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CHANGE REQUEST

⌘ **33.220** **CR** **045** ⌘ **rev** **-** ⌘ Current version: **6.3.0** ⌘

*For **HELP** on using this form, see bottom of this page or look at the pop-up text over the ⌘ symbols.*

Proposed change affects: | UICC apps⌘ ☐ ME ☐ Radio Access Network ☐ Core Network **X**

Title:	⌘	Security capability negotiation in GBA
Source:	⌘	ZTE Corporation
Work item code:	⌘	GBA-SSC
Date:	⌘	27/01/2005
Category:	⌘	**F**
Release:	⌘	Rel-6

Use one of the following categories:
F *(correction)*
A *(corresponds to a correction in an earlier release)*
B *(addition of feature),*
C *(functional modification of feature)*
D *(editorial modification)*
Detailed explanations of the above categories can be found in 3GPP TR 21.900.

Use one of the following releases:

Ph2	*(GSM Phase 2)*
R96	*(Release 1996)*
R97	*(Release 1997)*
R98	*(Release 1998)*
R99	*(Release 1999)*
Rel-4	*(Release 4)*
Rel-5	*(Release 5)*
Rel-6	*(Release 6)*
Rel-7	*(Release 7)*

Reason for change:	⌘	There is no security capability negotiation between UE and NAF in current version, so UE and NAF shall use identical security algorithm that is decided by two entities before communication. Specification's flexibility and extension ability is not so good in this aspect.
Summary of change:	⌘	Adding the procedure of security capability negotiation.
Consequences if not approved:	⌘	It isn't easy to extend GBA to various situations.

Clauses affected:	⌘	4.2.3 4.2.4 4.5.1 4.5.2 4.5.3 5.3.2 5.3.3

Other specs affected: ⌘	Y	N	
		X	Other core specifications
		X	Test specifications
		X	O&M Specifications

⌘

Other comments: ⌘

*** BEGIN OF FIRST CHANGE ***

4.2.3 HSS

The set of all user security settings (USSs), i.e. GUSS, is stored in the HSS. In the case where the subscriber has multiple subscriptions, i.e. multiple ISIM or USIM applications on the UICC, the HSS shall contain one or more GUSSs that can be mapped to one or more private identities, i.e. IMPIs and IMSIs.

The requirements on the HSS are:

- HSS shall provide the only persistent storage for GUSSs;
- GUSS shall be defined in such a way that interworking of different operators for standardised application profiles is possible;
- GUSS shall be defined in such a way that profiles for operator specific applications and extensions to existing application profiles are supported without need for standardisation of these elements.
- GUSS shall be able to contain application-specific USSs that contain parameters that are related to identification or authorization information of one or more applications hosted by one ore more NAFs. Any other types of parameters are not allowed in the application-specific USS.

NOTE 1: The necessary subscriber profile data may be fetched by the NAF directly from HSS or from its local database using identity information provided by the application-specific USS.

NOTE 2: The HSS may temporarily remove an application-specific USS from the GUSS if the service is temporarily revoked from the subscriber.

- GUSS shall be able to contain parameters intended for the BSF usage:
  - the type of the UICC the subscriber is issued (i.e. is it GBA_U aware or not, cf. subclause 5);
  - subscriber specific key lifetime.
  - subscriber specific security grade.

NOTE 3: These parameters are optional and if they are missing from subscriber's GUSS or subsciber does not have GUSS then the BSF will use the default values in the BSF local policy defined by the particular MNO.

NOTE 4: Operator classifies and combines security algorithms supported by almost all UEs and NAFs, transform to security grades based on need of application. Each grade may include an authentication algorithm, an encryption algorithm, and other parameter, it is suitable to protecting certain application. Operator may advise subscriber how to use these grades;

- HSS shall be able to assign application-specific USSs to a NAF Group. This shall be defined in such a way that different USSs for the same application, but for different groups of NAFs, are possible. The restrictions on the number of USSs per GUSS are dependent on the usage of NAF Groups by the operator:
  - if no NAF Groups are defined for this application then at most one USS per application is stored in GUSS;
  - if NAF Groups are defined for this application then at most one USS per application and NAF Group is stored in GUSS.
  - NAF Group definitions in the HSS and all connected BSFs belonging to the same operator's network shall be equal.

4.2.4 UE

The required functionalities from the UE are:

- the support of HTTP Digest AKA protocol;
- the capability to use both a USIM and an ISIM in bootstrapping;

- the capability to select either a USIM or an ISIM to be used in bootstrapping, when both of them are present;
- the capability for a Ua application on the ME to indicate to the GBA Function on the ME the type or the name of UICC application to use in bootstrapping (see clause 4.4.8);
- the capability to derive new key material to be used with the protocol over Ua interface from CK and IK;
- support of NAF-specific application protocol (For an example see TS 33.221 [5]).
- support of configuration of security grade.

A GBA-aware ME shall support both GBA_U, as specified in clause 5.2.1 and GBA_ME procedures, as specified in clause 4.5.

NOTE: There are security grades on UE, the method of definition is the same as operator's. The grade list on UE is a subset of grades of operator's, and shall be stored in USS; Subscriber configures security grade based on his/her need of application, the grade could be a single value, or a range with priority;

*** END OF FIRST CHANGE ***

*** BEGIN OF SECOND CHANGE ***

4.5.1 Initiation of bootstrapping

Before communication between the UE and the NAF can start, the UE and the NAF first have to agree whether to use the GBA. When a UE wants to interact with a NAF, but it does not know if the NAF requires the use of shared keys obtained by means of the GBA, the UE shall contact the NAF for further instructions (see figure 4.2).

