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CR-Form-v5

**CHANGE REQUEST**⌘ **33.203** CR ⌘ rev - ⌘ Current version: **5.2.0** ⌘For **HELP** on using this form, see bottom of this page or look at the pop-up text over the ⌘ symbols.Proposed change affects: ⌘ (U)SIM  ME/UE  Radio Access Network  Core Network **Title:** ⌘ Update of SA handling procedures**Source:** ⌘ Hutchison 3G UK**Work item code:** ⌘ **Date:** ⌘ 4/7/2002**Category:** ⌘ **F**Use one of the following categories:**F** (correction)**A** (corresponds to a correction in an earlier release)**B** (addition of feature),**C** (functional modification of feature)**D** (editorial modification)Detailed explanations of the above categories can be found in 3GPP [TR 21.900](#).**Release:** ⌘ Rel-5Use one of the following releases:

2 (GSM Phase 2)

R96 (Release 1996)

R97 (Release 1997)

R98 (Release 1998)

R99 (Release 1999)

REL-4 (Release 4)

REL-5 (Release 5)

**Reason for change:** ⌘ Current security association (SA) handling procedures do not cover all the possible cases that can occur**Summary of change:** ⌘ Update the way the P-CSCF handles security associations to deal with some cases that are not already covered. Also describes the behaviour of the UE and P-CSCF in isolation of each other.**Consequences if not approved:** ⌘ Some behaviour of the P-CSCF is not described, which means that different P-CSCF may take different action possibly causing the UE to become unreachable.**Clauses affected:** ⌘ 7.4**Other specs affected:** ⌘  Other core specifications ⌘  Test specifications  O&M Specifications**Other comments:** ⌘

## 7.4 Authenticated re-registration

Every registration that includes a user authentication attempt produces new security associations. If the authentication is successful, then these new security associations shall replace the previous ones. This clause describes how the UE and P-CSCF handle this replacement and which SAs to apply to which message.

~~If the registration is a re-registration, a pair of security associations between UE has an and P-CSCF is already active security association, then it should use this to protect REGISTER message. If the S-CSCF is notified by the P-CSCF that the REGISTER message from the UE was integrity-protected it may decide not to authentication the user by means of the AKA protocol. However, the UE may send unprotected REGISTER messages at any time. In this case, the S-CSCF shall authenticate the user by means of the AKA protocol. This means the UE shall send an unprotected REGISTER, if it requires an authentication to refresh its security associations, for example it believes there is a problem with the security associations it holds or it only has security association that will expire soon and a registration that will expire later. The authenticated re-registration shall initially utilize the existing SA. This is the normal case. However, in the event the UE originates the (SM1) Register message using no protection, the P-CSCF shall still accept it and forward the request to the S-CSCF, indicating that the register message was not integrity protected. This shall trigger the S-CSCF to challenge the subscriber with the execution of a new IMS AKA authentication procedure as described in clause 6.1.1.~~

~~[Editors Note: The exact mechanism for changing SAs is currently under investigation.]~~

~~Before SM7 is sent by the UE, both peers shall replace the existing SA by the new SA negotiated during these first two messages:~~

~~[Editors Note: The exact mechanism when to change SA1 to SA2 under certain error conditions is FFS.]~~

Security associations may be unidirectional or bi-directional. This section assumes that security associations are unidirectional, as this is the general case. For IP layer SAs, it is possible that there needs to be parallel SAs for each available transport protocol. Whenever a user is registered there are **current SAs** at both the P-CSCF and the UE. At the UE, there may also be either **registration SAs** or **inbound old SAs**. Whilst at the P-CSCF, there may also be **registration SAs** and/or a **valid SAs**. They are denoted as follows:

<u>SA_in_cur</u>	current inbound SA
<u>SA_out_cur</u>	current outbound SA
<u>SA_in_reg</u>	registration inbound SA
<u>SA_out_reg</u>	registration outbound SA
<u>SA_in_old</u>	old inbound SA (in UE only)
<u>SA_in_val</u>	valid inbound SA (in P-CSCF only)
<u>SA_out_val</u>	valid outbound SA (in P-CSCF only)

This notation has local significance only. That means that SA\_in\_cur at the UE is not always the same as SA\_out\_cur at the P-CSCF and similarly for other SAs.

