

# 3G TR 33.900 V1.2.0 (2000-01)

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*Technical Specification*

## **3rd Generation Partnership Project; Technical Specification Group SA WG3; A Guide to 3<sup>rd</sup> Generation Security (3G TR 33.900 version 1.2.0)**



The present document has been developed within the 3<sup>rd</sup> Generation Partnership Project (3GPP™) and may be further elaborated for the purposes of 3GPP.

The present document has not been subject to any approval process by the 3GPP Organisational Partners and shall not be implemented.

This Specification is provided for future development work within 3GPP only. The Organisational Partners accept no liability for any use of this Specification.

Specifications and reports for implementation of the 3GPP™ system should be obtained via the 3GPP Organisational Partners' Publications Offices.

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Reference

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Keywords

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## 2 Foreword

The 3GPP have produced this Technical Report.

The contents of the present document are subject to continuing work within the TSG and may change following formal TSG approval. Should the TSG modify the contents of this TS, it will be re-released by the TSG with an identifying change of release date and an increase in version number as follows:

Version 3.y.z

where:

- x the first digit:
  - 1 presented to TSG for information;
  - 2 presented to TSG for approval;
  - 3 Indicates TSG approved document under change control.
- y the second digit is incremented for all changes of substance, i.e. technical enhancements, corrections, updates, etc.
- z the third digit is incremented when editorial only changes have been incorporated in the specification;

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## 3 Introduction

This document is intended to offer security guidance to those involved in 3GPP systems. All specifications have to take into account the cost and feasibility of security features and functions. Inevitably, some compromise is necessary, and so it is important to realise where possible risks and threats may exist. The document describes those security issues that have been identified in the formulation of the standards.

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## 4 Scope

The present document gives a general description of the security architecture and features of 3<sup>rd</sup> Generation Security. It is intended to provide an overview of security, for detailed explanation and the actual standards the reader is referred to the appropriate standards.

It also serves the purpose of identifying the potential risks and threats that have been highlighted and require careful consideration when implementing a third generation mobile system.

The document attempts to identify whether the security features and mechanism provided in the latest version of the 3G security architecture specification {1} address the 2G security weaknesses.

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## 5 References

Reference to an ETS shall also be taken to refer to later versions published as an EN with the same number.

- [1] 3G TS 33.102, 3G Security; Security Architecture, version 3.0.0
- [2] 3G TS 33.120, 3G Security; Security Principles and Objectives, version 3.0.0
- [3] 3G TS 21.133, 3G Security; Security Threats and Requirements, version 3.0.01

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## 6 Definitions, symbols and abbreviations

### 6.1 Definitions

### 6.2 Symbols

### 6.3 Abbreviations

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## 7 A brief overview of 3GPP Security

3GPP security was based on GSM security, with the following important changes:

- A change was made to defeat the false base station attack. The security mechanisms include a sequence number that ensures that the mobile can identify the network.
- Key lengths were increased to allow for the possibility of stronger algorithms for encryption and integrity.
- Mechanisms were included to support security within and between networks.
- Security is based within the switch rather than the base station as in GSM. Therefore links are protected between the base station and switch.
- Integrity mechanisms for the terminal identity (IMEI) have been designed in from the start, rather than that introduced late into GSM.
- The authentication algorithm has not been defined, but guidance on choice will be given.
- When roaming between networks, such as between a GSM and 3GPP, only the level of protection supported by the smart card will apply. Therefore a GSM smart card will not be protected against the false base station attack when in a 3GPP network.

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## 8 Counteracting envisaged 3G attacks

Many of the security enhancements required to 2G systems are intended to counteract attacks which were not perceived to be feasible in 2G systems. This includes attacks that are, or are perceived to be, possible now or very soon because intruders have access to more computational capabilities, new equipment has become available, and the physical security of certain network elements is questioned.

In order to perform the attacks the intruder has to possess one or more of the following capabilities:

- **Eavesdropping.** This is the capability that the intruder eavesdrops signalling and data connections associated with other users. The required equipment is a *modified MS*.
- **Impersonation of a user.** This is the capability whereby the intruder sends signalling and/or user data to the network, in an attempt to make the network believe they originate from the target user. The required equipment is again a *modified MS*.
- **Impersonation of the network.** This is the capability whereby the intruder sends signalling and/or user data to the target user, in an attempt to make the target user believe they originate from a genuine network. The required equipment is modified *BS*.
- **Man-in-the-middle.** This is the capability whereby the intruder puts itself in between the target user and a genuine network and has the ability to eavesdrop, modify, delete, re-order, replay, and spoof signalling and user data messages exchanged between the two parties. The required equipment is modified *BS in conjunction with a modified MS*.
- **Compromising authentication vectors in the network.** The intruder possesses a *compromised authentication vector*, which may include challenge/response pairs,

cipher keys and integrity keys. This data may have been obtained by compromising network nodes or by intercepting signalling messages on network links.