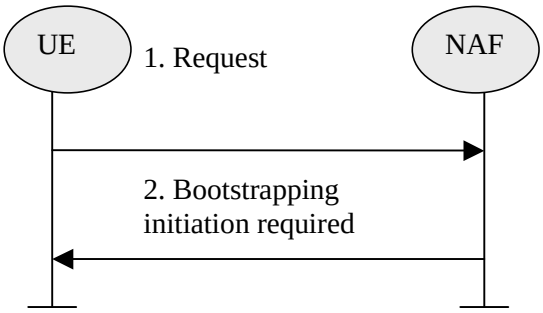


Figure 4.2: Initiation of bootstrapping

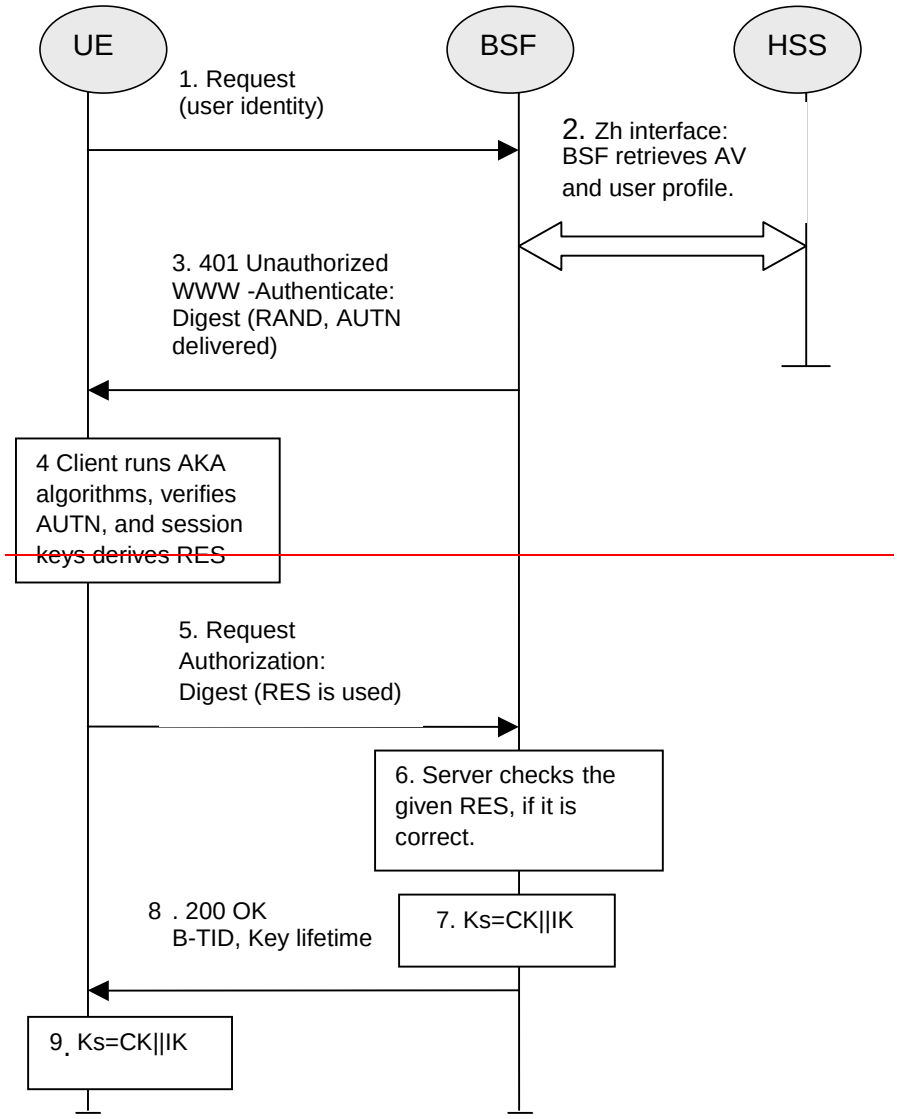
1. UE starts communication over reference point Ua with the NAF without any GBA-related parameters.
2. If the NAF requires the use of shared keys obtained by means of the GBA,but the request from UE does not include GBA-related parameters, the NAF replies with a bootstrapping initiation message. In order to negotiate security capability with UE, the response message should include integrity and encryption algorithm list supported by NAF, ordered by priority. The form of this indication may depend on the particular reference point Ua and is specified in the relevant stage 3-specifications.

4.5.2 Bootstrapping procedures

When a UE wants to interact with a NAF, and it knows that the bootstrapping procedure is needed, it shall first perform a bootstrapping authentication (see figure 4.3). Otherwise, the UE shall perform a bootstrapping authentication only

when it has received bootstrapping initiation required message or a bootstrapping negotiation indication from the NAF, or when the lifetime of the key in UE has expired (cf. subclause 4.5.3).

NOTE 1: The main steps from the specifications of the AKA protocol in TS 33.102 [2] and the HTTP digest AKA protocol in RFC 3310 [4] are repeated in figure 3 for the convenience of the reader. In case of any potential conflict, the specifications in TS 33.102 [2] and RFC 3310 [4] take precedence.


UE
BSF
HSS
1. Request
(user identity)
2. Zh interface:
BSF retrieves AV
and user profile.
3. 401 Unauthorized
WWW -Authenticate:
Digest (RAND, AUTN
delivered)
4 Client runs AKA
algorithms, verifies
AUTN, and session
keys derives RES
5. Request
Authorization:
Digest (RES is used)
6. Server checks the
given RES, if it is
correct.
8 . 200 OK
B-TID, Key lifetime
7. Ks=CK||IK
9. Ks=CK||IK

