

Agenda item:

Source: Ericsson

Title: Text proposal for RACH channelization code allocation

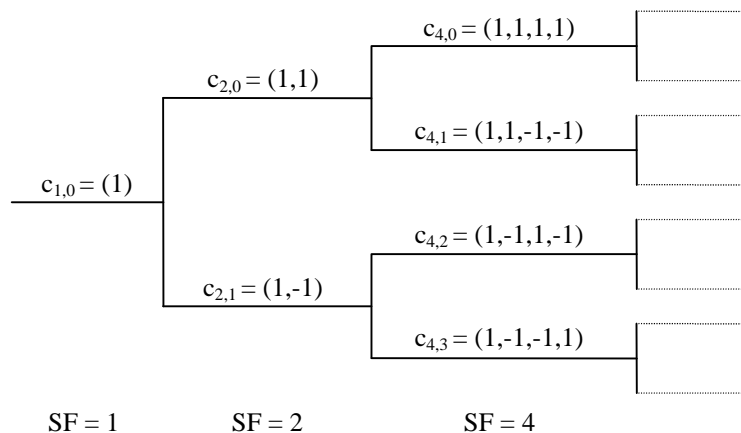
Document for: Decision

1 Introduction

The current section describing the channelization code assignment for the RACH message part is not 100% clear, and requires some further detailed specification. This paper proposes such text.

2 Text proposal for TS 25.213 V2.1.2

The text proposal in the following assumes the numbering of channelization codes proposed by Ericsson in TSGR1#6(99)845, i.e. the following numbering of the codes in the code tree:



4.3.3.4 Channelization codes for the message part

The preamble signature $s, 1 \leq s \leq 16$, in the preamble specifies points to one of the 16 nodes in the code-tree that corresponds to channelization codes of length 16, as shown in Figure 14. The sub-tree below the specified node is used for spreading of the message part. The control part(Q-branch) is spread with the channelization code $C_{ch,c}$ of spreading factor 256 in the lowest branch of the sub-tree, i.e. $C_{ch,c} = c_{256,m}$ where $m = 16(s - 1) + 15$. The spread control part is mapped to the Q-branch, similar to the DPCCH for dedicated channels.

-The data part(I-branch) can use any of the channelization codes from spreading factor 32 to 256 in the upper-most branch of the sub-tree. To be exact, the data part is spread by channelization code $C_{ch,d}$, where $C_{ch,d} = c_{SF,m}$ and SF is the spreading factor used for the data part and $m = SF \times (s - 1) / 16$.

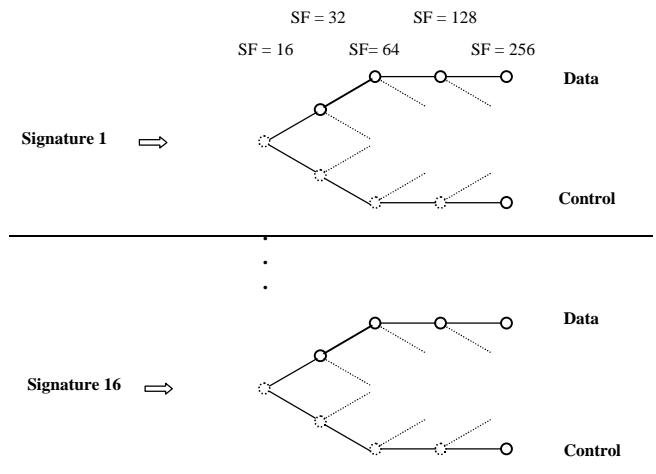


Figure 11. Channelization codes for the random access message part.