

PSEUDO-CHANGE REQUEST

⌘ **33.878 CR CRNum** ⌘ rev **-** ⌘ Current version: **1.0.0** ⌘

For **HELP** on using this form, see bottom of this page or look at the pop-up text over the ⌘ symbols.

Proposed change affects: | UICC apps ME Radio Access Network Core Network

Title:	⌘ Architecture of early IMS Security		
Source:	⌘ ZTE Corporation		
Work item code:	⌘ Early IMS	Date:	⌘ 27/12/2004
Category:	⌘ B	Release:	⌘ Rel-6
	Use <u>one</u> of the following categories: F (correction) A (corresponds to a correction in an earlier release) B (addition of feature), C (functional modification of feature) D (editorial modification) Detailed explanations of the above categories can be found in 3GPP TR 21.900 .		Use <u>one</u> of the following releases: <i>Ph2</i> (GSM Phase 2) <i>R96</i> (Release 1996) <i>R97</i> (Release 1997) <i>R98</i> (Release 1998) <i>R99</i> (Release 1999) <i>Rel-4</i> (Release 4) <i>Rel-5</i> (Release 5) <i>Rel-6</i> (Release 6) <i>Rel-7</i> (Release 7)

Reason for change:	⌘ There is not an architecture to reflect security mechanism in TR 33.878.
Summary of change:	⌘ Adding an architecture of early IMS security
Consequences if not approved:	⌘ An architecture of early IMS security is not included in current draft specification.

Clauses affected:	⌘ 6.1										
Other specs affected:	<table border="1" style="display: inline-table; border-collapse: collapse;"> <tr> <td style="width: 20px; text-align: center;">Y</td> <td style="width: 20px; text-align: center;">N</td> </tr> <tr> <td style="text-align: center;"><input type="checkbox"/></td> <td style="text-align: center;"><input checked="" type="checkbox"/></td> </tr> <tr> <td style="text-align: center;"><input type="checkbox"/></td> <td style="text-align: center;"><input checked="" type="checkbox"/></td> </tr> <tr> <td style="text-align: center;"><input type="checkbox"/></td> <td style="text-align: center;"><input checked="" type="checkbox"/></td> </tr> </table>	Y	N	<input type="checkbox"/>	<input checked="" type="checkbox"/>	<input type="checkbox"/>	<input checked="" type="checkbox"/>	<input type="checkbox"/>	<input checked="" type="checkbox"/>	Other core specifications ⌘ Test specifications O&M Specifications	⌘
Y	N										
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Other comments:	⌘										

*** BEGIN SET OF CHANGES ***

6 Specification

6.1 Overview

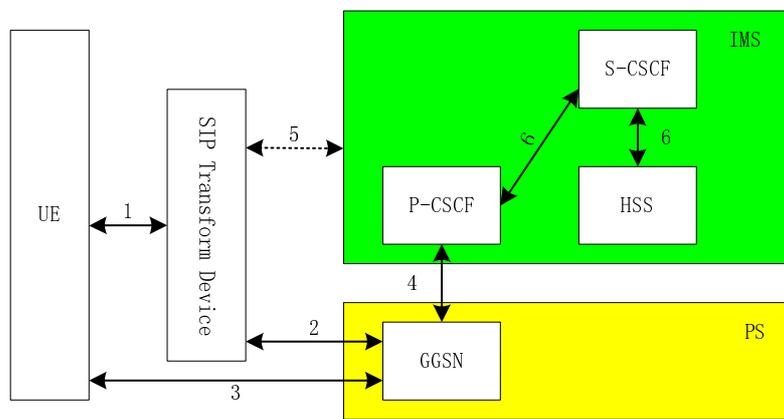


figure X: architecture of early IMS security

There are six different security associations and different needs for security protection in architecture of early IMS security and they are numbered 1, 2, 3, 4, 5 and 6 in figure 1 where:

1. Provides the security protection of message exchanged between UE and SIP transform device. The security mechanism may be out of early IMS specification.
2. Provides to check the binding between the IP address that the GGSN allocated the UE in the PDP context activation and the source IP address in subsequent packets.
3. Provides security in GPRS bearer level
4. Provides security between SIP capable nodes. This security association is covered by TS 33.210.
5. Provides to check the binding between the IP address on the bearer level, and the user identities.
6. Provides security within IMS. This security association is covered by TS 33.210.