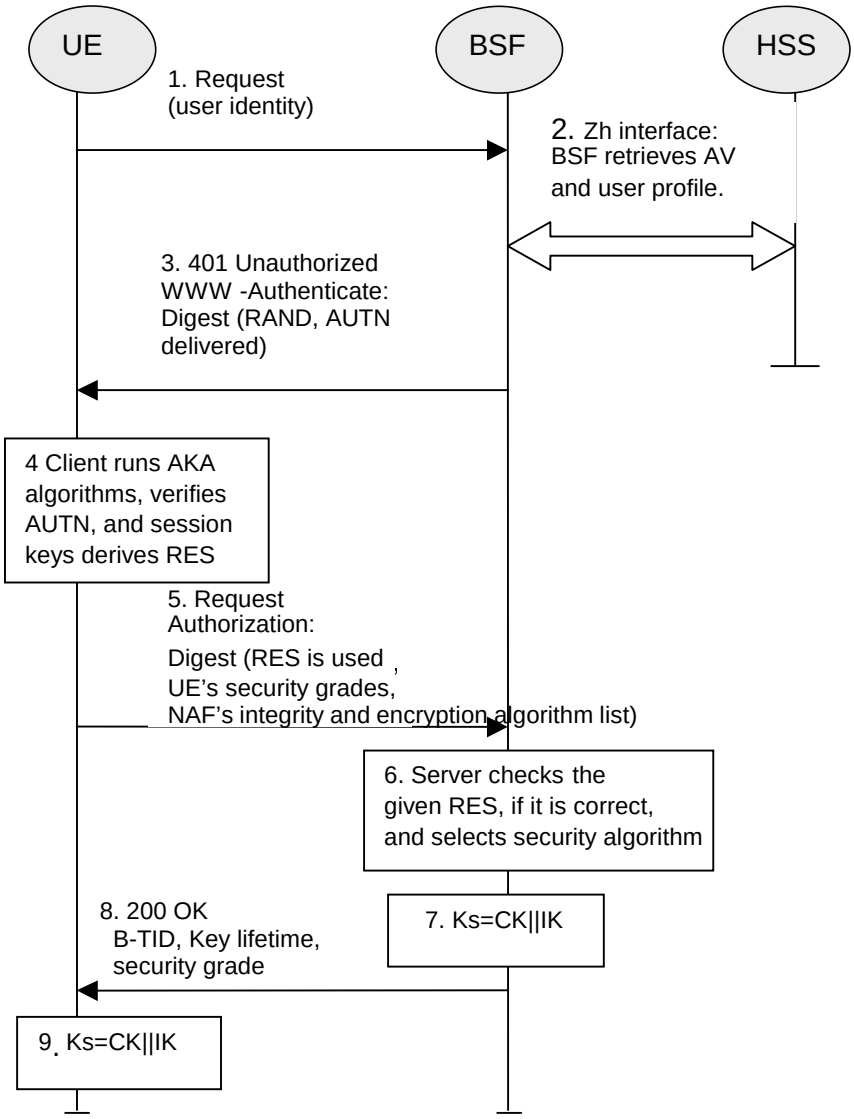


Figure 4.3: The bootstrapping procedure

1. The UE sends an HTTP request towards the BSF.
2. BSF retrieves the complete set of GBA user security settings and one Authentication Vector (AV, AV = RAND||AUTN||XRES||CK||IK) over the reference point Zh from the HSS.
3. Then BSF forwards the RAND and AUTN to the UE in the 401 message (without the CK, IK and XRES). This is to demand the UE to authenticate itself.
4. The UE checks AUTN to verify that the challenge is from an authorised network; the UE also calculates CK, IK and RES. This will result in session keys IK and CK in both BSF and UE.
5. The UE sends another HTTP request, containing the Digest AKA response (calculated using RES), to the BSF. In order to negotiate security capability with NAF, the HTTP request should include the security grades configured by user, and NAF's security algorithm list.
6. The BSF authenticates the UE by verifying the Digest AKA response, and selects the first algorithm combination on NAF's list that is also supported by the UE. If BSF doesn't know the relation of UE's security grade and security algorithms, it could get it from USS.

7. The BSF generates key material Ks by concatenating CK and IK. The B-TID value shall be also generated in format of NAI by taking the base64 encoded [12] RAND value from step 3, and the BSF server name, i.e. base64encode(RAND)@BSF_servers_domain_name.

8. The BSF shall send a 200 OK message, including a B-TID and selected security grade, to the UE to indicate the success of the authentication. In addition, in the 200 OK message, the BSF shall supply the lifetime of the key Ks. The key material Ks is generated in UE by concatenating CK and IK.

9. Both the UE and the BSF shall use the Ks to derive the key material Ks_NAF during the procedures as specified in clause 4.5.3. Ks_NAF shall be used for securing the reference point Ua. UE shall use security grade sent by BSF communicate with NAF.

   Ks_NAF is computed as Ks_NAF = KDF (Ks, "gba-me" || RAND || IMPI || NAF_Id), where KDF is the key derivation function as specified in Annex B, and the key derivation parameters consist of the user's IMPI, the NAF_Id and RAND. The NAF_Id consists of the full DNS name of the NAF. KDF shall be implemented in the ME.

NOTE 2: To allow consistent key derivation based on NAF name in UE and BSF, at least one of the three following prerequisites shall be fulfilled:

(1) The NAF is known in DNS under one domain name (FQDN) only, i.e. no two different domain names point to the IP address of the NAF. This has to be achieved by administrative means.
This prerequisite is not specific to 3GPP, as it is necessary also under other circumstances, e.g. for TLS V1.0 without use of wildcard or multiple-name certificates.

(2) Each DNS entry of the NAF points to a different IP address. The NAF responds to all these IP addresses. Each IP address is tied to the corresponding FQDN by NAF configuration. The NAF can see from the IP address, which FQDN to use for key derivation.

(3) Ua uses a protocol which transfers the host name (FQDN of NAF as used by UE) to NAF (e.g. HTTP/1.1 with mandatory Host request header field). This requires the NAF to check the validity of the host name, to use this name in all communication with UE where appropriate, and to transfer this name to BSF to allow for correct derivation of Ks_NAF.
In case of a TLS tunnel this requires either multiple-identities certificates or the deployment of RFC 3546 [9] or other protocol means with similar purpose.

The UE and the BSF shall store the key Ks with the associated B-TID for further use, until the lifetime of Ks has expired, or until the key Ks is updated.

4.5.3 Procedures using bootstrapped Security Association

Before communication between the UE and the NAF can start, the UE and the NAF first have to agree whether to use shared keys obtained by means of the GBA. If the UE does not know whether to use GBA with this NAF, it uses the Initiation of Bootstrapping procedure described in clause 4.5.1.

Once the UE and the NAF have established that they want to use GBA then every time the UE wants to interact with an NAF the following steps are executed as depicted in figure 4.4.

1. UE starts communication over reference point Ua with the NAF:
   - in general, UE and NAF will not yet share the key(s) required to protect the reference point Ua. If they already do (i.e. if a key Ks_NAF for the corresponding key derivation parameter NAF_Id is already available),, the UE and the NAF can start to securely communicate right away. If the UE and the NAF do not yet share a key, the UE proceeds as follows:

   - if a key Ks for the selected UICC application is available in the UE, the UE derives the key Ks_NAF from Ks, as specified in clause 4.5.2;

   - if no key Ks for the selected UICC application is available in the UE, the UE first agrees on a new key Ks with the BSF over the reference point Ub, and then proceeds to derive Ks_NAF;

NOTE 1: If it is not desired by the UE to use the same Ks for the selected UICC application to derive more than one Ks_NAF then the UE should agree on a new key Ks with the BSF over the reference point Ub, and then proceed to derive Ks_NAF;

- if the NAF shares a key with the UE, but the NAF requires an update of that key, e.g. because the key's lifetime has expired or will expire soon, or the key can not meet the NAF local validity condition, it shall send a suitable bootstrapping renegotiation request to the UE, see figure 4.5. If the key's lifetime has expired the protocol used over reference point Ua shall be terminated. The form of this indication depends on the particular protocol used over reference point Ua. If the UE receives a bootstrapping renegotiation request, it starts a run of the protocol over reference point Ub, as specified in clause 4.5.2, in order to obtain a new key Ks.