For IP layer SAs, the lifetime mentioned in the following clauses is the lifetime held at the application layer. Furthermore deleting an SA means deleting the SA from both the application and network layer. If parallel SAs are needed for more than more transport, the SA management procedures in the following clauses need to be applied for each parallel set of SAs. The message numbers, e.g. SM1, used in the following clauses relate to the message flow given in 6.1.1.

### 7.4.1 ~~7.4.1~~ Management of security associations in the UE ~~Handling of security associations in authenticated re-registrations (successful case)~~

The UE shall be involved in only one registration procedure at a time. Upon starting a new registration procedure, any existing registration SAs shall be deleted. The UE shall delete any SA whose lifetime is exceeded. If the wrong SA is used to protect any message, the message shall be discarded.

A successful authentication proceeds in the following steps:

- The UE sends the SM1 message to register with the IMS. If protected it should be integrity-protected using SA\_out\_cur.

- The UE receives an authentication challenge in a message (SM6) from the P-CSCF. This message shall be integrity-protected using SA\_in\_cur if SM1 was integrity-protected.
- If this message SM6 can be successfully processed by the UE, the UE deletes SA\_out\_old if it exists and creates the new SAs, SA\_in\_reg and SA\_out\_reg, which are derived according to section 7.1. The lifetime of the registration SAs should be set to allow enough time to complete the registration procedure. The UE then sends its response (SM7) to the P-CSCF, which shall be protected with SA\_out\_reg.
- The UE receives an authentication successful message (SM12) from the P-CSCF, which shall be protected using SA\_in\_reg.
- After the successful processing of this message by the UE, the registration is complete. The UE sets the lifetime of the registration SAs using the registration timer in the message. SA\_in\_cur becomes the new SA\_in\_old if it exists and SM1 was protected. SA\_out\_reg becomes the new SA\_out\_cur and SA\_in\_reg becomes the new SA\_in\_cur.

A failure in the authentication means the UE shall delete SA\_in\_reg and SA\_out\_reg. If SM1 was protected, the UE shall protect all outbound failure messages in the authentication with SA\_out\_cur and ensure that SA\_in\_cur was applied to protect all inbound failure messages in the authentication. If the SM1 was not protected, then no protection shall be applied to the failure messages.

When a SIP message protected with SA\_in\_cur is successfully received from the P-CSCF, the UE shall delete SA\_in\_old if it exists.

For messages outside an authentication, the UE shall use SA\_out\_cur to protect all outbound traffic and ensure that all inbound traffic is protected with either SA\_in\_cur or SA\_in\_old.

~~Before re-registration begins the following SAs exist:~~

- ~~— SA1 from UE to P-CSCF;~~
- ~~— SA2 from P-CSCF to UE.~~

~~The re-registration then is as follows:~~

- ~~1) The UE sends SM1 to re-register with the IMS, using the existing SA1 to the P-CSCF. As in the case of a new registration, a list of parameters to be negotiated in a security association set-up is included.~~

~~{Editors Note: It is FFS if the SA1 shall be used for SM1 or not}~~

- ~~2) The P-CSCF waits for the response SM4 from the S-CSCF and then sends SM6 to the UE, using SA2. As in the case of a new registration, the parameters selected for the new security associations are included. The P-CSCF then creates two new security associations, in parallel to the existing ones, in its database:~~

- ~~— SA11 from UE to P-CSCF;~~
- ~~— SA12 from P-CSCF to UE.~~

- ~~3) If SM6 could be successfully processed by the UE, the UE also creates the new SAs SA11 and SA12 in its database. The UE then sends SM7 to the P-CSCF. As in the case of a new registration, the authentication response and the list of parameters repeated from message 1 are included. SM7 is protected with the new SA11.~~

- ~~4) The P-CSCF waits for the response SM10 from the S-CSCF and then sends SM12 to the UE, using the new SA12.~~

