

**3GPP TSG RAN WG1 Meeting #9**  
**Dresden, Germany, 30 NOV 1999 - 03 DEC 1999**

**Document R1-99I37**

e.g. for 3GPP use the format TP-99xxx  
or for SMG, use the format P-99-xxx

<h2 style="margin: 0;">CHANGE REQUEST</h2>		<i>Please see embedded help file at the bottom of this page for instructions on how to fill in this form correctly.</i>
<b>25.222</b>	<b>CR 009r1</b>	Current Version: <b>V3.0.0</b>
GSM (AA.BB) or 3G (AA.BBB) specification number ↑	↑ CR number as allocated by MCC support team	
For submission to: <b>TSG RAN #6</b> <i>list expected approval meeting # here</i> ↑	for approval <input checked="" type="checkbox"/> for information <input type="checkbox"/>	strategic <input type="checkbox"/> non-strategic <input type="checkbox"/> <i>(for SMG use only)</i>

Form: CR cover sheet, version 2 for 3GPP and SMG The latest version of this form : <ftp://ftp.3gpp.org/Information/CR-Form-v2.doc>

**Proposed change affects:** (U)SIM  ME  UTRAN / Radio  Core Network   
(at least one should be marked with an X)

**Source:** Nokia, Siemens **Date:** 2 Dec 1999

**Subject:** Modified physical channel mapping scheme

**Work item:** TS25.222

<b>Category:</b>	F Correction <input type="checkbox"/> A Corresponds to a correction in an earlier release <input type="checkbox"/> B Addition of feature <input type="checkbox"/> C Functional modification of feature <input checked="" type="checkbox"/> D Editorial modification <input type="checkbox"/>	<b>Release:</b>	Phase 2 <input type="checkbox"/> Release 96 <input type="checkbox"/> Release 97 <input type="checkbox"/> Release 98 <input type="checkbox"/> Release 99 <input checked="" type="checkbox"/> Release 00 <input type="checkbox"/>
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(only one category shall be marked with an X)

**Reason for change:** The current physical channel mapping scheme does not support the multi-code transmission optimally

**Clauses affected:** 6.2.11.1, 6.2.11.2

**Other specs affected:**

Other 3G core specifications	<input type="checkbox"/>	→ List of CRs:
Other GSM core specifications	<input type="checkbox"/>	→ List of CRs:
MS test specifications	<input type="checkbox"/>	→ List of CRs:
BSS test specifications	<input type="checkbox"/>	→ List of CRs:
O&M specifications	<input type="checkbox"/>	→ List of CRs:

**Other comments:**



<----- double-click here for help and instructions on how to create a CR.

## 6.2.11 Physical channel mapping

The PhCH for both uplink and downlink is defined in [6]. The bits after physical channel mapping are denoted by  $W_{p1}, W_{p2}, \dots, W_{pU_p}$ , where  $p$  is the PhCH number and  $U_p$  is the number of bits in one radio frame for the respective PhCH. The bits  $W_{pk}$  are mapped to the PhCHs so that the bits for each PhCH are transmitted over the air in ascending order with respect to  $k$ . The mapping scheme depends on the applied 2<sup>nd</sup> interleaving scheme.

### 6.2.11.1 Mapping scheme after frame related 2<sup>nd</sup> interleaving

#### 6.2.11.1.1 Mapping scheme after frame related 2<sup>nd</sup> interleaving in uplink

In uplink there are at most two codes allocated ( $P \leq 2$ ). If there is only one code, the same mapping as for downlink is applied, see section 6.2.11.1.2. Denote SF1 and SF2 the spreading factors used for code 1 and 2, respectively. Then denote the inverse relation of the spreading factors  $s1: s2 = SF2: SF1$ , where the smallest possible integers are used for  $s1$  and  $s2$ .

The following mapping rule is applied:

Bits are mapped on the first PhCH (in forward order) if  $(k-1) \bmod (s1+s2) = 0, \dots, s1-1$  after physical channel mapping:

$$W_{1k} = V_k \quad k = 1, 2, \dots, U_1$$

$$W_{1, (k \operatorname{div} (s1+s2)) \cdot s1 + k \operatorname{mod} (s1+s2)} = V_k$$

else  $B$  bits are mapped on the second PhCH (in reverse order) after physical channel mapping:

$$W_{2k} = V_{(k+U_1)} \quad k = U_2, U_2-1, \dots, 1, 2, \dots, U_2$$

$$W_{2, U_2 - (k \operatorname{div} (s1+s2)) \cdot s2 + k \operatorname{mod} (s1+s2) - s1} = V_k$$

... This formula is applied starting with  $k=1$  and increasing  $k$  until one of the PhCH is completely filled. From then on, the remaining bits are mapped on the PhCH which has not been filled in the same order (forward or reverse depending on the PhCH) as used previously on that PhCH.