NOTE 2: To allow for consistent key derivation in BSF and UE, both have to use the same FQDN for derivation (see NOTE 2 of clause 4.5.2). For each protocol used over Ua it shall be specified if only cases (1) and (2) of NOTE 2 of clause 4.5.2 are allowed for the NAF or if the protocol used over Ua shall transfer also the FQDN used for key derivation by UE to NAF.

NOTE 3: If the shared key between UE and NAF is invalid , the NAF can set deletion conditions to the corresponding security association for subsequent removal.

- the UE supplies the B-TID to the NAF, in the form as specified in clause 4.3.2, to allow the NAF to retrieve the corresponding keys from the BSF;

NOTE 4: The UE may adapt the key material Ks_NAF to the specific needs of the reference point Ua. This adaptation is outside the scope of this specification.

- key management for GBA related keys in the ME (i.e. Ks and Ks_NAF keys):
  - all GBA related keys shall be deleted from the ME when a different UICC is inserted. Therefore the ME needs to store in non-volatile memory the last inserted UICC-identity to be able to compare that with the used UICC-identity at UICC insertion and power on;
  - the Key Ks shall be deleted from the ME when the ME is powered down;
  - all other GBA related keys may be deleted from the ME when the ME is powered down. If the ME does not delete the GBA keys at power down then the GBA keys need to be stored in non-volatile memory.
- when a new Ks is agreed over the reference point Ub and a key Ks_NAF, derived from one NAF_Id, is updated, the other keys Ks_NAF, derived from different values NAF_Id, stored on the UE shall not be affected;

NOTE 5: According to the procedures defined in clauses 4.5.2 and 4.5.3, in the UE there is at most one Ks_NAF key stored per NAF-Id.

2. NAF starts communication over reference point Zn with BSF
   - The NAF requests key material corresponding to the B-TID supplied by the UE to the NAF over reference point Ua. If the NAF has several FQDNs, which may be used in conjunction with this specification, then the NAF shall transfer in the request over Zn the same FQDN, which was used over Ua (see NOTE 2 on key derivation in this clause);
   - The NAF may also request one or more application-specific USSs for the applications, which the request received over Ua from UE may access;
   - With the key material request, the NAF shall supply NAF's public hostname that UE has used to access NAF to BSF, and BSF shall be able verify that NAF is authorized to use that hostname;
3. The BSF derives the keys required to protect the protocol used over reference point Ua from the key Ks and the key derivation parameters, as specified in clause 4.5.2, and supplies to NAF the requested key Ks_NAF, as well as the lifetime of that key, the integrity and encryption algorithm selected by BSF, and the requested application-specific and potentially NAF group specific USSs if they are available in subscriber's GUSS and if the NAF is authorized to receive the requested USSs. If the key identified by the B-TID supplied by the NAF is not available at the BSF, the BSF shall indicate this in the reply to the NAF. The NAF then indicates a bootstrapping renegotiation request to the UE.

NOTE 6: The NAF can further set the local validity condition of the Ks_NAF according to the local policy,for example a limitation of reuse times of a Ks_NAF.

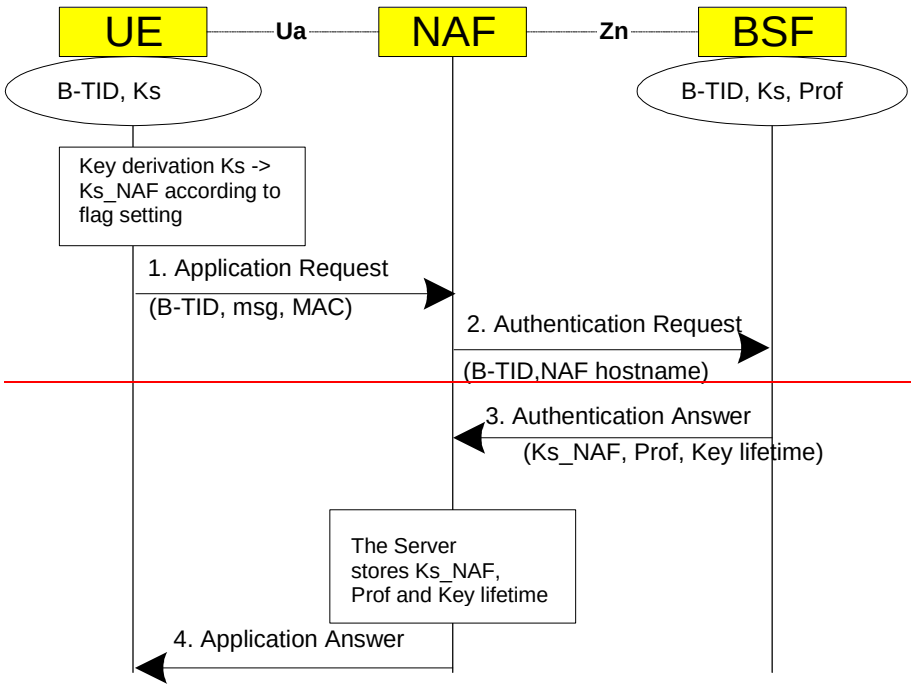
NOTE 7: The NAF shall adapt the key material Ks_NAF to the specific needs of the reference point Ua in the same way as the UE did. This adaptation is outside the scope of this specification.

- NAF shall use security algorithms sent by BSF to communicate with UE.

- The BSF may require that one or more application-specific and potentially NAF group specific USSs shall be present in subscriber's GUSS for the NAF (see clause 4.4.6). If one or more of these required settings are missing from the GUSS, the BSF shall indicate this in the reply to the NAF.

  - The BSF may also send the private user identity (IMPI) and requested USSs to NAF according to the BSF's policy;

4. NAF continues with the protocol used over the reference point Ua with the UE.
   Once the run of the protocol used over reference point Ua is completed the purpose of bootstrapping is fulfilled as it enabled UE and NAF to use reference point Ua in a secure way.



