

6.3.3 Authentication and key agreement

The purpose of this procedure is to authenticate the user and establish a new pair of cipher and integrity keys between the VLR/SGSN and the USIM. During the authentication, the USIM verifies the freshness of the authentication vector that is used.

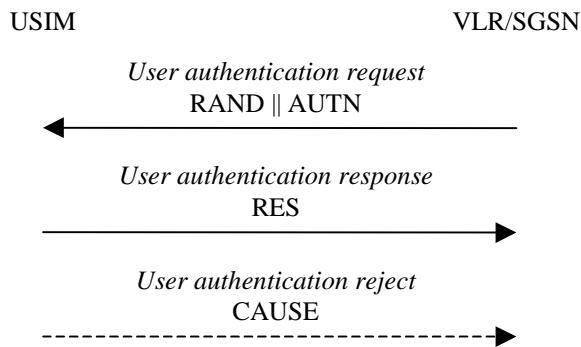


Figure 8: Authentication and key establishment

The VLR/SGSN invokes the procedure by selecting the next unused authentication vector from the ordered array of authentication vectors in the VLR/SGSN database. The VLR/SGSN sends to the USIM the random challenge $RAND$ and an authentication token for network authentication $AUTN$ from the selected authentication vector.

Upon receipt the user proceeds as shown in Figure 9.

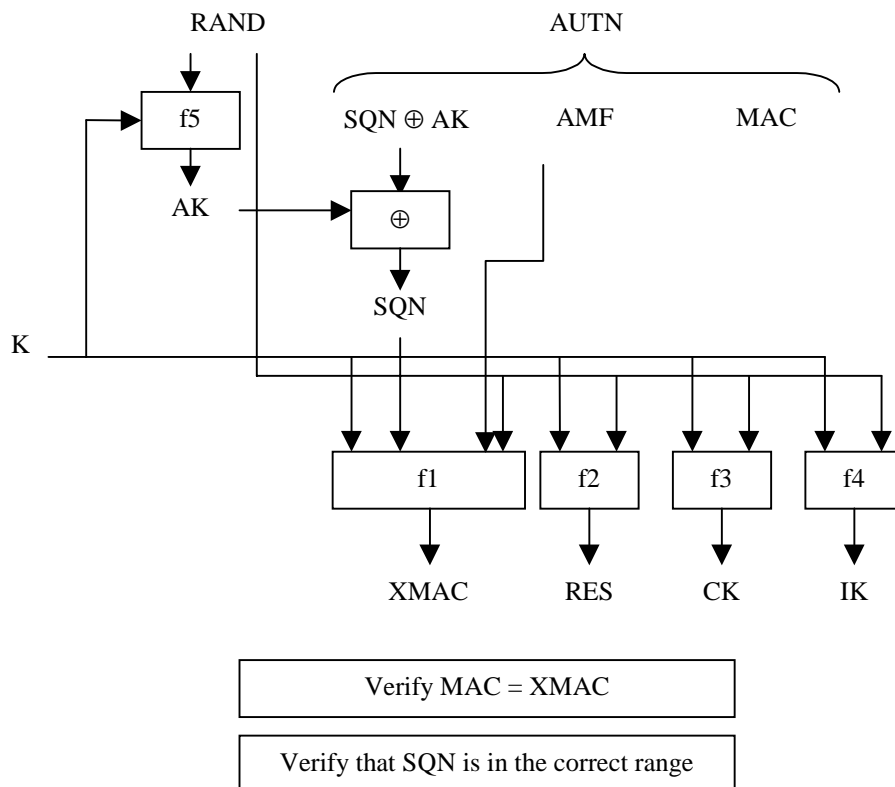


Figure 9: User authentication function in the USIM

Upon receipt of RAND and AUTN the USIM first computes the anonymity key $AK = f5_K(RAND)$ and retrieves the sequence number $SQN = (SQN \oplus AK) \oplus AK$.

Next the USIM computes $XMAC = f1_K(SQN \parallel RAND \parallel AMF)$ and compares this with MAC which is included in AUTN. If they are different, the user sends *user authentication reject* back to the VLR/SGSN with an indication of the cause and the user abandons the procedure. In this case, VLR/SGSN shall initiate an Authentication Failure Report procedure towards the HLR as specified in section 6.3.6. VLR/SGSN may also decide to initiate a new identification and authentication procedure towards the user.

Next the USIM verifies that the received sequence number SQN is in the correct range.

If the USIM considers the sequence number to be not in the correct range, it sends *synchronisation failure* back to the VLR/SGSN including an appropriate parameter, and abandons the procedure.

The synchronisation failure message contains the parameter AUTS. It is $AUTS = Conc(SQN_{MS}) \parallel MAC-S$. $Conc(SQN_{MS}) = SQN_{MS} \oplus f5_K(MAC-S \parallel 0...0)$ is the concealed value of the counter SEQ_{MS} in the MS, and $MAC-S = f1^*_K(SEQ_{MS} \parallel RAND \parallel AMF)$ where RAND is the random value received in the current user authentication request. $f1^*$ is a message authentication code (MAC) function with the property that no valuable information can be inferred from the function values of $f1^*$ about those of $f1, \dots, f5$ and vice versa.