- ~~5) After the reception of SM12 by the UE, the re-registration is complete.~~

~~The UE now uses the new SAs for all subsequent messages. The old (outbound) SA1 is deleted. The old (inbound) SA2 must be kept until a further SIP message protected with the new inbound SA12 is successfully received from the P-CSCF.~~

~~The P-CSCF keeps all four SAs with the UE active until a further SIP message protected with the new inbound SA11 is successfully received from the UE. In the meantime, the P-CSCF continues to use the old SA2 for outbound traffic to the UE.~~

## 7.4.2 Management of security associations in the P-CSCF ~~Error cases related to authenticated re-registration~~

~~Whenever an expected message is not received after a time-out the receiving entity considers the registration to have failed. The receiving entity then deletes any new security associations it may have established and continues to use the old ones if they have not yet expired.~~

~~If the registration protocol goes well up to the last message SM12, and SM12 is sent by the P-CSCF, but not received by the UE, then the UE has only the old SAs available (after the time-out), but the P-CSCF cannot know this. Therefore, the P-CSCF continues to use the old SA2 for outbound traffic to the UE, but keeps both, old and new SAs. The new SAs are deleted when a message is received from the UE which is protected with the old SA, or if a REGISTER message is received on the port where the P-CSCF accepts specific unprotected messages.~~

The P-CSCF shall delete any SA whose lifetime is exceeded. If the current SAs are deleted and there exist valid SAs, then the P-CSCF makes the SA\_out\_val the new SA\_out\_cur and SA\_in\_val the new SA\_in\_cur, and removes the valid SAs. If the wrong SA is used to protect any message, the message shall be discarded.

The P-CSCF associates the IMPI and IMPU given in the registration procedure with the registration SAs created during that registration procedure. The P-CSCF associates the IMPI given in the registration procedure and all the successfully registered IMPUs related to that IMPI with current and valid SAs.

A successful authentication proceeds in the following steps:

- The P-CSCF receives the SM1 message. If it is protected, it should be integrity-protected using SA\_in\_cur or SA\_in\_val.
- The P-CSCF forwards the message containing the challenge (SM6) to the UE. This shall be integrity-protected using SA\_out\_cur, if SM1 was protected and unprotected otherwise.
- The P-CSCF then creates the new SAs, SA\_in\_reg and SA\_out\_reg, which are derived according to section 7.1. The expiry time of the registration SAs should be set to allow just enough time to complete the registration procedure. These SAs shall overwrite any previous registration SAs related to that IMPU and may overwrite any previous registration SAs related to that IMPI
- The P-CSCF receives the message carrying the response (SM7) from the UE. It shall be protected using SA\_in\_reg related to that IMPU.
- The P-CSCF forwards the successful registration message (SM12) to the UE, which shall be protected using SA\_out\_reg related to that IMPU. This completes the registration procedure for the P-CSCF. The P-CSCF sets the expiry time of the registration SAs equal to the registration timer in the message. If SM1 was protected and the current SAs exist, then SA\_out\_reg becomes SA\_out\_val, SA\_in\_reg becomes SA\_in\_val (overwriting any previous valid SAs) and the expiry times of SA\_in\_cur and SA\_out\_cur should be shortened to allow only enough time for a further authentication in case of lost messages. Otherwise SA\_out\_reg becomes SA\_out\_cur and SA\_in\_reg becomes SA\_in\_cur, and all valid and registration SAs are deleted.

A failure in the authentication means the P-CSCF shall delete SA\_in\_reg and SA\_out\_reg. If SM1 was protected, the P-CSCF shall protect all outbound failure messages in the authentication with SA\_out\_cur and ensure that SA\_in\_cur was applied to protect all inbound failure messages in the authentication. If the SM1 was not protected, then no protection shall be applied to the failure messages.

When the P-CSCF successfully receives a SIP message protected with SA\_in\_val from the UE, then SA\_in\_val and SA\_out\_val becomes the new SA\_in\_cur and SA\_out\_cur respectively, and there are no more valid SAs.

For messages outside an authentication, the P-CSCF shall use SA\_out\_cur to protect all outbound traffic and ensure that all inbound traffic is protected with either SA\_in\_cur or SA\_in\_val.