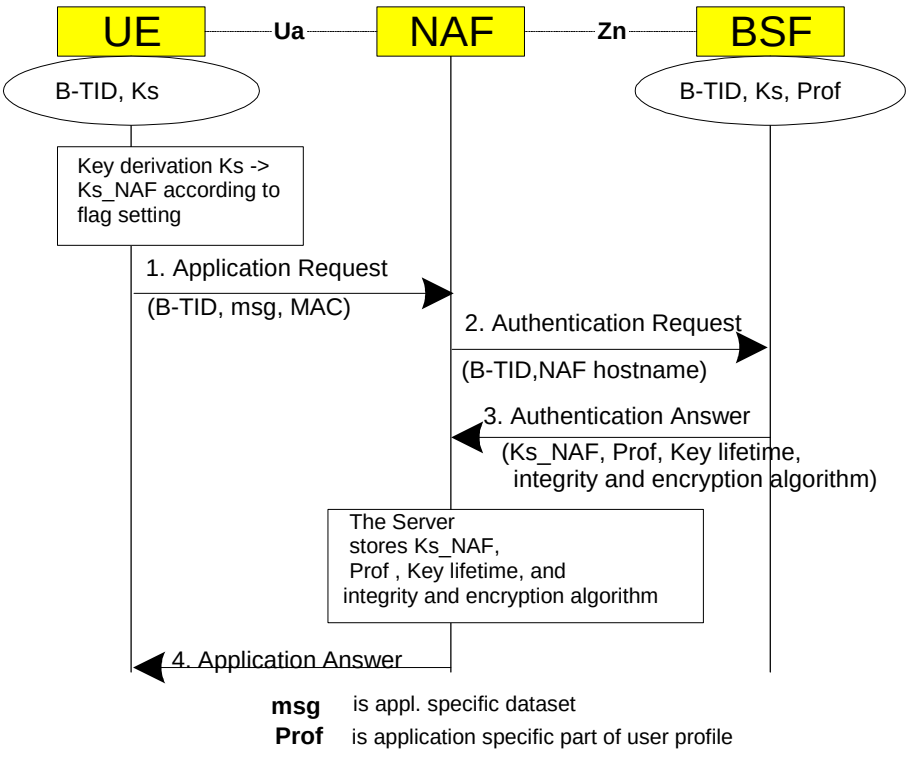


Figure 4.4: The bootstrapping usage procedure

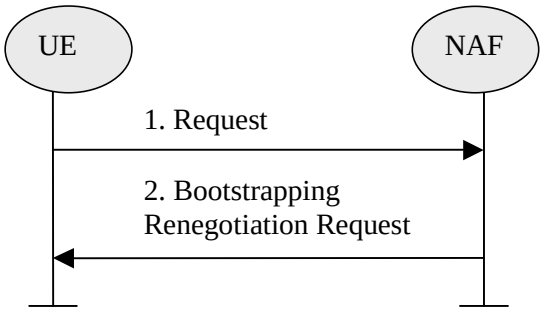


Figure 4.5: Bootstrapping renegotiation request

*** END OF SECOND CHANGE ***

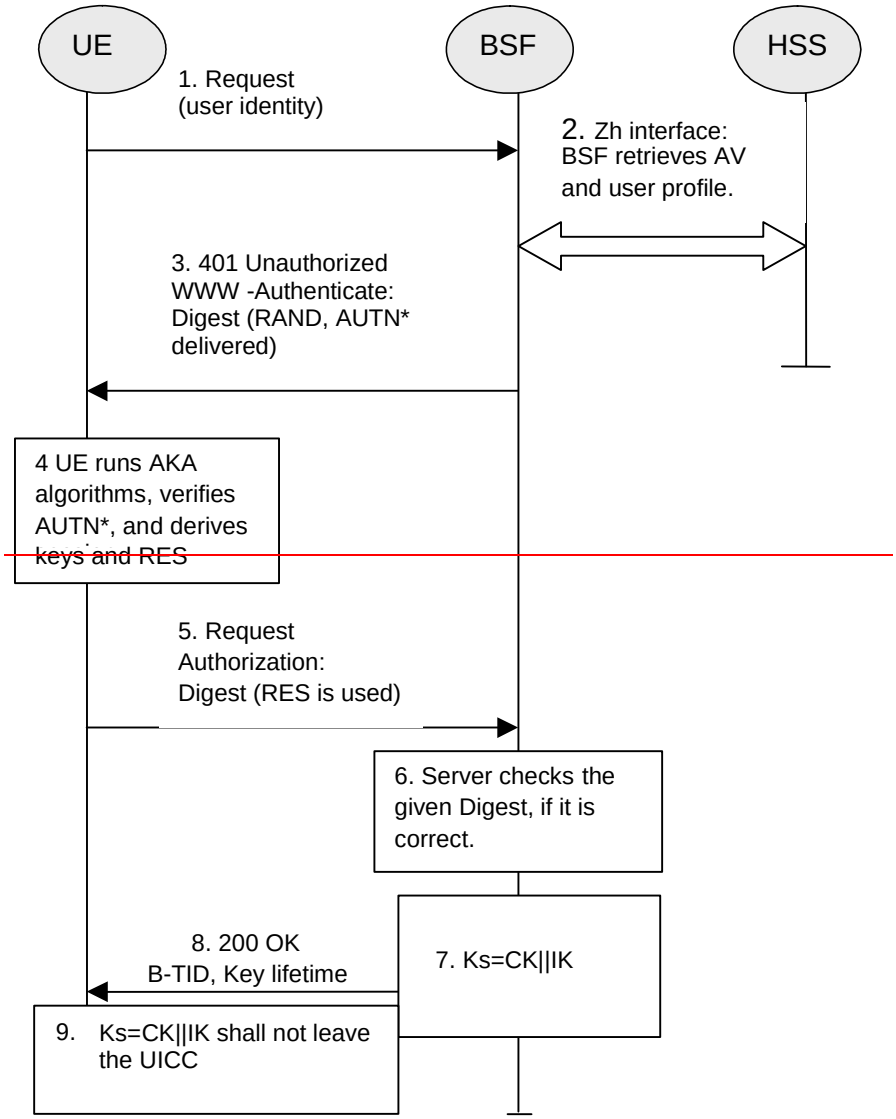
*** BEGIN OF THIRD CHANGE ***

5.3.2 Bootstrapping procedure

The procedure specified in this clause differs from the procedure specified clause 4.5.2 in the local handling of keys and Authentication Vectors in the UE and the BSF. The messages exchanged over the Ub reference point are identical for both procedures.

When a UE wants to interact with a NAF, and it knows that the bootstrapping procedure is needed, it shall first perform a bootstrapping authentication (see figure 5.1). Otherwise, the UE shall perform a bootstrapping authentication only when it has received bootstrapping initiation required message or a bootstrapping renegotiation indication from the NAF, or when the lifetime of the key in UE has expired (see clause 5.3.3).