The first capability is the easiest to achieve the following capabilities are gradually more complex and require more investment by the attacker. Therefore, in general, an intruder having a certain capability is assumed also to have the capabilities positioned above that capability in the list. The first two capabilities were acknowledged in the design of 2G systems. 3G security however should thwart all five types of attacks.

In the following we consider several attacks to 3G systems which may not have been fully addressed in 2G systems and attempt to identify whether the security features and mechanisms provided in the latest version of the 3G security architecture specification counteracts each of these attacks.

## 8.1 Denial of service

We distinguish between the following denial of service attacks:

### 8.1.1 User de-registration request spoofing

#### *Description:*

An attack that requires a *modified MS* and exploits the weakness that the network cannot authenticate the messages it receives over the radio interface. The intruder spoofs a de-registration request (IMSI detach) to the network. The network de-registers the user from the visited location area and instructs the HLR to do the same. The user is subsequently unreachable for mobile terminated services.

*Does 3G security architecture counteract the attack: Yes*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the de-registration request allows the serving network to verify that the de-registration request is legitimate.

### 8.1.2 Location update request spoofing

#### *Description:*

An attack that requires a *modified MS* and exploits the weakness that the network cannot authenticate the messages it receives over the radio interface. Instead of the de-registration request, the user spoofs a location update request in a different location area from the one in which the user is roaming. The network registers in the new location area and the target user will be paged in that new area. The user is subsequently unreachable for mobile terminated services.

*Does 3G security architecture counteract the attack: Yes*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the location update request allows the serving network to verify that the location update request is legitimate.

### 8.1.3 Camping on a false BS

#### *Description:*



An attack that requires a *modified BS* and exploits the weakness that a user can be enticed to camp on a false base station. Once the target user camps on the radio channels of a false base station, the target user is out of reach of the paging signals of the serving network in which he is registered.

*Does 3G security architecture counteract the attack: No*

The security architecture does not counteract this attack. However, the denial of service in this case only persists for as long as the attacker is active unlike the above attacks which persist beyond the moment where intervention by the attacker stops. These attacks are comparable to radio jamming which is very difficult to counteract effectively in any radio system.

#### 8.1.4 Camping on a false BS/MS

*Description:*

An attack that requires a *modified BS/MS* and exploits the weakness that a user can be enticed to camp on a false base station. A false BS/MS can act as a repeater for some time and can relay some requests in between the network and the target user, but subsequently modify or ignore certain service requests and/or paging messages related to the target user.

*Does 3G security architecture counteract the attack: No*

The security architecture does not prevent a false BS/MS relaying messages between the network and the target user, neither does it prevent the false BS/MS ignoring certain service requests and/or paging requests. Integrity protection of critical message may however help to prevent some denial of service attacks, which are induced by modifying certain messages. Again, the denial of service in this case only persists for as long as the attacker is active unlike the above attacks, which persist beyond the moment where intervention by the attacker stops. These attacks are comparable to radio jamming which is very difficult to counteract effectively in any radio system.

## 8.2 Identity catching

We identify the following types of attacks against the user identity confidentiality:

### 8.2.1 Passive identity catching

*Description:*

A passive attack that requires a *modified MS* and exploits the weakness that the network may sometimes request the user to send its identity in cleartext.

*Does 3G security architecture counteract the attack: Yes*

The identity confidentiality mechanism counteracts this attack. The use of temporary identities allocated by the serving network makes passive eavesdropping inefficient since the user must wait for a new registration or a mismatch in the serving network database before he can capture the user's permanent identity in plaintext. The inefficiency of this attack given the likely rewards to the attacker would make this scenario unlikely. (Note however that the permanent identity may be protected in the event of new registrations or serving network database failure in order to guard against more efficient active attacks.)

## 8.2.2 Active identity catching

### *Description:*

An active attack that requires a *modified BS* and exploits the weakness that the network may request the MS to send its permanent user identity in cleartext. An intruder entices the target user to camp on its false BS and subsequently requests the target user to send its permanent user identity in cleartext perhaps by forcing a new registration or by claiming a temporary identity mismatch due to database failure.

*Does 3G security architecture counteract the attack: Yes*

The identity confidentiality mechanism counteracts this attack by using an encryption key shared by a group of users to protect the user identity in the event of new registrations or temporary identity database failure in the serving network. Note however that the size of the groups should be chosen carefully: too small and the group identify may compromise the user identity itself; too large and the group encryption key might be vulnerable to attack.

## 8.3 Impersonation of the network

We identify the following attacks with the objective of impersonating a genuine network. The ultimate aim of such attacks is usually to eavesdrop on user data (see section 2.4), or to send the user information that he subsequently believes to originate from a genuine network or user with whom he is connected through that network.

