

**3GPP TSG RAN WG1 Meeting #9
Dresden, Germany, 30 Nov - 3 Dec 1999**

Document R1-99k05

e.g. for 3GPP use the format TP-99xxx
or for SMG, use the format P-99-xxx

<h2 style="margin: 0;">CHANGE REQUEST</h2>		<i>Please see embedded help file at the bottom of this page for instructions on how to fill in this form correctly.</i>
25.213	CR	005r1
GSM (AA.BB) or 3G (AA.BBB) specification number ↑		Current Version: v3.0.0
↑ CR number as allocated by MCC support team		
For submission to: RAN #6 <i>list expected approval meeting # here ↑</i>	for approval for information	strategic <input type="checkbox"/> non-strategic <input type="checkbox"/> <i>(for SMG use only)</i>
Form: CR cover sheet, version 2 for 3GPP and SMG		The latest version of this form is available from: ftp://ftp.3gpp.org/Information/CR-Form-v2.doc

Proposed change affects: (U)SIM ME UTRAN / Radio Core Network
(at least one should be marked with an X)

Source: Nokia **Date:** 1 Nov 1999

Subject: Harmonization of notations for downlink scrambling codes

Work item:

Category:	F Correction <input type="checkbox"/> A Corresponds to a correction in an earlier release <input type="checkbox"/> B Addition of feature <input type="checkbox"/> C Functional modification of feature <input type="checkbox"/> D Editorial modification <input checked="" type="checkbox"/>	Release:	Phase 2 <input type="checkbox"/> Release 96 <input type="checkbox"/> Release 97 <input type="checkbox"/> Release 98 <input type="checkbox"/> Release 99 <input checked="" type="checkbox"/> Release 00 <input type="checkbox"/>
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(only one category shall be marked with an X)

Reason for change: The current text in 5.2.2 is misleading, notation is clarified.

Clauses affected: 5.2.2 of TS25.213

Other specs affected:	Other 3G core specifications <input type="checkbox"/> Other GSM core specifications <input type="checkbox"/> MS test specifications <input type="checkbox"/> BSS test specifications <input type="checkbox"/> O&M specifications <input type="checkbox"/>	→ List of CRs: → List of CRs: → List of CRs: → List of CRs: → List of CRs:	
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Other comments:

<----- double-click here for help and instructions on how to create a CR.

In case the OVSF code on the PDSCH varies from frame to frame, the OVSF codes shall be allocated such a way that the OVSF code(s) below the smallest spreading factor will be from the branch of the code tree pointed by the smallest spreading factor used for the connection. This means that all the codes for UE for the PDSCH connection can be generated according to the OVSF code generation principle from smallest spreading factor code used by the UE on PDSCH.

In case of mapping the DSCH to multiple parallel PDSCHs, the same rule applies, but all of the branches identified by the multiple codes, corresponding to the smallest spreading factor, may be used for higher spreading factor allocation.

5.2.2 Scrambling code

A total of $2^{18}-1 = 262,143$ scrambling codes, numbered $0 \dots 262,142$ can be generated. However not all the scrambling codes are used. The scrambling codes are divided into 512 sets each of a primary scrambling code and 15 secondary scrambling codes.

The primary scrambling codes consist of scrambling codes $n=16*i$ where $i=0 \dots 511$. The i :th set of secondary scrambling codes consists of scrambling codes $16*i+k$, where $k=1 \dots 15$.

There is a one-to-one mapping between each primary scrambling code and 15 secondary scrambling codes in a set such that i :th primary scrambling code corresponds to i :th set of scrambling codes.

Hence, according to the above, scrambling codes $k = 0, 1, \dots, 8191$ are used. Each of these codes are associated with an even alternative scrambling code and an odd alternative scrambling code, that may be used for compressed frames. The even alternative scrambling code corresponding to scrambling code k is scrambling code number $k + 8192$, while the odd alternative scrambling code corresponding to scrambling code k is scrambling code number $k + 16384$.

The set of primary scrambling codes is further divided into 64 scrambling code groups, each consisting of 8 primary scrambling codes. The j :th scrambling code group consists of primary scrambling codes $16*8*j+16*k$, where $j=0 \dots 63$ and $k=0 \dots 7$.

Each cell is allocated one and only one primary scrambling code. The primary CCPCCH [and primary CPICH are](#) always transmitted using the primary scrambling code. The other downlink physical channels can be transmitted with either the primary scrambling code or a secondary scrambling code from the set associated with the primary scrambling code of the cell.

The mixture of primary scrambling code and secondary scrambling code for one CCTrCH is allowable.