NOTE: The main steps from the specifications of the AKA protocol in TS 33.102 [2] and the HTTP digest AKA protocol in RFC 3310 [4] are repeated in figure 5.1 for the convenience of the reader. In case of any potential conflict, the specifications in TS 33.102 [2] and RFC 3310 [4] take precedence.


UE
BSF
HSS
1. Request
(user identity)
2. Zh interface:
BSF retrieves AV
and user profile.
3. 401 Unauthorized
WWW -Authenticate:
Digest (RAND, AUTN*
delivered)
4 UE runs AKA
algorithms, verifies
AUTN*, and derives
keys and RES
5. Request
Authorization:
Digest (RES is used)
6. Server checks the
given Digest, if it is
correct.
7. Ks=CK||IK
8. 200 OK
B-TID, Key lifetime
9. Ks=CK||IK shall not leave
the UICC

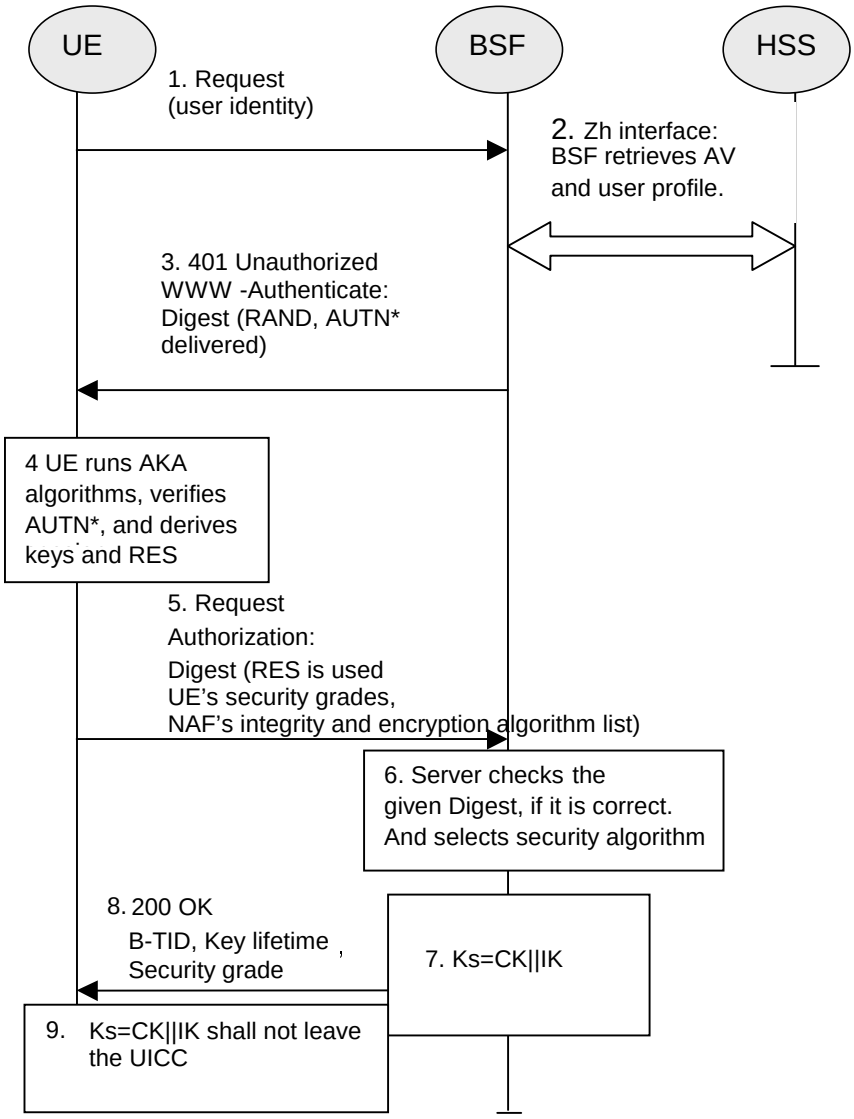


Figure 5.1: The bootstrapping procedure with UICC-based enhancements

1. The ME sends an HTTP request towards the BSF.

2. The BSF retrieves the complete set of GBA user security settings and one Authentication Vector (AV, AV = RAND||AUTN||XRES||CK||IK) over the Zh reference point from the HSS. The BSF can then decide to perform GBA_U, based on the user security settings (USSs). In this case, the BSF proceeds in the following way:

- BSF computes MAC* = MAC⊕ Trunc(SHA-1(IK))

NOTE: Trunc denotes that from the 160 bit output of SHA-1 [21], the 64 bits numbered as [0] to [63] are used within the * operation to MAC.

The BSF stores the XRES after flipping the least significant bit.

3. Then BSF forwards the RAND and AUTN* (where AUTN* = SQN ⊕ AK || AMF || MAC*) to the UE in the 401 message (without the CK, IK and XRES). This is to demand the UE to authenticate itself.

4. The ME sends RAND and AUTN* to the UICC. The UICCcalculates IK and MAC (by performing MAC= MAC* ⊕ Trunc(SHA-1(IK))). Then the UICC checks AUTN(i.e. SQN ⊕ AK || AMF || MAC) to verify that the challenge is from an authorised network; the UICC also calculates CK and RES. This will result in session keys

CK and IK in both BSF and UICC. The UICC then transfers RES (after flipping the least significant bit) to the ME and stores Ks, which is the concatenation of CK and IK, on the UICC.

5. The ME sends another HTTP request, containing the Digest AKA response (calculated using RES), to the BSF. In order to negotiate security capability with NAF, the HTTP request should include the security grades configured by user, and NAF's security algorithm list.

6. The BSF authenticates the UE by verifying the Digest AKA response. and selects the first algorithm combination on NAF's list that is also supported by the UE. If BSF doesn't know the relation of UE's security grade and security algorithms, it could get it from USS.

7. The BSF generates the key Ks by concatenating CK and IK. The B-TID value shall be also generated in format of NAI by taking the base64 encoded [12] RAND value from step 3, and the BSF server name, i.e. base64encode(RAND)@BSF_servers_domain_name.

8. The BSF shall send a 200 OK message, including the B-TID and selected security grade, to the UE to indicate the success of the authentication. In addition, in the 200 OK message, the BSF shall supply the lifetime of the key Ks.