The AMF used to calculate MAC-S assumes a dummy value of all zeros so that it does not need to be transmitted in the clear in the re-synch message.

The construction of the parameter AUTS is shown in the following Figure 10:

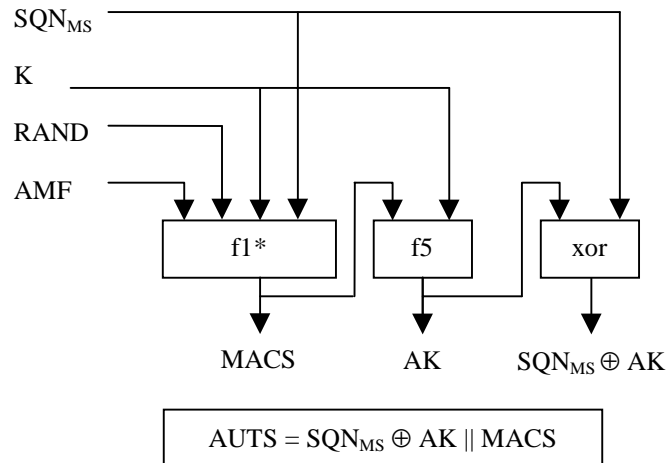


Figure 10: Construction of the parameter AUTS

If the sequence number is considered to be in the correct range however, the USIM computes $RES = f2_K(RAND)$ and includes this parameter in a *user authentication response* back to the VLR/SGSN. Finally the USIM computes the cipher key $CK = f3_K(RAND)$ and the integrity key $IK = f4_K(RAND)$. Note that if this is more efficient, RES, CK and IK could also be computed earlier at any time after receiving RAND. If the USIM also supports conversion function c3, it shall derive the GSM cipher key Kc from the UMTS cipher/integrity keys CK and IK. UMTS keys are sent to the MS along with the derived GSM key for UMTS-GSM interoperability purposes. USIM shall store original CK, IK until the next successful execution of AKA.

Upon receipt of *user authentication response* the VLR/SGSN compares RES with the expected response XRES from the selected authentication vector. If XRES equals RES then the authentication of the user has passed. The VLR/SGSN also selects the appropriate cipher key CK and integrity key IK from the selected authentication vector. If XRES and RES are different, VLR/SGSN shall initiate an Authentication Failure

Report procedure towards the HLR as specified in section 6.3.6. VLR/SGSN may also decide to initiate a new identification and authentication procedure towards the user.

Conditions on the use of authentication information by the VLR/SGSN: The VLR/SGSN shall use a UMTS authentication vector (i.e. a ~~quintuplet~~quintet) only once and, hence, shall send out each user authentication request *RAND // AUTN* only once no matter whether the authentication attempt was successful or not. A consequence is that UMTS authentication vectors (~~quintuplets~~quintets) cannot be reused.

6.8.1.2 R99+ HLR/AuC

Upon receipt of an *authentication data request* from a R99+ VLR/SGSN for a UMTS subscriber, a R99+ HLR/AuC shall send ~~quintuplets~~quintets, generated as specified in 6.3.

Upon receipt of an *authentication data request* from a R98- VLR/SGSN for a UMTS subscriber, a R99+ HLR/AuC shall send triplets, derived from ~~quintuplets~~quintets using the following conversion functions:

- a) $c1: RAND_{[GSM]} = RAND$
- b) $c2: SRES_{[GSM]} = XRES_1 [xor XRES_2 [xor XRES_3 [xor XRES_4]]]$
- c) $c3: Kc_{[GSM]} = CK_1 xor CK_2 xor IK_1 xor IK_2$

whereby $XRES_i$ are all 32 bit long and $XRES = XRES_1 [|| XRES_2 [|| XRES_3 [|| XRES_4]]]$ dependent on the length of $XRES$, and CK_i and IK_i are both 64 bits long and $CK = CK_1 || CK_2$ and $IK = IK_1 || IK_2$.

6.8.1.3 R99+ VLR/SGSN

The AKA procedure will depend on the terminal capabilities, as follows:

- UMTS subscriber with R99+ ME

When the user has R99+ ME, UMTS AKA shall be performed using a ~~quintuplet~~quintet that is either:

- a) retrieved from the local database,
- b) provided by the HLR/AuC, or
- c) provided by the previously visited R99+ VLR/SGSN.

Note: Originally all ~~quintuplets~~quintets are provided by the HLR/AuC.

UMTS AKA results in the establishment of a UMTS security context; the UMTS cipher/integrity keys CK and IK and the key set identifier KSI are stored in the VLR/SGSN.

When the user is attached to a UTRAN, the UMTS cipher/integrity keys are sent to the RNC, where the cipher/integrity algorithms are allocated.

When the user is attached to a GSM BSS, UMTS AKA is followed by the derivation of the GSM cipher key from the UMTS cipher/integrity keys. When the user receives service from an MSC/VLR, the derived cipher key Kc is then sent to the BSC (and forwarded to the BTS). When the user receives service from an SGSN, the derived cipher key Kc is applied in the SGSN itself.