### 8.3.1 Impersonation of the network by suppressing encryption between the target user and the intruder

#### *Description:*

An attack that requires a *modified BS* and that exploits the weakness that the MS cannot authenticate messages received over the radio interface. The target user is enticed to camp on the false BS. When the intruder or the target user initiates a service, the intruder does not enable encryption by spoofing the cipher mode command. The intruder maintains the call as long as it is required or as long as his attack remains undetected.

*Does 3G security architecture counteract the attack: Yes*

A mandatory cipher mode command with message authentication and replay inhibition allows the mobile to verify that encryption has not been suppressed by an attacker.

### 8.3.2 Impersonation of the network by suppressing encryption between the target user and the true network

#### *Description:*

An attack that requires a *modified BS/MS* and that exploits the weakness that the network cannot authenticate messages received over the radio interface. The target user is enticed to camp on the false BS/MS. When a call is set-up the false BS/MS modifies the ciphering capabilities of the MS to make it appear to the network that a genuine incompatibility exists between the network and the mobile station. The network may then decide to establish an un-enciphered connection. After the decision not to cipher has been taken,

the intruder cuts the connection with the network and impersonates the network to the target user.

*Does 3G security architecture counteract the attack: Yes*

A mobile station classmark with message authentication and replay inhibition allows the network to verify that encryption has not been suppressed by an attacker.

### 8.3.3 Impersonation of the network by forcing the use of a compromised cipher key

*Description:*

An attack that requires a *modified BS* and the possession by the intruder of a *compromised authentication vector* and thus exploits the weakness that the user has no control upon the cipher key. The target user is enticed to camp on the false BS/MS. When a call is set-up the false BS/MS forces the use of a compromised cipher key on the mobile user. The intruder maintains the call as long as it is required or as long as his attack remains undetected.

*Does 3G security architecture counteract the attack: Yes*

The presence of a sequence number in the challenge allows the USIM to verify the freshness of the cipher key to help guard against forced re-use of a compromised authentication vector. However, the architecture does not protect against force use of compromised authentication vectors which have not yet been used to authenticate the USIM. Thus, the network is still vulnerable to attacks using compromised authentication vectors which have been intercepted between generation in the authentication centre and use or destruction in the serving network.

The user must trust the SN (transitively via the HE) to handle authentication vectors securely. For instance, an attacker with a false BS may work in collusion with an SN to intercept unused authentication vectors, or the SN may expose itself to undue risks because it stockpiles large numbers of authentication vectors before they need to be used.

## 8.4 Eavesdropping on user data

We identify the following attacks with the objective of eavesdropping on user data which is transmitted through the genuine network to the intended recipient.

### 8.4.1 Eavesdropping on user data by suppressing encryption between the target user and the intruder

*Description:*

An attack that requires a *modified BS/MS* and that exploits the weakness that the MS cannot authenticate messages received over the radio interface. The target user is enticed to camp on the false BS. When the target user or the intruder initiates a call the network does not enable encryption by spoofing the cipher mode command. The attacker however sets up his own connection with the genuine network using his own subscription. The attacker may then subsequently eavesdrop on the transmitted user data.

*Does 3G security architecture counteract the attack: Yes*

A mandatory cipher mode command with message authentication and replay inhibition allows the mobile to verify that encryption has not been suppressed by an attacker.

#### 8.4.2 Eavesdropping on user data by suppression of encryption between the target user and the true network

*Description:*

An attack that requires a *modified BS/MS* and that exploits the weakness that the network cannot authenticate messages received over the radio interface. The target user is enticed to camp on the false BS/MS. When the target user or the genuine network sets up a connection, the false BS/MS modifies the ciphering capabilities of the MS to make it appear to the network that a genuine incompatibility exists between the network and the mobile station. The network may then decide to establish an un-enciphered connection. After the decision not to cipher has been taken, the intruder may eavesdrop on the user data.

*Does 3G security architecture counteract the attack: Yes*

Message authentication and replay inhibition of the mobile's ciphering capabilities allows the network to verify that encryption has not been suppressed by an attacker.

#### 8.4.3 Eavesdropping on user data by forcing the use of a compromised cipher key

*Description:*

An attack that requires a *modified BS/MS* and the possession by the intruder of a *compromised authentication vector* and thus exploits the weakness that the user has no control the cipher key. The target user is enticed to camp on the false BS/MS. When the target user or the intruder set-up a service, the false BS/MS forces the use of a compromised cipher key on the mobile user while it builds up a connection with the genuine network using its own subscription.

*Does 3G security architecture counteract the attack: Yes*

The presence of a sequence number in the challenge allows the USIM to verify the freshness of the cipher key to help guard against forced re-use of a compromised authentication vector. However, the architecture does not protect against force use of compromised authentication vectors, which have not yet been used to authenticate the USIM. Thus, the network is still vulnerable to attacks using compromised authentication vectors, which have been intercepted between generation in the authentication centre and use and destruction in the serving network.