The early IMS security solution works by creating a secure binding in the HSS between the public/private user identity (SIP-level identity) and the IP address currently allocated to the user at the GPRS level (bearer/network level identity). Therefore, IMS level signaling, and especially the IMS identities claimed by a user, can be connected securely to the PS domain bearer level security context.

A SIP Transform Device can transform other message into SIP message so that the UE without SIP module can access the IMS service. The SIP Transform Device should support the security mechanism of early IMS. It should support GGSN to perform source IP address spoofing check.

The GGSN, terminates each user's PDP context and has assurance that the IMSI used within this PDP context is authenticated. The GGSN shall provide the user's IP address, IMSI and MSISDN to a RADIUS server in the HSS over the Gi interface when a PDP context is activated towards the IMS system. The HSS has a binding between the IMSI and/or MSISDN and the IMPI and IMPU(s), and is therefore able to store the currently assigned IP address from the GGSN against the user's IMPI and/or IMPU(s). The precise way of the handling of these identities in the HSS is outside the scope of standardization. The GGSN informs the HSS when the PDP context is deactivated/modified so that the stored IP address can be updated in the HSS. When the S-CSCF receives a SIP registration request or any subsequent requests for a given IMPU, it checks that the IP address in the SIP header (verified by the network) matches the IP address that was stored against that subscriber's IMPU in the HSS.

The mechanism assumes that the GGSN does not allow a UE to successfully transmit an IP packet with a source IP address that is different to the one assigned during PDP context activation. In other words, the GGSN must prevent

"source IP spoofing". The mechanism also assumes that the P-CSCF checks that the source IP address in the SIP header is the same as the source IP address in the IP header received from the UE (the assumption here, as well as for the full security solution, is that no NAT is present between the GGSN and the P-CSCF).

The mechanism prevents an attacker from using his own IP address in the IP header but spoofing someone else's IMS identity or IP address in the SIP header, so that he pays for GPRS level charges, but not for IMS level charges. The mechanism also prevents an attacker spoofing the address in the IP header so that he does not pay for GPRS charges. It therefore counters the threat scenarios given in clause 5 above.

The mechanism assumes that only one contact IP address is associated with one IMPI. Furthermore, the mechanism supports the case that there may be several IMPUs associated with one IMPI, but one IMPU is associated with only one IMPI.

In early IMS the IMS user authentication is performed by linking the IMS registration (based on an IMPI) to a PDP context (based on an authenticated IMSI). The mechanism here assumes that there is a one-to-one relationship between the IMSI for bearer access and the IMPI for IMS access.

For the purposes of this present document, an APN, which is used for IMS services, is called an IMS APN. An IMS APN may be also used for non-IMS services. The mechanism described in this present document further adds a restriction that there is only one APN for accessing IMS for a PLMN and that all active PDP contexts, for a single UE, associated with that IMS APN use the same IP address at any given time.

In the following we use the terms P-CSCF and S-CSCF in a general sense to refer to components of an early IMS system. We note however that early IMS solutions may not have the same functionality split between SIP entities as defined in TS 23.228 [3]. Therefore, the requirements imposed on the SIP/IP core are specified in such a way that they are independent of the functionality split between SIP entities as far as possible. While the exact functionality split of the SIP/IP core may be left open, it is important that any changes to the Cx interface towards the HSS and changes to the interface towards the UE are standardised for vendor interoperability reasons.

*** END SET OF CHANGES ***