9. Both the UICC and the BSF shall use the Ks to derive NAF-specific keys Ks_ext_NAF and Ks_int_NAF during the procedures as specified in clause 5.3.3, if applicable. Ks_ext_NAF and Ks_int_NAF are used for securing the Ua reference point. UE shall use security grade sent by BSF communicate with NAF.

   Ks_ext_NAF is computed in the UICC as Ks_ext_NAF = KDF(Ks, h1-key derivation parameters), and Ks_int_NAF is computed in the UICC as Ks_int_NAF = KDF(Ks, h1-key derivation parameters), where KDF is the key derivation function as specified in Annex B, and the key derivation parameters include the user's IMPI, the NAF_Id and RAND. The NAF_Id consists of the full DNS name of the NAF. The key derivation parameters used for Ks_ext_NAF derivation must be different from those used for Ks_int_NAF derivation. This is done by adding a static string "gba-me" in Ks_ext_NAF and "gba-u" in Ks_int_NAF as an input parameter to the key derivation function.

NOTE: The NOTE 2 of clause 4.5.2 also applies here.

The UICC and the BSF store the key Ks with the associated B-TID for further use, until the lifetime of Ks has expired, or until the key Ks is updated.

5.3.3 Procedures using bootstrapped Security Association

Before communication between the UE and the NAF can start, the UE and the NAF first have to agree whether to use shared keys obtained by means of the GBA. If the UE does not know whether to use GBA with this NAF, it uses the Initiation of Bootstrapping procedure described in clause 5.3.1.

Once the UE and the NAF have established that they want to use GBA then every time the UE wants to interact with a NAF the following steps are executed as depicted in figure 5.3.

Next, the UE and the NAF have to agree, which type of keys to use, Ks_ext_NAF or Ks_int_NAF, or both. The default is the use of Ks_ext_NAF only. This use is also supported by MEs and NAFs, which are GBA_U unaware. If Ks_int_NAF, or both Ks_ext_NAF and Ks_int_NAF are to be used, this use has to be agreed between UE and NAF prior to the execution of the procedure described in the remainder of this clause 5.3.3. Any such agreement overrules the default use of the keys. How this agreement is reached is application-specific and is not within the scope of this document.

NOTE 1: This agreement may be mandated by the specification, which defines the Ua reference point between UE and NAF, e.g. TS 33.246 for the use of GBA in MBMS, or negotiated by the NAF and the UE over the Ua reference point, or reached by configuration.

In general, UE and NAF will not yet share the key(s) required to protect the Ua reference point. If they do not, the UE proceeds as follows:

- if Ks_ext_NAF is required and a key Ks for the selected UICC application is available in the UICC, the ME requests the UICC to derive the key Ks_ext_NAF from Ks, as specified in clause 5.3.2;

- if Ks_int_NAF is required and a key Ks for the selected UICC application is available in the UICC, the ME requests the UICC to derive the key Ks_int_NAF from Ks, as specified in clause 5.3.2;

NOTE 2: If it is not desired by the UE to use the same Ks for the selected UICC application to derive more than one Ks_ext/int_NAF then the UE should first agree on new key Ks with the BSF over the Ub reference point, as specified in clause 5.3.2, and then proceeds to derive Ks_ext_NAF or Ks_int_NAF, or both, as required.

- if Ks for the selected UICC application is not available in the UE, the UE first agrees on a new key Ks with the BSF over the Ub reference point, as specified in clause 5.3.2, and then proceeds to derive Ks_ext_NAF or Ks_int_NAF, or both, as required;

- if the NAF shares a key with the UE, but the NAF requires an update of that key, it shall send a suitable bootstrapping renegotiation request to the UE. If the key's lifetime has expired the protocol used over reference point Ua shall be terminated. The form of this indication depends on the particular protocol used over Ua reference point. If the UE receives a bootstrapping renegotiation request, it starts a run of the protocol over Ub, as specified in clause 5.3.2, in order to obtain new keys.

NOTE 3: If the shared keys between UE and NAF become invalid, the NAF can set deletion conditions to the corresponding security association for subsequent removal.

NOTE 4: If it is not desired by the NAF to use the same Ks to derive more than one Ks_int/ext_NAF then the NAF should always reply to the first request sent by a UE by sending a key update request to the UE.

1. UE and NAF can now start the communication over Ua reference point using the keys Ks_ext_NAF or Ks_int_NAF, or both, as required. They proceed as follows:

   - The UE supplies the B-TID to the NAF, as specified in clause 5.3.2, to allow the NAF to retrieve the corresponding keys from the BSF

   NOTE 5: To allow for consistent key derivation in BSF and UE, both have to use the same FQDN for derivation (cf. NOTE 2 of clause 4.5.2). For each protocol used over Ua it shall be specified if only cases (1) and (2) of NOTE 2 of clause 4.5.2 are allowed for the NAF or if the protocol used over Ua shall transfer also the FQDN used for key derivation by UE to NAF.

   NOTE 6: The UE may adapt the keys Ks_ext_NAF or Ks_int_NAF to the specific needs of the Ua reference point. This adaptation is outside the scope of this specification.

   - key management for GBA related keys in the ME (i.e. Ks_ext_NAF keys):

   - all GBA related keys shall be deleted from the ME when a different UICC is inserted. Therefore the ME needs to store in non-volatile memory the last inserted UICC-identity to be able to compare that with the used UICC-identity at UICC insertion and power on.

   - all GBA related keys may be deleted from the ME when the ME is powered down. If the ME does not delete the GBA keys at power down then the GBA keys need to be stored in non-volatile memory.

   - all GBA related keys in the UICC do not need to be deleted when the ME is powered down.

   NOTE 7: After each run of the protocol over the Ub reference point, a new key Ks, associated with a new B-TID, are derived in the UE according to clause 5.3.2, so that it can never happen, that key Ks with different B-TIDs simultaneously exist in the UE.

   - When new key Ks is agreed over the Ub reference point and new NAF-specific keys need to be derived for one NAF_Id, then both, Ks_ext_NAF and Ks_int_NAF (if present), shall be updated for this NAF_Id, but further keys Ks_ext_NAF or Ks_int_NAF relating to other NAF_Ids, which may be stored on the UE, shall not be affected.

   NOTE 8: According to the procedures defined in clauses 5.3.2 and 5.3.3, in the UE there is at most one Ks_int_NAF/Ks_ext_NAF key pair stored per NAF_Id.