The user must trust the SN (transitively via the HE) to handle authentication vectors securely. For instance, an attacker with a false BS may work in collusion with an SN to intercept unused authentication vectors, or the SN may expose itself to undue risks because it stockpiles large numbers of authentication vectors before they need to be used.

## 8.5 Impersonation of the user

### 8.5.1 Impersonation of the user through the use of by the network of a compromised authentication vector

*Description:*

An attack that requires a *modified MS* and the possession by the intruder of a *compromised authentication vector* which is intended to be used by the network to authenticate a legitimate user. The intruder uses that data to impersonate the target user towards the network and the other party.

*Does 3G security architecture counteract the attack: Yes*

The presence of a sequence number in the challenge means that authentication vectors cannot be re-used to authenticate USIMs. This helps to reduce the opportunity of using a compromised authentication vector to impersonate the target user. However, the network is still vulnerable to attacks using compromised authentication vectors, which have been intercepted between generation in the authentication centre and use and destruction in the serving network.

The user must trust the SN (transitively via the HE) to handle authentication vectors securely. For instance, an attacker with a false BS may work in collusion with an SN to intercept unused authentication vectors, or the SN may expose itself to undue risks because it stockpiles large numbers of authentication vectors before they need to be used.

### 8.5.2 Impersonation of the user through the use by the network of an eavesdropped authentication response

*Description:*

An attack that requires a *modified MS* and exploits the weakness that an authentication vector may be used several times. The intruder eavesdrops on the authentication response sent by the user and uses that when the same challenge is sent later on. Subsequently, ciphering has to be avoided by any of the mechanisms described above. The intruder uses the eavesdropped response data to impersonate the target user towards the network and the other party.

*Does 3G security architecture counteract the attack: Yes*

The presence of a sequence number in the challenge means that authentication vectors cannot be re-used to authenticate USIMs.

### 8.5.3 Hijacking outgoing calls in networks with encryption disabled

*Description:*

This attack requires a *modified BS/MS*. While the target user camps on the false base station, the intruder pages the target user for an incoming call. The user then initiates the call set-up procedure, which the intruder allows to occur between the serving network and the target user, modifying the signalling elements such that for the serving network it appears as if the target user wants to set-up a mobile originated call. The network does not enable encryption. After authentication the intruder cuts the connection with the target

user, and subsequently uses the connection with the network to make fraudulent calls on the target user's subscription.

*Does 3G security architecture counteract the attack: Partly*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the connection set-up request allows the serving network to verify that the request is legitimate. In addition, periodic integrity protected messages during a connection helps protect against hijacking of un-enciphered connections after the initial connection establishment. However, hijacking the channel between periodic integrity protection messages is still possible, although this may be of limited use to attackers. In general, connections with ciphering disabled will always be vulnerable to some degree of channel hijacking.

#### **8.5.4 Hijacking outgoing calls in networks with encryption enabled**

*Description:*

This attack requires a *modified BS/MS*. In addition to the previous attack this time the intruder has to attempt to suppress encryption by modification of the message in which the MS informs the network of its ciphering capabilities.

*Does 3G security architecture counteract the attack: Yes*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the MS station classmark and the connection set-up request helps prevent suppression of encryption and allows the serving network to verify that the request is legitimate.

#### **8.5.5 Hijacking incoming calls in networks with encryption disabled**

*Description:*

This attack requires a *modified BS/MS*. While the target user camps on the false base station, an associate of the intruder makes a call to the target user's number. The intruder acts as a relay between the network and the target user until authentication and call set-up has been performed between target user and serving network. The network does not enable encryption. After authentication and call set-up the intruder releases the target user, and subsequently uses the connection to answer the call made by his associate. The target user will have to pay for the roaming leg.

*Does 3G security architecture counteract the attack: Partly*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the connection accept message allows the serving network to verify that the request is legitimate. In addition, periodic integrity protected messages during a connection helps protect against hijacking of un-enciphered connections after the initial connection establishment. However, hijacking the channel between periodic integrity protection messages is still possible, although this may be of limited use to attackers. In general, connections with ciphering disabled will always be vulnerable to some degree of channel hijacking.

#### **8.5.6 Hijacking incoming calls in networks with encryption enabled**

*Description:*

This attack requires a *modified BS/MS*. In addition to the previous attack this time the intruder has to suppress encryption.

*Does 3G security architecture counteract the attack: Yes*

Integrity protection of critical signalling messages protects against this attack. More specifically, data authentication and replay inhibition of the MS station classmark and the connection accept message helps prevent suppression of encryption and allows the serving network to verify that the connection accept is legitimate.