The scrambling code sequences are constructed by combining two real sequences into a complex sequence. Each of the two real sequences are constructed as the position wise modulo 2 sum of 38400 chip segments of two binary m -sequences generated by means of two generator polynomials of degree 18. The resulting sequences thus constitute segments of a set of Gold sequences. The scrambling codes are repeated for every 10 ms radio frame. Let x and y be the two sequences respectively. The x sequence is constructed using the primitive (over GF(2)) polynomial $1+X^7+X^{18}$. The y sequence is constructed using the polynomial $1+X^5+X^7+X^{10}+X^{18}$.

The sequence depending on the chosen scrambling code number n is denoted z_n , in the sequel. Furthermore, let $x(i)$, $y(i)$ and $z_n(i)$ denote the i :th symbol of the sequence x , y , and z_n , respectively

The m -sequences x and y are constructed as:

Initial conditions:

x is constructed with $x(0)=1, x(1)=x(2)=\dots=x(16)=x(17)=0$

$y(0)=y(1)=\dots=y(16)=y(17)=1$

Recursive definition of subsequent symbols:

$x(i+18) = x(i+7) + x(i) \text{ modulo } 2, i=0, \dots, 2^{18}-20,$

$y(i+18) = y(i+10)+y(i+7)+y(i+5)+y(i) \text{ modulo } 2, i=0, \dots, 2^{18}-20.$

The n:th Gold code sequence $z_n, n=0,1,2,\dots,2^{18}-2$, is then defined as

$$z_n(i) = x(i+n) \text{ modulo } (2^{18} - 1) + y(i) \text{ modulo } 2, i=0,\dots, 2^{18}-2.$$

These binary sequences code words are converted to real valued sequences Z_n by the following transformation: $0 \rightarrow +1, 1 \rightarrow -1$.

$$Z_n(i) = \begin{cases} +1 & \text{if } z_n(i) = 0 \\ -1 & \text{if } z_n(i) = 1 \end{cases} \text{ for } i = 0,1,\dots,2^{18} - 2.$$

Finally, the n:th complex scrambling code sequence $S_{dl,n}$ is defined as (the lowest index corresponding to the chip scrambled first in each radio frame) (where N is the period in chips and M is 131,072):

$$S_{dl,n}(i) = Z_n(i) + j Z_n((i + 131072M) \text{ modulo } (2^{18} - 1)), i=0,1,\dots,38399N-1.$$

Note that the pattern from phase 0 up to the phase of 38399 is repeated.

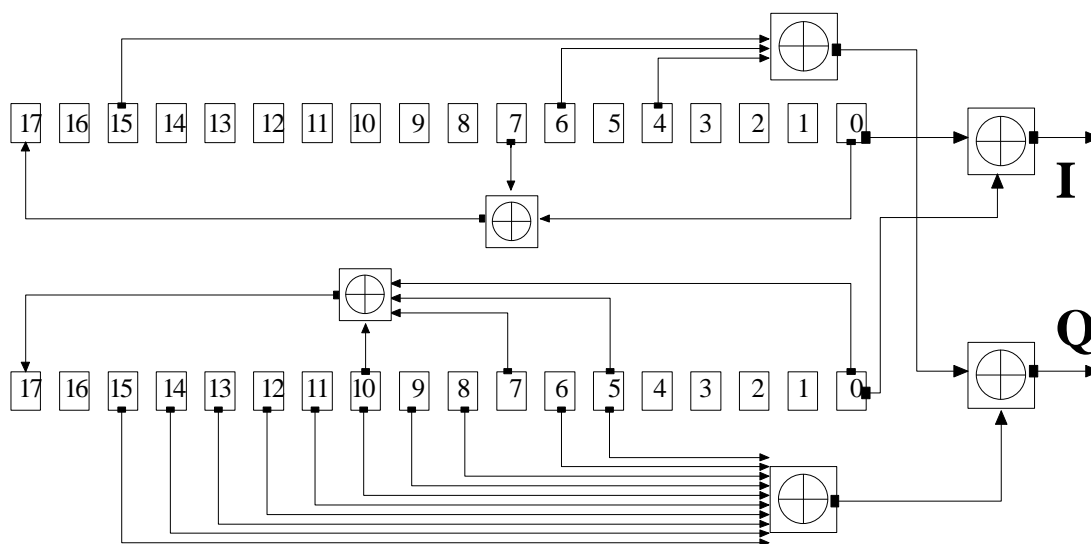


Figure 11: Configuration of downlink scrambling code generator