UMTS authentication and key freshness is always provided to UMTS subscribers with R99+ ME independently of the radio access network.

- UMTS subscriber with R98- ME

When the user has R98- ME, the R99+ VLR/SGSN shall perform GSM AKA using a triplet that is either

- a) derived by means of the conversion functions $c2$ and $c3$ in the R99+ VLR/SGSN from a ~~quintuplet~~quintet that is:
 - i) retrieved from the local database,
 - ii) provided by the HLR/AuC, or

- iii) provided by the previously visited R99+ VLR/SGSN, or
- b) provided as a triplet by the previously visited MSC/VLR or SGSN.

NOTE: R99+ VLR/SGSN will always provide ~~triplets~~ quintets for UMTS subscribers.

NOTE: For a UMTS subscriber, all triplets are derived from ~~triplets~~ quintets, be it in the HLR/AuC or in an VLR/SGSN.

GSM AKA results in the establishment of a GSM security context; the GSM cipher key Kc and the cipher key sequence number CKSN are stored in the VLR/SGSN.

In this case the user is attached to a GSM BSS. When the user receives service from an MSC/VLR, the GSM cipher key is sent to the BSC (and forwarded to the BTS). When the user receives service from an SGSN, the derived cipher key Kc is applied in the SGSN itself.

UMTS authentication and key freshness cannot be provided to UMTS subscriber with R98- ME.

6.8.3 Distribution and use of authentication data between VLRs/SGSNs

The distribution of authentication data (unused authentication vectors and/or current security context data) between R99+ VLRs/SGSNs of the same service network domain is performed according to chapter 6.3.4. The following four cases are distinguished related to the distribution of authentication data between VLRs/SGSNs (of the same or different releases). Conditions for the distribution of such data and for its use when received at VLRn/SGSNn are indicated for each case:

a) R99+ VLR/SGSN to R99+ VLR/SGSN

UMTS and GSM authentication vectors can be distributed between R99+ VLRs/SGSNs. Note that originally all authentication vectors (~~quintuplets~~quintets for UMTS subscribers and triplets for GSM subscribers) are provided by the HLR/AuC.

Current security context data can be distributed between R99+ VLRs/SGSNs. VLRn/SGSNn shall not use current security context data received from VLRO/SGSN0 to authenticate the subscriber using local authentication in the following cases:

- i) Security context to be established at VLRn/SGSNn requires a different set of keys than the one currently in use at VLRO/SGSN0. This change of security context is caused by a change of ME release ($R'99 \text{ ME} \leftrightarrow R'98 \text{ ME}$) when the user registers at VLRn/SGSNn.
- ii) Authentication data from VLRO includes Kc+CKSN but no unused AVs and the subscriber has a R'99 ME (under GSM BSS or UTRAN). In this situation, VLRn have no indication of whether the subscriber is GSM or UMTS and it is not able to decide whether Kc received can be used (in case the subscriber were a GSM subscriber).

In these two cases, received current security context data shall be discarded and a new AKA procedure shall be performed.

b) R98- VLR/SGSN to R98- VLR/SGSN

Only triplets can be distributed between R98- VLRs/SGSNs. Note that originally for GSM subscribers, triplets are generated by HLR/AuC and for UMTS subscribers, they are derived from UMTS authentication vectors by R99+ HLR/AuC. UMTS AKA is not supported and only GSM security context can be established by a R98- VLR/SGSN.

R98- VLRs are not prepared to distribute current security context data.

Since only GSM security context can be established under R98- SGSNs, security context data can be distributed and used between R98- SGSNs.

c) R99+ VLR/SGSN to R98- VLR/SGSN

R99+ VLR/SGSN can distribute to a new R98- VLR/SGSN triplets originally provided by HLR/AuC for GSM subscribers or can derive triplets from stored ~~quintuplets~~quintets originally provided by R99+ HLR/AuC for UMTS subscribers. Note that R98- VLR/SGSN can only establish GSM security context.

R99+ VLRs shall not distribute current security context data to R98- VLRs.

Since R98- SGSNs are only prepared to handle GSM security context data, R99+ SGSNs shall only distribute GSM security context data (Kc, CKSN) to R98- SGSNs.

d) R98- VLR/SGSN to R99+ VLR/SGSN

In order to not establish a GSM security context for a UMTS subscriber, triplets provided by a R98-VLR/SGSN can only be used by a R99+ VLR/SGSN to establish a GSM security context under GSM-BSS with a R98- ME.

In all other cases, R99+ VLR/SGSN shall request fresh AVs (either triplets or ~~quintuplets~~quintets) to HE. In the event, the R99+ VLR/SGSN receives ~~quintuplets~~quintets, it shall discard the triplets provided by the R98- VLR/SGSN.

R98- VLRs are not prepared to distribute current security context data.

R98- SGSNs can distribute GSM security context data only. The use of this information at R99+ SGSNs shall be performed according to the conditions stated in a).