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## 9 Network issues

### 9.1 Security policy

#### 9.1.1 Access control policy

Access control policy with respect to 3GPP network elements should be consistent with general access control policy as defined in the particular operator's security policy. As a basis, the following rules should apply:

1. In granting users access rights to 3GPP networks elements or supporting IT systems the following principles should be followed:
  - every employee should only have access to those resources necessary for the completion of the work-related tasks set,
  - the "positive access control" principle should be applied, meaning it shall be assumed that an employee is authorised to carry out only those operations for which he has obtained authority,
  - The right of access to resources should be granted only at the moment when it is actually necessary and should be rescinded when no longer necessary for the completion of work-related tasks.
2. Operator's employees should be made responsible for the secure storing and use of access control executive components entrusted to them (badges, cards). Access control executive components should not be stored together with a computer used to access the network element or IT system.
3. Every user of a given system should be provided with an identification (log-in name, account name) that is unique within the framework or the Company. The following principles apply:
  - a user's identification on its own should not be sufficient for granting access authority,
  - an identification should not give any indication of the user's authority within the system,
  - The use of forms of group identification should only be admissible in exceptional circumstances.

Granting of full or very wide rights of access to resources should be limited and strictly controlled.

## 9.2 Secure network elements interconnection

3GPP network elements must provide means for remote management, maintenance and communication with IT systems (e.g. the billing system). Often an operator's corporate computer network is used for this purpose. This considerably lower infrastructure costs but poses significant security threats for 3GPP system entities. If no security is applied, usually each user of corporate network can try to access remotely a 3GPP network element, provided its network address is known.

As a principle, 3GPP network elements should be separated, at least logically, from an operator's corporate computer network. A unique username and password should identify each employee who is authorised to access to network element. Proper application and system logs should be maintained, reviewed and protected.

Remote access to network entities should be, subject to the operator's security policy, protected from eavesdropping and session hijacking.

Physical access to 3GPP network elements should be controlled by appropriate physical security measures. It is advisable that physical location of network elements be treated as protected information.

## 9.3 Communications node security

To countermeasure the threats described in this document an operator should define and implement proper security measures. The following section specifies the desirable security features that any 3GPP Network Element (NE), Network System (NS), Operations System (OS) or Data Communications Network (DCN) should provide in order to reduce the risk of potentially service affecting security compromises. The term "3GPP node" in the following section is used to imply a NE, NS, OS, or a DCN and its nodes.

### 9.3.1 Identification

Each operations related process running in the 3GPP node should be associated with the corresponding user-ID (so that an audit trail can be established if there is a need).

The 3GPP node should disable a user-ID if it has remained inactive (i.e., never used) over a specified time period.

### 9.3.2 Authentication

All Operations, Administration, Maintenance and Provisioning (OAM&P) input ports of the 3GPP node (including direct, dial-up and network access) should require authentication of a session requester, without any provision for a bypass mechanism.

A single stored password entry (e.g., in a password file) should not be allowed to be shared by multiple user-IDs. However, the 3GPP node should not prevent a user from choosing (unknowingly) a password that is already being used by some other user. Nor should the 3GPP node volunteer this information to either user.



Passwords should be stored in a one-way encrypted form, and should not be retrievable by any user including managers or administrators (of system and security). Also, there should be no clear text display (on a device such as a screen, typewriter, or printer) of a password at any time (e.g., login, file dump, etc.).

The 3GPP node should allow passwords to be user changeable (requiring re-authentication), and should require that the user change it the first time he/she establishes a session with the password assigned to him/her. The default should be non-trivial in nature, ideally random.

The password should have an "ageing" feature, and it should have a complexity requirement to make it not easily guessed. The 3GPP node should not accept common words or names as valid passwords. Also, it should not allow a recently obsolete password to be readily reselected by the said user.

### 9.3.3 System Access Control

The 3GPP node should not allow access to any session requester unless identified and authenticated. There should be no default mechanism to circumvent it.

The 3GPP node should not allow any session to be established via a port that is not authorised to accept input commands. For example, if an output port receives a login request, the 3GPP node should not respond.

The entire login procedure should be allowed to be completed without interruption, even if incorrect parameters (such as an incorrect user-ID or an incorrect password) are entered, and no "help message" should be transmitted to the session requester as to whom part of the authentication is incorrect. The only information to be conveyed at the end of the login attempt is that the login is invalid.

After a specified number of incorrect login attempts carried out in succession, the 3GPP node should lock out the channel and raise an alarm in real time for the administrator.

Before the session begins, the 3GPP node should provide a warning message explicitly alerting the user of the consequences of unauthorised access and use.

At the beginning of the session, the 3GPP node should display the date and time of the user's last successful access and the number of unsuccessful attempts, if any, that have been made to establish a session since the last successful access.