   NOTE 9: This rule ensures that the keys Ks_ext_NAF and Ks_int_NAF are always in synch at the UE and the NAF.

2. NAF now starts communication over the Zn reference point with the BSF.

   - The NAF requests from the BSF the keys corresponding to the B-TID, which was supplied by the UE to the NAF over the Ua reference point. If the NAF is GBA_U aware it indicates this by including a corresponding flag in the request. If the NAF has several FQDNs, which may be used in conjunction with this specification, then the

NAF shall transfer in the request over Zn the same FQDN, which was used over Ua (see note above on key derivation in this clause).

- The NAF may also request one or more application-specific USSs for the applications, which the request received over Ua from UE may access;

- With the keys request over the Zn reference point, the NAF shall supply NAF's public hostname that UE has used to access NAF to BSF, and BSF shall be able to verify that NAF is authorized to use that hostname.

3. The BSF derives the keys Ks_ext_NAF, and Ks_int_NAF (if additionally required), as specified in clause 5.3.2. If the NAF indicated in its request that it is GBA_U aware, the BSF supplies to NAF both keys, Ks_ext_NAF, and Ks_int_NAF, otherwise the BSF supplies only Ks_ext_NAF. In addition, the BSF supplies the lifetime time of these keys, the integrity and encryption algorithm selected by BSF, and the requested application-specific and potentially NAF group specific USSs if they are available in subscriber's GUSS and if the NAF is authorized to receive the requsted USSs. If the key identified by the B-TID supplied by the NAF is not available at the BSF, the BSF shall indicate this in the reply to the NAF. The NAF then indicates a bootstrapping renegotiation request (See figure 4.5) to the UE;

NOTE 10: The NAF can further set the local validity condition of the Ks_NAF according to the local policy, for example a limitation of reuse times of a Ks_NAF.

NOTE 11: The NAF may adapt the keys Ks_ext_NAF and Ks_int_NAF to the specific needs of the Ua reference point in the same way as the UE did. This adaptation is outside the scope of this specification.

- NAF shall use security algorithms sent by BSF to communicate with UE.

- The BSF may require that one or more application-specific and potentially NAF group specific USSs shall be present in subscriber's GUSS for the NAF (see clause 4.4.6). If one or more of these required settings are missing from the GUSS, the BSF shall indicate this in the reply to the NAF.

- The BSF may also send the private user identity (IMPI) and requested USSs to NAF according to the BSF's policy.

4. The NAF now continues with the protocol used over the Ua reference point with the UE.

Once the run of the protocol used over Ua reference point is completed the purpose of bootstrapping is fulfilled as it enabled the UE and NAF to use Ua reference point in a secure way.

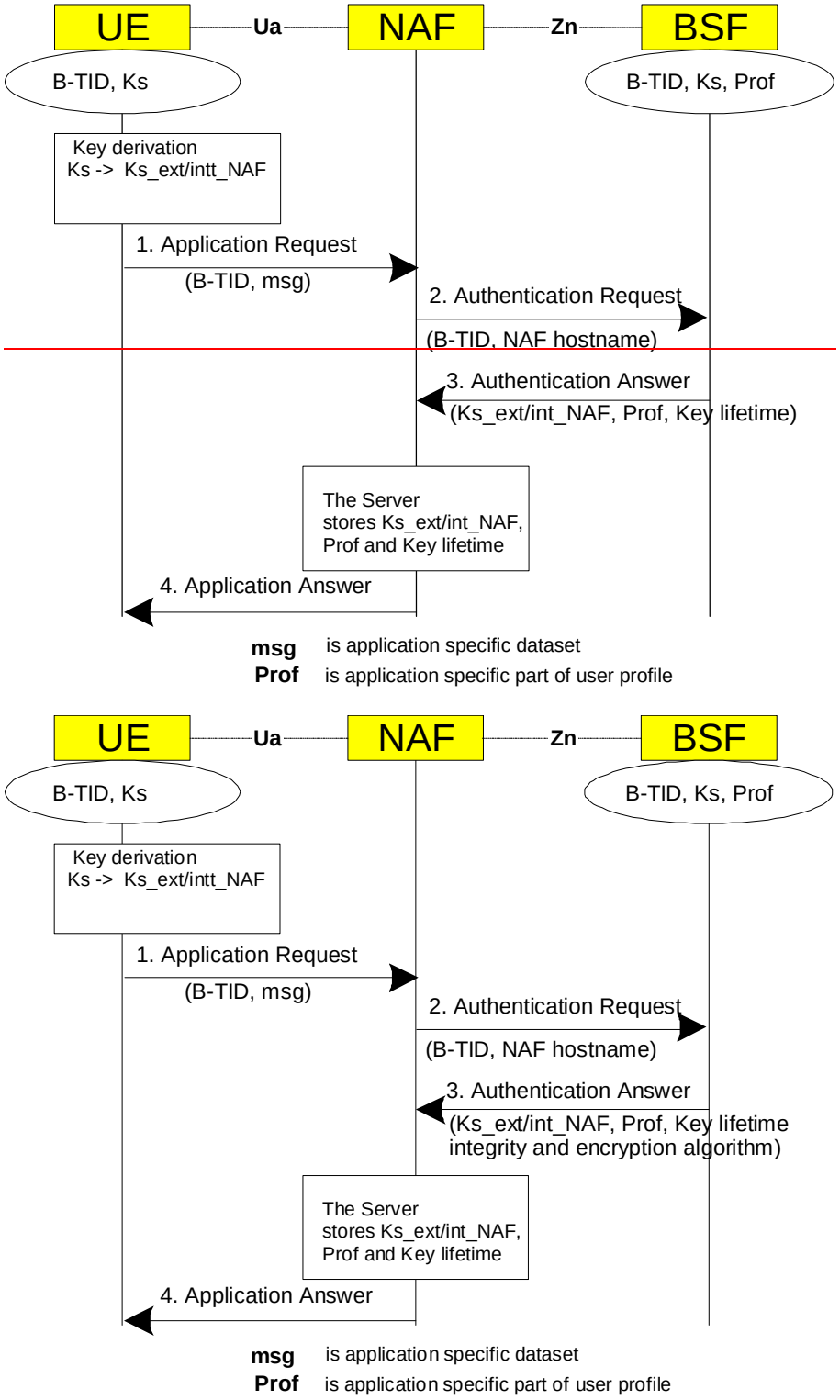


Figure 5.3: The bootstrapping usage procedure with UICC-based enhancements

*** END OF THIRD CHANGE ***