7 (99)a86, 'TS 25.213 V2.0.1 (1999-08) Spreading and modulation(FDD)'