There should be a "time-out" feature - i.e., the 3GPP node should disconnect or re-authenticated users after a specified time interval during which no messages were exchanged. Also, there should be a mechanism for user-initiated keyboard locking.

The 3GPP node should provide a mechanism to end a session through a secure logoff procedure. If a session gets interrupted due to reasons such as time-out, power failure, link disconnection, etc., the port should be dropped immediately.

For dial-up access over untrusted channels, authentication involving one time passwords should be required (e.g., smart card, etc.).

### 9.3.4 Resource Access Control

Access to resources should be controlled on the basis of “privilege” (i.e., access permission) associated with user-ID and channel. It should not be based on a “password” associated with the access function, because that password will have to be necessarily shared among all users requiring such access. Neither should encryption be used as a primary access control mechanism (though encryption may be used to enhance it).

The granularity of resource access control should be such that for each resource it should be possible to grant (or deny) access privilege to any single user (or a prescribed group of users). For example, the control should be adequately fine-grained so that user access and channel access can be restricted on the basis of commands, database views (i.e., objects), records (i.e., object instances), and fields (i.e., attributes).

If external entities - e.g., customers, are allowed access to the resources, each 3GPP node's resource (e.g., proprietary data) should be protected from access by unauthorised persons.

Executable/loadable/fetchable software should be access controlled for overwrite, update, and execution rights.

### 9.3.5 Accountability and Audit

The 3GPP node should generate a security log containing information sufficient for after-the-fact investigation of loss or impropriety.

The security log should be protected from unauthorized access. No user should be allowed to modify or delete a security log. There should be no mechanism to disable the security log. There should be an alarm in real time if the security log does not function properly.

The security log should, as a minimum, record events such as:

- all sessions established,
- invalid user authentication attempts,
- unauthorized attempts to access resources (including data and transactions),
- changes in users' security profiles and attributes,
- changes in access rights to resources,
- changes in the 3GPP node security configuration,
- And modification of 3GPP node software.

For each such event, the record should, as a minimum, include date and time of event, initiator of the event such as: user-ID, terminal, port, network address, etc., names of resources accessed, and success or failure of the event.

Actual or attempted passwords should not be recorded in the security log

There should be audit tools to produce exception reports, summary reports, and detailed reports on specifiable data items, users, or communication facilities.

### 9.3.6 Security Administration

The 3GPP node should support functions for the “management” of security related data (e.g., security parameters such as user-IDs, passwords, privileges, etc.) as “separate”

from other user functions. Security administration should be reserved only for an appropriate administrator.

The administrator should be able to display all currently logged-in users as well as a list of all authorised user-IDs.

The administrator should be able to independently and selectively monitor, in real time, the actions of any one or more users based on respective user-IDs, terminals, ports, or network addresses.

The administrator should be able to identify all resources owned by or accessible to any specific user along with the associated access privileges.

The administrator should be able to enter, edit, delete or retrieve all attributes of a user-ID (except for a password, which should not be retrievable).

The administrator should limit the use of a "null password" during system login on a per user or per port basis (i.e., during new release installation).

The administrator should be able to save the security log for safe storage, so that it is not written over when the buffer is full.

All security parameters (e.g., password-ageing interval, time-out interval, and various alarm conditions) should be specifiable and adjustable by the administrator. This implies that the 3GPP node should not have any security parameters hard coded.

### 9.3.7 Documentation

Any 3GPP node supplier/vendor should provide documentation on security considerations for administrators, operators, and users. They can be stand-alone documents or sections incorporated in appropriate vendor manuals.

The administrator's guide should contain items such as: functions and privileges that need to be controlled to secure the facility, proper usage of security audit tools, procedures for examining and maintaining audit files, procedures for periodic saving and backup of security logs, recommendations on setting the minimum access permissions on all files, directories, and databases, guidelines on security assessment techniques.

The operator's guide should contain procedures necessary to initially start the 3GPP node in a secure manner and to resume secure operation after any lapse that may have occurred.

The user's guide should describe the protection mechanisms that are non-transparent to the user, should explain their purpose, and provide guidelines on their use. It should not contain any information that could jeopardise the security of the 3GPP node if made public.

Passwords should be stored in a one-way encrypted form, and should not be retrievable by any user including managers or administrators (of system and security). Also, there should be no clear text display (on a device such as a screen, typewriter, or printer) of a password at any time (e.g., login, file dump, etc.).

The 3GPP node should allow passwords to be user changeable (requiring reauthentication), and should require that the user change it the first time he/she establishes a session with the password assigned to him/her. The default should be non-trivial in nature, ideally random.

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## 10 Inter Network Security

### 10.1 Signalling system Number 7

Mobile networks primarily use Signaling System no. 7 (SS7) for communication between networks for such activities as authentication, location update, and supplementary services and call control. The messages unique to 3GPP are MAP messages.