Bits on the odd numbered  $P^{\text{th}}$  PhCH after physical channel mapping ( $P = 1, 3, 5, \dots$ ):

$$W_{Pk} = V_{(k+U_1+\dots+U_{P-1})} \quad k = 1, 2, \dots, U_P$$

Bits on the even numbered  $P^{\text{th}}$  PhCH after physical channel mapping ( $P = 2, 4, 6, \dots$ ):

$$W_{Pk} = V_{(k+U_1+\dots+U_{P-1})} \quad k = U_P-1, U_P-2, \dots, 1$$

#### 6.2.11.1.2 Mapping scheme after frame related 2<sup>nd</sup> interleaving in downlink

The mapping is equivalent to block interleaving, writing in columns, but a PhCH with an odd number is filled in forward order, were as a PhCH with an even number is filled in reverse order.

The following mapping rule is applied:

Bits are mapped on an odd numbered PhCH (in forward order) according to the following rule, if  $(k \bmod P)+1$  is odd:

$$W_{k \operatorname{mod} P+1, k \operatorname{div} P} = V_k$$

Bits are mapped on an even numbered PhCH (in reverse order) according to the following rule, if  $(k \bmod P)+1$  is even:

$$W_{k \operatorname{mod} P+1, U_P-1-k \operatorname{div} P} = V_k$$

This formula is applied starting with  $k=1$  and increasing  $k$  until all the PhCHs which carry TFCI are completely filled. From then on, the remaining bits are mapped on the remaining PhCHs in the same order (forward or reverse depending on the PhCH) as previously on these PhCHs.

### 6.2.11.2 Mapping scheme after timeslot related 2<sup>nd</sup> interleaving

For each timeslot only those physical channels with  $p = 1, 2, \dots, P_t$  are considered respectively, which are transmitted in that timeslot, and the following mapping scheme is applied:

### 6.2.11.2.1 Mapping scheme after timeslot related 2nd interleaving in uplink

In uplink there are at most two codes allocated ( $P \leq 2$ ). If there is only one code, the same mapping as for downlink is applied, see section 6.2.11.1.2. Denote SF1 and SF2 the spreading factors used for code 1 and 2, respectively. Then denote the inverse relation of the spreading factors  $s1: s2 = SF2: SF1$ , where the smallest possible integers are used for  $s1$  and  $s2$ .

The following mapping rule is applied:

Bits are mapped on the first PhCH (in forward order) if  $(k-1) \bmod (s1+s2) = 0, \dots, s1-1$ :

$$W_{1, (k \operatorname{div} (s1+s2)) \cdot s1 + k \operatorname{mod} (s1+s2)} = V_{tk}$$

else bits are mapped on the second PhCH (in reverse order):

$$W_{2, U_2 - (k \operatorname{div} (s1+s2)) \cdot s2 + k \operatorname{mod} (s1+s2) - s1} = V_{tk}$$

This formula is applied starting with  $k=1$  and increasing  $k$  until one of the PhCH is completely filled. From then on, the remaining bits are mapped on the PhCH which has not been filled in the same order (forward or reverse depending on the PhCH) as used previously on that PhCH.

### 6.2.11.2.2 Mapping scheme after timeslot related 2nd interleaving in downlink

The mapping is equivalent to block interleaving, writing in columns, but a PhCH with an odd number is filled in forward order, were as a PhCH with an even number is filled in reverse order.

The following mapping rule is applied:

Bits are mapped on an odd numbered PhCH (in forward order) according to the following rule, if  $(k \bmod P_t)+1$  is odd:

$$W_{k \bmod P_t + 1, k \operatorname{div} P_t} = V_{tk}$$

Bits are mapped on an even numbered PhCH (in reverse order) according to the following rule, if  $(k \bmod P_t)+1$  is even:

$$W_{k \bmod P_t + 1, U_{P_t} - 1 - k \operatorname{div} P_t} = V_{tk}$$

This formula is applied starting with  $k=1$  and increasing  $k$  until all the PhCHs which carry TFCI are completely filled. From then on, the remaining bits are mapped on the remaining PhCHs in the same order (forward or reverse depending on the PhCH) as previously on these PhCHs.

Bits on first PhCH in timeslot  $t$  after physical channel mapping:

$$W_{1k} = v_{tk} \quad k = 1, 2, \dots, U_1$$

Bits on second PhCH in timeslot  $t$  after physical channel mapping:

$$W_{2k} = v_{t(k+U_1)} \quad k = U_2 - U_2 - 1, \dots, 1, 2, \dots, U_2$$

...

Bits on the odd numbered PhCH  $P_t$  in timeslot  $t$  after physical channel mapping ( $P = 1, 3, 5, \dots$ ):

$$W_{P_t k} = v_{t(k+U_1+\dots+U_{P_t-1})} \quad k = 1, 2, \dots, U_{P_t}$$

Bits on the even numbered  $P_t^{\text{th}}$  PhCH  $P_t$  in timeslot  $t$  after physical channel mapping ( $P = 2, 4, 6, \dots$ ):

$$W_{P_t k} = v_{t(k+U_1+\dots+U_{P_t-1})} \quad k = U_{P_t} - 1, U_{P_t} - 2, \dots, 1$$