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Source: Gemplus, Axalto, Oberthur

Title: GBA: GBA\_U derivations

Document for: Discussion and decision

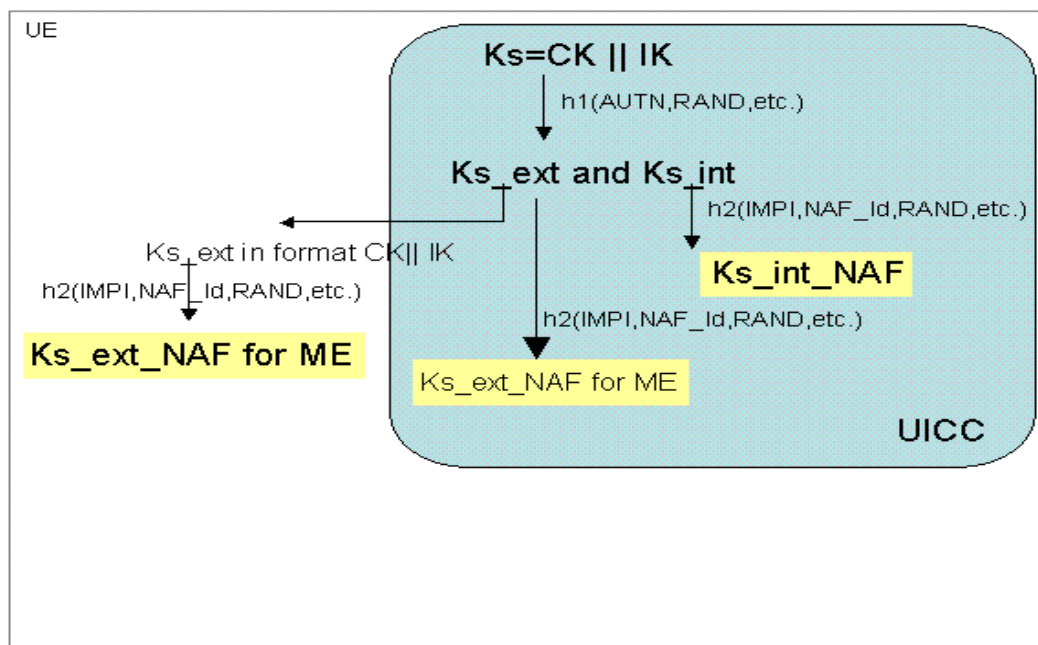
Agenda Item: GBA

## 1. Introduction

The current version of TS 33.220 “Generic Bootstrapping Architecture” mentions that the location (whether in the UICC or in the ME) of the storage of  $Ks_{ext}$  is for further study. This contribution proposes some elements of comparison in order to select the location of  $Ks_{ext}$  storage.

## 2. Current status of GBA-U key derivation

The current TS 33.220 [1] proposes the following key derivations for GBA-U:



The TS states that:

- The location (whether in the UICC or in the ME) of the storage of Ks\_ext is for further study.
  
- When the UE is powered down, or when the UICC is removed, any GBA\_U keys shall be deleted from storage in the ME. There is no need to delete the keys Ks\_int and Ks\_int\_NAF from storage in the UICC.

### **3. Elements of comparison**

During SA3#33 meeting some elements of comparison were provided in order to select the location of the storage of Ks\_ext for GBA-U.

#### **Ks\_ext key lifetime:**

- Scenario 1: Ks\_ext stored on the ME  
Ks\_ext shall be deleted when the UE is powered down or when the UICC is removed.
  
- Scenario 2: Ks\_ext stored on the UICC  
There is no need to delete Ks\_int and Ks\_ext from the UICC in case of UE power down or UICC removal. Only, the Ks\_ext\_NAF key stored on the ME shall be deleted.

So, the key lifetime of Ks\_ext is longer with K\_ext storage on the UICC, this leads to:

- Decrease the frequency of bootstrapping procedures
- Decrease the consumption of authentication vectors

#### **Ks\_ext availability**

Ks\_ext\_NAF are derived from Ks\_ext, they are computed as  $Ks\_ext\_NAF = h2(Ks\_ext, h2\text{-key derivation parameters})$  where the h2-key derivation parameters include IMPI, NAF\_Id and RAND.

- Scenario 1: Ks\_ext stored on the ME  
Attacks on ME are possible, so the retrieval of Ks\_ext, Ks\_ext\_NAF keys, IMPI and RAND is possible. The same value of IMPI and RAND is used as h2-key derivation parameters for all Ks\_ext\_NAF computation. So, an attacker, who accesses one time Ks\_ext, IMPI and RAND on the ME during the key lifetime of Ks\_ext, is able to compute/deduce Ks\_ext\_NAF of any NAF since he only requires the NAF\_ID value to perform  $h2(Ks\_ext, IMPI, NAF\_ID, RAND)$ . The attacker could use these Ks\_ext\_NAFs to send authenticated application requests to new NAFs.

Moreover, the level of security of Ks\_ext\_NAF with GBA\_U is the same than the security level of Ks\_NAF with GBA-ME.

- Scenario 2: Ks\_ext stored on the UICC  
Attacks on ME allow retrieving the IMPI and RAND values used as h2-key derivation parameters and all the Ks\_ext\_NAF keys available on the ME. But, an attacker cannot compute/deduce the value of a Ks\_ext\_NAF corresponding to a new NAF, since the attacker does not know Ks\_ext. So, he cannot send an authenticated application request to new NAFs.

So, the security level is higher in case of Ks\_ext storage on the UICC.

### **ME not implementing GBA-U**

- **Scenario 1: Ks\_ext stored on the ME**  
A GBA-aware UICC gives Ks\_ext to the ME. This scenario allows handling the situation where an ME (with only GBA\_ME function implemented) does not know that a GBA-aware UICC has been inserted.
- **Scenario 2: Ks\_ext stored on the UICC**  
A GBA-U aware UICC does not provide Ks\_ext to the ME and sends Ks\_ext\_NAF during the second run of the bootstrapping procedure. A ME, which does not support GBA\_U capabilities, could not derive the correct Ks\_ext\_NAF.

This issue only exists in case of Rel-6 GBA-capable MEs not supporting GBA-U. But this scenario is not yet decided and depends on SA3 decision, a SA3#34 contribution on “GBA\_U in Rel 6 MEs » [2] is proposed for discussion. So, if SA3 decide that Rel-6 GBA-capable MEs shall support both GBA-ME and GBA-U, the issue described in previous paragraph is not longer relevant.

## **4. Conclusion**

In case of Rel-6 ME supporting both GBA\_U and GBA\_ME, the location of Ks\_ext storage in the UICC provides a higher level of security for Ks\_ext\_NAF usage and extends the key life time of Ks\_ext.

So, we kindly recommend SA3 to require the storage of Ks\_ext on the UICC for GBA\_U.  
A CR [3] implements this proposal.

## **5. References**

- [1] TS 33.220, v6.1.0
- [2] TD S3-040xxx, “GBA\_U in Rel 6 Mes”, Axalto, Gemplus, SA3#34
- [3] TD S3-040xxx, “CR to TS 33.220, GBA: GBA\_U derivations: ”, Axalto, Gemplus, SA3#34