The security of the global SS7 network as a transport system for signaling messages e.g. authentication and supplementary services such as call forwarding is open to major compromise.

The problem with the current SS7 system is that messages can be altered, injected or deleted into the global SS7 networks in an uncontrolled manner.

In the past, SS7 traffic was passed between major PTO's covered under treaty organization and the number of operators was relatively small and the risk of compromise was low.

Networks are getting smaller and more numerous. Opportunities for unintentional mishaps will increase, as will the opportunities for hackers and other abusers of networks.

With the increase in different types of operators and the increase in the number of interconnection circuits there is an ever-growing loss of control of security of the signaling networks.

There is also exponential growth in the use of interconnection between the telecommunication networks and the Internet .The IT community now has many protocol converters for conversion of SS7 data to IP, primarily for the transportation of voice and data over the IP networks. In addition new services such as those based on IN will lead to a growing use of the SS7 network for general data transfers.

There have been a number of incidents from accidental action, which have damaged a network. To date, there have been very few deliberate actions.

The availability of cheap PC based equipment that can be used to access networks and the ready availability of access gateways on the Internet will lead to compromise of SS7 signaling and this will effect mobile operators.

The risk of attack has been recognised in the USA at the highest level of the President's office indicating concern on SS7. It is understood that the T1, an American group is seriously considering the issue.

For the network operator there is some policing of incoming signaling on most switches already, but this is dependent on the make of switch as well as on the way the switch is configured by operators.

Some engineering equipment is not substantially different from other advanced protocol analysers in terms of its fraud potential, but is more intelligent and can be programmed more easily.

The SS7 network as presently engineered is insecure. It is vitally important that network operators ensure that signaling screening of SS7 incoming messages takes place at the entry points to their networks and that operations and maintenance systems alert against unusual SS7 messages. There are a number of messages that can have a significant effect on the operation of the network and inappropriate messages should be controlled at entry point.

Network operators network security engineers should on a regular basis carry out monitoring of signaling links for these inappropriate messages. In signing agreements with roaming partners and carrying out roaming testing, review of messages and also to seek appropriate confirmation that network operators are also screening incoming SS7 messages their networks to ensure that no rouge messages appear.

In summary there is no adequate security left in SS7. Mobile operators need to protect them selves from attack from hackers and inadvertent action that could stop a network or networks operating correctly.

Operators should note that HPLMN control over a subscriber roaming in a VPLMN using different MAP release could be limited. To avoid this, operators should assure that their roaming partners use the current MAP version, as specified by the 3GPP Association.

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## 11 Intra network security

### 11.1 3GPP Network elements and interfaces

Unauthorised, local or remote access to 3GPP network elements can result in access to confidential data stored by system entities, unauthorised access to services and resources, misuse of the network element to gain access to data or services or denial of service. The following section gives an outline of potential threats related to attacks on 3GPP network elements and recommendations.

#### 11.1.1 Home Location Register - HLR

An unauthorised access to HLR could result in activating subscribers not seen by the billing system, thus not chargeable. Services may also be activated or deactivated for each subscriber, thus allowing unauthorised access to services or denial of service attacks. In certain circumstances it is possible to use Man-Machine (MM) commands to monitor other HLR user's action - this would also often allow for unauthorised access to data.

An operator should not rely on the fact that an intruder's knowledge on particular vendor's MM language will be limited. Those attacks can be performed both by external intruders and by operator's employees.

Access control to HLRs should be based on user profiles, using at least a unique username and a password as authentication data. Remote access to HLR should be protected from eavesdropping, source and destination spoofing and session hijacking. An operator may therefore wish to limit the range of protocols available for communication with HLR..

#### 11.1.2 Authentication Centre - AuC

*An intruder who gains direct access to an AuC can effectively clone all subscribers whose data he had access to.*

Number of employees having physical and logical access to AuC should be limited. From security point of view it is then reasonable to use an AuC which is not integrated with HLR.

Operators should carefully consider the need for encryption of AuC data. Some vendors use default encryption, the algorithm being proprietary and confidential. It should be noted that strength of such encryption could be questionable.

If decided to use an add-on ciphering facility, attention should be paid to cryptographic key management. Careless use of such equipment could even lower AuC security.

Authentication triplets can be obtained from AuC by masquerading as another system entity (namely HLR). The threat is present when HLR and AuC are physically separated.

### 11.1.3 Mobile Switching Centre - MSC

An MSC is one of the most important nodes of any 3GPP network. It handles all calls incoming to, or originating from subscribers visiting the given switch area. Unauthorised, local or remote, access to an MSC would likely result in the loss of confidentiality of user data, unauthorised access to services or denial of service for large numbers of subscribers.

It is strongly recommended that access to MSCs is restricted, both in terms of physical and logical access. It is also recommended that their physical location is not made public.

When co-located, several MSCs should be independent (i.e. separated power, transmission,) in order to limit the impacts from accidents on one particular MSC (e.g. fire).

### 11.1.4 3GPP network interfaces

An intruder gaining access to 3GPP network interfaces would primary gain access to information sent on the interface targeted. However, playing denial-of-service attacks would also be feasible - dependent on how the interface is technically realised (e.g. cable or wireless).

Telecommunication networks are usually designed with necessary redundancy, allowing for reconfiguration in case of loss of a link or links. From security point of view it is particularly important to foresee alternate connection paths where links vulnerable to denial-of-service attacks (e.g., microwave links susceptible to jamming) are in use.

### 11.1.5 Billing system / Customer Care system

Billing/customer care systems are critical for maintaining the business continuity of a 3GPP Operator.

Unauthorised access to the billing or customer care system could result in

- loss of revenue due to manipulated CDRs (on the mediation device/billing system level)
- unauthorised applying of service discounts (customer care system level), unauthorised access to services (false subscriptions)
- and even denial of service - by repeated launching of resource - consuming system jobs.

Attention should be paid to the fact that access rights to the billing/customer care system are often granted to temporary employees.

As 3GPP network operators should introduce proper access control mechanisms, coherent with the Operator's general security policy. In particular, it would be advisable to:

- Control the access to the billing data on the database level.
- All users of the billing system should be authenticated by the billing database and access rights should be granted by the database upon successful authentication. Relying on the application-to-database authentication leaves the database open for a skilled attacker.
- Review the activation process.

The same employee should not carry out both tasks; data verification should involve a trusted employee. Activation should be made only upon confirmation of the person verifying the data entered.

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## 12 User Module and Smart Card

If a 3GPP SIM is integrated on a multi-application smart card, there should be sufficient guarantees that the Ki cannot be read or used by any application other than the 3GPP application. Also there should be clear and secure procedures for placing applications and information on the smart card, ensuring that 3GPP information cannot be changed in an unauthorised way. There should be clear responsibilities and procedures for dealing with stolen or malfunctioning cards.

The importance of secure management of Ki's is already detailed above. In addition it is important that SIM status lists are kept up to date and that operators define measures to detect and investigate the misuse of SIMs. There should be procedures to replace SIMs, for example at the end of their validity period, and to deal with stolen SIMs. It is particularly important that individual operators devise and operate secure SIM management processes with their SIM suppliers and throughout the SIM distribution channel.

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## 13 Algorithms

### 13.1 Authentication algorithm

3GPP does not define a standard authentication algorithm, allowing operators to choose their own versions, which comply with the published standards. However, in order to help operators guidelines are available as to how to develop a suitable algorithm. The authentication algorithm is contained within the smart card.

The individual key for each IMSI must be chosen to be random, and must be protected in order to prevent the user from being duplicated. Throughout the security process Ki should be protected.

### 13.2 Confidentiality algorithm

3GPP defines a standard confidentiality algorithm, which is contained within all mobiles, and protects user data from the mobile to the serving node. This is not only over the radio path as in GSM, but also continues back over the links to the serving node.

The confidentiality algorithm, called Kasumi, is expected to be published.

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## 14 Services

There are many value-added services within the ETSI standards, which will sometimes, when wrongly implemented or interpreted, can be used for fraud.

For example, call forwarding can be set which will then allow calls made to a mobile to be sent to expensive destination numbers. This could be done, for example, by ringing a



mobile customer and getting them to put in a call forward number themselves by persuading them that they are testing the mobile.

Many other similar problems exist, such as follow-me services, voicemail, and explicit call transfer. It is to expected that as the services offered by 3GPP become more complex (and include for example Internet connectivity, packet data services as well as MExE which runs code on the mobile, and Java multi application smart cards) then the problem can only become worse.

Operators should ensure that they look carefully at every new network feature and service product to ensure that such problems will not occur in their networks.

## 14.1 Location services

The location service feature in 3GPP depends on the accuracy of the mechanism used within the mobile equipment. It cannot be though of as accurate, as the mobile software can be modified, or the GPS (Global Positioning System by Satellite) could be displaced by a differential input.

## 14.2 Mobile Execution Environment - MExE

The ability to remotely modify remote and run code on a mobile clearly introduces a security risk. In the case of MExE it is up to the user to determine if a possible security risk is introduced, and stop the action from taking place. It is to be expected that a smart attacker will be able to introduce code that will fool a user into setting up services or connection that will compromise them or result them in losing money.

# 15 Index

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## 16History

<b>Document history</b>		
1.0.0	Oct 1999	Publication as first draft to 3GPP TSG SA WG3 Security
1.1.0	Nov 1999	Presented at No 6 for information
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